

Companions and conjoints in a virtual equipment

Jana K. Nickel, 17 February 2026

In the following notes, we will introduce companions and conjoints in a virtual equipment and state and prove some related results in great detail.

The notes are mainly based on:

- Nathanael Arkor and Dylan McDermott: The formal theory of relative monads, 2025.
- G.S.H Cruttwell and Michael A. Shulman: A unified framework for generalized multicategories, 2010.

Yet, I will present the material in a slightly different style and with many additional details.

Construction.

Let X be a virtual equipment.

Let $f: A \rightarrow B$ be any tight arrow in X .

1) Then we have two niche configurations in X

$$\begin{array}{ccc} A & & B \\ f \downarrow & & \parallel \\ B & \xrightarrow{e_B} & B \end{array} \quad \text{and} \quad \begin{array}{ccc} B & & A \\ \parallel & & \downarrow f \\ B & \xrightarrow{e_B} & B \end{array}$$

By assumption, they come along with cartesian 2-cells

$$\begin{array}{ccc} A & \xrightarrow{B(1,f)} & B \\ f \downarrow & \bar{B}(1,f) \parallel & \\ B & \xrightarrow{e_B} & B \end{array} \quad \text{and} \quad \begin{array}{ccc} B & \xrightarrow{B(f,1)} & A \\ \parallel & \bar{B}(f,1) \downarrow & \\ B & \xrightarrow{e_B} & B \end{array}$$

respectively.

We call the two loose arrows $B(1,f): A \rightarrow B$ and $B(f,1): B \rightarrow A$ in X also the base change loose arrows of the tight arrow $f: A \rightarrow B$.

2) Consider the identity 2-cell e_f of f in X :

$$\begin{array}{ccc} A & \xrightarrow{e_A} & A \\ f \downarrow & e_f & \downarrow f \\ B & \xrightarrow{e_B} & B \end{array}$$

This is the unique 2-cell in X of the form

$$\begin{array}{ccc} A & \xrightarrow{e_A} & A \\ f \downarrow & e_f & \downarrow f \\ B & \xrightarrow{e_B} & B \end{array}$$

satisfying $e_B \circ (f) = e_f \circ e_A$,

i.e.,

$$\begin{array}{ccc} A & & \\ f \downarrow & & \\ B & \xrightarrow{e_B} & B \\ \text{opcart} & & \end{array} = \begin{array}{ccc} A & \xrightarrow{e_A} & A \\ \text{opcart} & & \\ A & \xrightarrow{e_A} & A \\ f \downarrow & e_f & \downarrow f \\ B & \xrightarrow{e_B} & B \end{array}$$

2.1) We can write the 2-cell e_f as

$$\begin{array}{ccc}
 A \xrightarrow{e_A} A & & A \xrightarrow{e_A} A \\
 f \downarrow e_f \downarrow f & = & \parallel \downarrow f \\
 B \xrightarrow{e_B} B & & A \xrightarrow{e_f} B \\
 & & f \downarrow \parallel \\
 & & B \xrightarrow{e_B} B
 \end{array}$$

so by the universal property of the cartesian 2-cell

$$\begin{array}{ccc}
 A \xrightarrow{B(\lambda, f)} B & & \text{applied to } e_f, \\
 f \downarrow \bar{B}(\lambda, f) \parallel & & \\
 B \xrightarrow{e_B} B & &
 \end{array}$$

$\exists!$ 2-cell $v(\lambda, f)$ in \mathcal{X} of the form

$$\begin{array}{ccc}
 A \xrightarrow{e_A} A & & \\
 \parallel v(\lambda, f) \downarrow f & & \\
 A \xrightarrow{B(\lambda, f)} B & &
 \end{array}$$

with $e_f = \bar{B}(\lambda, f) \circ v(\lambda, f)$, i.e.,

$$\begin{array}{ccc}
 A \xrightarrow{e_A} A & & A \xrightarrow{e_A} A \\
 f \downarrow e_f \downarrow f & = & \parallel v(\lambda, f) \downarrow f \\
 B \xrightarrow{e_B} B & & A \xrightarrow{B(\lambda, f)} B \\
 & & f \downarrow \bar{B}(\lambda, f) \parallel \\
 & & B \xrightarrow{e_B} B
 \end{array}$$

2.2) Similarly, we can write the 2-cell e_f as

$$\begin{array}{ccc}
 A \xrightarrow{e_A} A & & A \xrightarrow{e_A} A \\
 f \downarrow e_f \downarrow f & = & f \downarrow \parallel \\
 B \xrightarrow{e_B} B & & B \xrightarrow{e_f} A \\
 & & \parallel \downarrow f \\
 & & B \xrightarrow{e_B} B
 \end{array}$$

so by the universal property of the cartesian 2-cell

$$\begin{array}{ccc}
 B \xrightarrow{B(f, \lambda)} A & & \text{applied to } e_f, \\
 \parallel \bar{B}(f, \lambda) \downarrow f & & \\
 B \xrightarrow{e_B} B & &
 \end{array}$$

$\exists!$ 2-cell $v(f, \lambda)$ in \mathcal{X} of the form

$$\begin{array}{ccc}
 A \xrightarrow{e_A} A & & \\
 f \downarrow v(f, \lambda) \parallel & & \\
 B \xrightarrow{B(f, \lambda)} A & &
 \end{array}$$

with $e_f = \bar{B}(f, \lambda) \circ v(f, \lambda)$, i.e.,

$$\begin{array}{ccc}
 A \xrightarrow{e_A} A & & A \xrightarrow{e_A} A \\
 f \downarrow e_f \downarrow f & = & f \downarrow v(f, \lambda) \parallel \\
 B \xrightarrow{e_B} B & & B \xrightarrow{B(f, \lambda)} A \\
 & & \parallel \bar{B}(f, \lambda) \downarrow f \\
 & & B \xrightarrow{e_B} B
 \end{array}$$

Remark.

Let again X be a virtual equipment and $f: A \rightarrow B$ any tight arrow in X .

1) Then we have the following equality of 2-cells in X :

$$\begin{array}{ccc}
 \begin{array}{c}
 A \xrightarrow{e_A} A \xrightarrow{B(\lambda, f)} B \\
 \parallel \nu(\lambda, f) \downarrow f \quad \bar{B}(\lambda, f) \text{ cart} \parallel \\
 A \xrightarrow{B(\lambda, f)} B \xrightarrow{e_B} B \\
 \parallel \lambda_{B(\lambda, f)} \text{ opcart} \parallel \\
 A \xrightarrow{B(\lambda, f)} B
 \end{array} & = & \begin{array}{c}
 A \xrightarrow{e_A} A \xrightarrow{B(\lambda, f)} B \\
 \parallel \rho_{B(\lambda, f)} \text{ opcart} \parallel \\
 A \xrightarrow{B(\lambda, f)} B,
 \end{array}
 \end{array}$$

i.e. $\lambda_{B(\lambda, f)} \circ (\nu(\lambda, f), \bar{B}(\lambda, f)) = \rho_{B(\lambda, f)}$.

Proof.

By the universal property of the cartesian 2-cell

$$\begin{array}{c}
 A \xrightarrow{B(\lambda, f)} B \\
 f \downarrow \quad \bar{B}(\lambda, f) \text{ cart} \parallel \\
 B \xrightarrow{e_B} B,
 \end{array}$$

it suffices to show that the post-compositions of the two above 2-cells with $\bar{B}(\lambda, f)$ agree.

Moreover, by the universal property of the opcartesian 2-cell

$$\begin{array}{c}
 A \\
 \parallel e_A \parallel \\
 A \xrightarrow{e_A} A,
 \end{array}$$

it suffices to show that their pre-compositions with $(e_A, \lambda_{B(\lambda, f)})$ agree, meaning the following two composite 2-cells:

$$\phi := \begin{array}{c}
 \begin{array}{c}
 \text{---} A \xrightarrow{B(\lambda, f)} B \\
 \text{---} \text{---} \text{---} \\
 \text{---} e_A \parallel \text{---} \text{---} \parallel \\
 A \xrightarrow{e_A} A \xrightarrow{B(\lambda, f)} B \\
 \parallel \nu(\lambda, f) \downarrow f \quad \bar{B}(\lambda, f) \parallel \\
 A \xrightarrow{B(\lambda, f)} B \xrightarrow{e_B} B \\
 \parallel \lambda_{B(\lambda, f)} \parallel \\
 A \xrightarrow{B(\lambda, f)} B \\
 f \downarrow \quad \text{---} \text{---} \text{---} \parallel \\
 \text{---} B \xrightarrow{e_B} B \\
 \text{---} \text{---} \text{---} \parallel \\
 \text{---} B \xrightarrow{e_B} B
 \end{array}
 \end{array}$$

and

$$\psi := \begin{array}{c}
 \begin{array}{c}
 \text{---} A \xrightarrow{B(\lambda, f)} B \\
 \text{---} \text{---} \text{---} \\
 \text{---} e_A \parallel \text{---} \text{---} \parallel \\
 A \xrightarrow{e_A} A \xrightarrow{B(\lambda, f)} B \\
 \parallel \rho_{B(\lambda, f)} \parallel \\
 A \xrightarrow{B(\lambda, f)} B \\
 f \downarrow \quad \text{---} \text{---} \text{---} \parallel \\
 \text{---} B \xrightarrow{e_B} B \\
 \text{---} \text{---} \text{---} \parallel \\
 \text{---} B \xrightarrow{e_B} B
 \end{array}
 \end{array}$$

2) Similarly, we have the following equality of 2-cells in X :

$$\begin{array}{ccc}
 B \xrightarrow{B(f,1)} A \xrightarrow{e_A} A & & B \xrightarrow{B(f,1)} A \xrightarrow{e_A} A \\
 \parallel \bar{B}(f,1) \downarrow f \wr(f,1) \parallel & = & \parallel \lambda_{B(f,1)} \parallel \\
 B \xrightarrow{e_B} B \xrightarrow{B(f,1)} A & & B \xrightarrow{B(f,1)} A \\
 \parallel \rho_{B(f,1)} \parallel & & \\
 B \xrightarrow{B(f,1)} A & &
 \end{array}$$

i.e., $\rho_{B(f,1)} \circ (\bar{B}(f,1), \wr(f,1)) = \lambda_{B(f,1)}$.

Proof.

This can be shown by an analogous reasoning as in 1).

Definition.

Let X be a virtual double category with nullary opcartesian 2-cells.

A companion tuple in X is a quadruple (f, p, α, β) consisting of

- a tight arrow $f: A \rightarrow B$ in X ,
- a loose arrow $p: A \rightarrow B$ in X ,
- a 2-cell in X of the form
$$\begin{array}{ccc}
 A & \xrightarrow{e_A} & A \\
 \parallel \alpha \downarrow f & & \\
 A & \xrightarrow{p} & B
 \end{array}$$

- a 2-cell in X of the form
$$\begin{array}{ccc}
 A & \xrightarrow{p} & B \\
 f \downarrow \beta \parallel & & \\
 B & \xrightarrow{e_B} & B
 \end{array}$$

satisfying the following equalities of 2-cells in X :

$$(i) \quad \begin{array}{ccc}
 A \xrightarrow{e_A} A & & A \xrightarrow{e_A} A \\
 \parallel \alpha \downarrow f & = & f \downarrow e_f \downarrow f \\
 A \xrightarrow{p} B & & B \xrightarrow{e_B} B \\
 f \downarrow \beta \parallel & &
 \end{array}$$

i.e., $\beta \circ \alpha = e_f$,

$$(ii) \quad \begin{array}{ccc} A \xrightarrow{e_A} A & \xrightarrow{p} & B \\ \parallel \alpha \downarrow f & & \parallel \\ A \xrightarrow{p} B & \xrightarrow{e_B} & B \\ \parallel & \lambda_p \text{ opcart} & \parallel \\ A & \xrightarrow{p} & B \end{array} = \begin{array}{ccc} A \xrightarrow{e_A} A & \xrightarrow{p} & B \\ \parallel & \rho_p \text{ opcart} & \parallel \\ A & \xrightarrow{p} & B \end{array}$$

i.e., $\lambda_p \circ (\alpha, \beta) = \rho_p$.

In this case, we also call (f, p) a companion pair in \mathcal{X} and say that $p: A \rightrightarrows B$ is a loose companion of $f: A \rightarrow B$ in \mathcal{X} and that $f: A \rightarrow B$ is a tight companion of $p: A \rightrightarrows B$ in \mathcal{X} . We also say that the tight arrow $f: A \rightarrow B$ and the loose arrow $p: A \rightrightarrows B$ are companions of each other in \mathcal{X} .

Definition.

Let \mathcal{X} be a virtual double category with nullary opcartesian 2-cells.

A conjunction in \mathcal{X} is a quadruple (f, p, α, β) consisting of

- a tight arrow $f: A \rightarrow B$ in \mathcal{X} ,
- a loose arrow $p: B \rightrightarrows A$ in \mathcal{X} , *n.b. the order of B and A!*
- a 2-cell in \mathcal{X} of the form
$$\begin{array}{ccc} A & \xrightarrow{e_A} & A \\ f \downarrow & \alpha & \parallel \\ B & \xrightarrow{p} & A \end{array},$$

- a 2-cell in \mathcal{X} of the form
$$\begin{array}{ccc} B & \xrightarrow{p} & A \\ \parallel & \beta & \downarrow f \\ B & \xrightarrow{e_B} & B \end{array}$$

satisfying the following equalities of 2-cells in \mathcal{X} :

$$(i) \quad \begin{array}{ccc} A \xrightarrow{e_A} A \\ f \downarrow \alpha \parallel \\ B \xrightarrow{p} A \\ \parallel \beta \downarrow f \\ B \xrightarrow{e_B} B \end{array} = \begin{array}{ccc} A \xrightarrow{e_A} A \\ f \downarrow e_f \downarrow f \\ B \xrightarrow{e_B} B \end{array}$$

i.e., $\beta \circ \alpha = e_f$,

$$(ii) \quad \begin{array}{ccc} B \xrightarrow{p} A \xrightarrow{e_A} A \\ \parallel \beta \downarrow f \alpha \parallel \\ B \xrightarrow{e_B} B \xrightarrow{p} A \\ \parallel \rho_p \parallel \\ B \xrightarrow{p} A \end{array} = \begin{array}{ccc} B \xrightarrow{p} A \xrightarrow{e_A} A \\ \parallel \lambda_p \parallel \\ B \xrightarrow{p} A \end{array}$$

i.e., $\rho_p \circ (\beta, \alpha) = \lambda_p$.

In this case, we also call (f, p) a conjoint pair in \mathcal{X} and say that $p: B \rightrightarrows A$ is a loose conjoint of $f: A \rightarrow B$ in \mathcal{X} and that $f: A \rightarrow B$ is a tight conjoint of $p: B \rightrightarrows A$ in \mathcal{X} . We also say that the tight arrow $f: A \rightarrow B$ and the loose arrow $p: B \rightrightarrows A$ are conjoints of each other in \mathcal{X} .

Example.

Let \mathcal{X} be a virtual equipment.

Let $f: A \rightarrow B$ be any tight arrow in \mathcal{X} .

Then f admits both a loose companion and a loose conjoint in \mathcal{X} as depicted below.

Specifically, the previous remark shows that the 2-cells $\begin{array}{ccc} A \xrightarrow{e_A} A \\ \parallel v(f, 1) \downarrow f \\ A \xrightarrow{B(f, 1)} B \end{array}$ and $\begin{array}{ccc} A \xrightarrow{B(f, 1)} B \\ f \downarrow \bar{B}(f, 1) \parallel \\ B \xrightarrow{e_B} B \end{array}$

witness $(f, B(f, 1))$ as a companion pair in \mathcal{X} , and the 2-cells $\begin{array}{ccc} A \xrightarrow{e_A} A \\ f \downarrow v(f, 1) \parallel \\ B \xrightarrow{B(f, 1)} A \end{array}$ and $\begin{array}{ccc} B \xrightarrow{B(f, 1)} A \\ \parallel \bar{B}(f, 1) \downarrow f \\ B \xrightarrow{e_B} B \end{array}$

witness $(f, B(f, 1))$ as a conjoint pair in \mathcal{X} .

Theorem.

Let \mathcal{X} be a virtual equipment.

Suppose we are given a niche configuration in \mathcal{X} of the form

$$\begin{array}{ccc} A & & C \\ f \downarrow & & \downarrow g \\ B & \xrightarrow{q} & D \end{array}$$

Then the loose restriction $q(g, p): A \rightrightarrows C$ of f and g along q is also a loose composite of the triple $(B(f, 1), q, D(g, 1))$ of loose arrows in \mathcal{X} as depicted below,

$$A \xrightarrow{B(f, 1)} B \xrightarrow{q} D \xrightarrow{D(g, 1)} C,$$

i.e., \exists opcartesian 2-cell ϕ in \mathcal{X} of shape

$$\begin{array}{ccccc} A & \xrightarrow{B(\lambda, f)} & B & \xrightarrow{q} & D & \xrightarrow{D(g, \lambda)} & C \\ \parallel & & & \text{opcart} & & & \parallel \\ A & \xrightarrow{q(g, f)} & & & & & C \end{array}$$

In particular, the triple $(B(\lambda, f), q, D(g, \lambda))$ admits a loose composite.

Hence, given any loose composite $B(\lambda, f) \circ q \circ D(g, \lambda)$ of $(B(\lambda, f), q, D(g, \lambda))$, then the loose arrows $q(g, f): A \rightarrow C$ and $B(\lambda, f) \circ q \circ D(g, \lambda): A \rightarrow C$ in \mathcal{X} are isomorphic as objects in the category $\mathcal{X}[A, C]$,

$$q(g, f) \cong B(\lambda, f) \circ q \circ D(g, \lambda).$$

Proof.

By assumption, \mathcal{X} is a virtual equipment, so by definition, the niche configuration

$$\begin{array}{ccc} A & & C \\ f \downarrow & & \downarrow g \\ B & \xrightarrow{q} & D \end{array} \text{ in } \mathcal{X} \text{ is equipped with a cartesian 2-cell } \begin{array}{ccc} A & \xrightarrow{q(g, f)} & C \\ f \downarrow & \bar{q}(g, f) \text{ cart} & \downarrow g \\ B & \xrightarrow{q} & D \end{array} \text{ in } \mathcal{X},$$

and so are the niche configurations

$$\begin{array}{ccc} A & & B \\ f \downarrow & & \parallel \\ B & \xrightarrow{e_B} & B \end{array} \text{ and } \begin{array}{ccc} D & & C \\ \parallel & & \downarrow g \\ D & \xrightarrow{e_D} & D \end{array},$$

yielding the cartesian 2-cells

$$\begin{array}{ccc} A & \xrightarrow{B(\lambda, f)} & B \\ f \downarrow & \bar{B}(\lambda, f) \text{ cart} & \parallel \\ B & \xrightarrow{e_B} & B \end{array} \text{ and } \begin{array}{ccc} D & \xrightarrow{D(g, \lambda)} & C \\ \parallel & \bar{D}(g, \lambda) \text{ cart} & \downarrow g \\ D & \xrightarrow{e_D} & D \end{array},$$

Consider the following composite 2-cell in \mathcal{X} :

$$\gamma := \begin{array}{ccccc} A & \xrightarrow{B(\lambda, f)} & B & \xrightarrow{q} & D & \xrightarrow{D(g, \lambda)} & C \\ f \downarrow & \bar{B}(\lambda, f) \text{ cart} & \parallel & = & \parallel & \bar{D}(g, \lambda) \text{ cart} & \downarrow g \\ B & \xrightarrow{e_B} & B & \xrightarrow{q} & D & \xrightarrow{e_D} & D \\ \parallel & = & \parallel & & \lambda_q & & \parallel \\ B & \xrightarrow{e_B} & B & \xrightarrow{q} & D & & \parallel \\ \parallel & & \rho_q & \text{opcart} & & & \parallel \\ B & \xrightarrow{q} & & & & & D \end{array}$$

By the universal property of the cartesian 2-cell

$$\begin{array}{ccc} A & \xrightarrow{q(g, f)} & C \\ f \downarrow & \bar{q}(g, f) \text{ cart} & \downarrow g \\ B & \xrightarrow{q} & D \end{array},$$

$$\begin{array}{c}
 X_0 \xrightarrow{\Gamma_1} X_1 \dots X_{k-1} \xrightarrow{\Gamma_k} X_k = A \xrightarrow{B(\lambda, \mu)} B \xrightarrow{q} D \xrightarrow{D(g, \lambda)} C = Y_0 \xrightarrow{S_1} Y_1 \dots Y_{e-1} \xrightarrow{S_e} Y_e \\
 \parallel \quad \quad \quad \parallel \\
 \begin{array}{c}
 \parallel e_A \parallel \\
 A \xrightarrow{e_A} A \xrightarrow{B(\lambda, \mu)} B \xrightarrow{q} D \xrightarrow{D(g, \lambda)} C \xrightarrow{e_C} C \\
 \parallel \nu(\lambda, \mu) \parallel \downarrow \parallel \nu(g, \lambda) \parallel \\
 A \xrightarrow{B(\lambda, \mu)} B \xrightarrow{e_B} B = D \xrightarrow{e_D} D \xrightarrow{D(g, \lambda)} C = \\
 \parallel \lambda_{B(\lambda, \mu)} \parallel \parallel \parallel \parallel \parallel \parallel \parallel \parallel \\
 \rho_{D(g, \lambda)} \\
 A \xrightarrow{B(\lambda, \mu)} B \xrightarrow{q} D \xrightarrow{D(g, \lambda)} C = Y_0 \xrightarrow{S_1} Y_1 \dots Y_{e-1} \xrightarrow{S_e} Y_e
 \end{array} \\
 \parallel \quad \quad \quad \parallel \\
 X_0 \xrightarrow{\Gamma_1} X_1 \dots X_{k-1} \xrightarrow{\Gamma_k} X_k = A \xrightarrow{B(\lambda, \mu)} B \xrightarrow{q} D \xrightarrow{D(g, \lambda)} C = Y_0 \xrightarrow{S_1} Y_1 \dots Y_{e-1} \xrightarrow{S_e} Y_e \\
 \downarrow h \quad \quad \quad \downarrow \alpha \quad \quad \quad \downarrow k \\
 Z_0 \xrightarrow{t} Z_1
 \end{array}$$

by an earlier remark concerning $\nu(\lambda, \mu)$ and $\nu(g, \lambda)$

$$\begin{array}{c}
 X_0 \xrightarrow{\Gamma_1} X_1 \dots X_{k-1} \xrightarrow{\Gamma_k} X_k = A \xrightarrow{B(\lambda, \mu)} B \xrightarrow{q} D \xrightarrow{D(g, \lambda)} C = Y_0 \xrightarrow{S_1} Y_1 \dots Y_{e-1} \xrightarrow{S_e} Y_e \\
 \parallel \quad \quad \quad \parallel \\
 \begin{array}{c}
 \parallel e_A \parallel \\
 A \xrightarrow{e_A} A \xrightarrow{B(\lambda, \mu)} B \xrightarrow{q} D \xrightarrow{D(g, \lambda)} C \xrightarrow{e_C} C \\
 \parallel \rho_{B(\lambda, \mu)} \parallel \parallel \parallel \parallel \parallel \parallel \parallel \parallel \\
 \lambda_{D(g, \lambda)} \\
 A \xrightarrow{B(\lambda, \mu)} B \xrightarrow{q} D \xrightarrow{D(g, \lambda)} C = Y_0 \xrightarrow{S_1} Y_1 \dots Y_{e-1} \xrightarrow{S_e} Y_e
 \end{array} \\
 \parallel \quad \quad \quad \parallel \\
 X_0 \xrightarrow{\Gamma_1} X_1 \dots X_{k-1} \xrightarrow{\Gamma_k} X_k = A \xrightarrow{B(\lambda, \mu)} B \xrightarrow{q} D \xrightarrow{D(g, \lambda)} C = Y_0 \xrightarrow{S_1} Y_1 \dots Y_{e-1} \xrightarrow{S_e} Y_e \\
 \downarrow h \quad \quad \quad \downarrow \alpha \quad \quad \quad \downarrow k \\
 Z_0 \xrightarrow{t} Z_1
 \end{array}$$

$$\begin{array}{c}
 X_0 \xrightarrow{\Gamma_1} X_1 \dots X_{k-1} \xrightarrow{\Gamma_k} X_k = A \xrightarrow{B(\lambda, \mu)} B \xrightarrow{q} D \xrightarrow{D(g, \lambda)} C = Y_0 \xrightarrow{S_1} Y_1 \dots Y_{e-1} \xrightarrow{S_e} Y_e \\
 \parallel \quad \quad \quad \parallel \\
 X_0 \xrightarrow{\Gamma_1} X_1 \dots X_{k-1} \xrightarrow{\Gamma_k} X_k = A \xrightarrow{B(\lambda, \mu)} B \xrightarrow{q} D \xrightarrow{D(g, \lambda)} C = Y_0 \xrightarrow{S_1} Y_1 \dots Y_{e-1} \xrightarrow{S_e} Y_e \\
 \downarrow h \quad \quad \quad \downarrow \alpha \quad \quad \quad \downarrow k \\
 Z_0 \xrightarrow{t} Z_1
 \end{array}$$

constructions of $\rho_{B(\lambda, \mu)}$ and $\lambda_{D(g, \lambda)}$

$$\begin{array}{c}
 X_0 \xrightarrow{\Gamma_1} X_1 \dots X_{k-1} \xrightarrow{\Gamma_k} X_k = A \xrightarrow{B(\lambda, \mu)} B \xrightarrow{q} D \xrightarrow{D(g, \lambda)} C = Y_0 \xrightarrow{S_1} Y_1 \dots Y_{e-1} \xrightarrow{S_e} Y_e \\
 \downarrow h \quad \quad \quad \downarrow \alpha \quad \quad \quad \downarrow k \\
 Z_0 \xrightarrow{t} Z_1
 \end{array}$$

$$= \alpha_1$$

So $\check{\alpha}_g(\lambda_{\Gamma_1}, \dots, \lambda_{\Gamma_k}, \hat{\gamma}_1, \lambda_{S_1}, \dots, \lambda_{S_e}) = \alpha$.

Proof of (*):

It suffices to show that the following two composite 2-cells in \mathcal{X} coincide:

$$\begin{array}{c}
 X_0 \xrightarrow{\Gamma_1} X_1 \dots X_{k-1} \xrightarrow{\Gamma_k} X_k = A \xrightarrow{B(1, f)} B \xrightarrow{e_B} B \xrightarrow{g} D \xrightarrow{e_D} D \xrightarrow{D(g, \lambda)} C = Y_0 \xrightarrow{s_1} Y_1 \dots Y_{e-1} \xrightarrow{s_p} Y_e \\
 \parallel = \parallel = \parallel = \parallel = \parallel = \parallel = \parallel \\
 X_0 \xrightarrow{\Gamma_1} X_1 \dots X_{k-1} \xrightarrow{\Gamma_k} X_k = A \xrightarrow{B(1, f)} B \xrightarrow{g} D \xrightarrow{D(g, \lambda)} C = Y_0 \xrightarrow{s_1} Y_1 \dots Y_{e-1} \xrightarrow{s_p} Y_e \\
 \downarrow h \qquad \qquad \qquad \downarrow \alpha \qquad \qquad \qquad \downarrow k \\
 Z_0 \xrightarrow{t} Z_1
 \end{array}$$

and

$$\begin{array}{c}
 X_0 \xrightarrow{\Gamma_1} X_1 \dots X_{k-1} \xrightarrow{\Gamma_k} X_k = A \xrightarrow{B(1, f)} B \xrightarrow{e_B} B \xrightarrow{g} D \xrightarrow{e_D} D \xrightarrow{D(g, \lambda)} C = Y_0 \xrightarrow{s_1} Y_1 \dots Y_{e-1} \xrightarrow{s_p} Y_e \\
 \parallel = \parallel = \parallel = \parallel = \parallel = \parallel = \parallel \\
 X_0 \xrightarrow{\Gamma_1} X_1 \dots X_{k-1} \xrightarrow{\Gamma_k} X_k = A \xrightarrow{B(1, f)} B \xrightarrow{g} D \xrightarrow{D(g, \lambda)} C = Y_0 \xrightarrow{s_1} Y_1 \dots Y_{e-1} \xrightarrow{s_p} Y_e \\
 \downarrow h \qquad \qquad \qquad \downarrow \alpha \qquad \qquad \qquad \downarrow k \\
 Z_0 \xrightarrow{t} Z_1
 \end{array}$$

(Remark:

We will not directly make use of the 2-cell α . Yet, we include it to ensure that we have well-defined 2-cells in \mathcal{X} as the loose target of a 2-cell has to be unary.)

By the universal properties of the opcartesian 2-cells $\parallel \begin{array}{c} B \\ \xrightarrow{e_B} \\ B \end{array} \parallel$ and $\parallel \begin{array}{c} D \\ \xrightarrow{e_D} \\ D \end{array} \parallel$,

it is sufficient to prove equality of the two composite 2-cells below:

$$\begin{array}{c}
 X_0 \xrightarrow{\Gamma_1} X_1 \dots X_{k-1} \xrightarrow{\Gamma_k} X_k = A \xrightarrow{B(1, f)} B \xrightarrow{g} D \xrightarrow{D(g, \lambda)} C = Y_0 \xrightarrow{s_1} Y_1 \dots Y_{e-1} \xrightarrow{s_p} Y_e \\
 \parallel = \parallel = \parallel = \parallel = \parallel = \parallel = \parallel \\
 A \xrightarrow{B(1, f)} B \xrightarrow{g} D \xrightarrow{e_D} D \xrightarrow{D(g, \lambda)} C \\
 \parallel = \parallel \begin{array}{c} B \\ \xrightarrow{e_B} \\ B \end{array} \parallel = \parallel \begin{array}{c} D \\ \xrightarrow{e_D} \\ D \end{array} \parallel = \parallel \\
 A \xrightarrow{B(1, f)} B \xrightarrow{e_B} B \xrightarrow{g} D \xrightarrow{e_D} D \xrightarrow{D(g, \lambda)} C \\
 \parallel = \parallel = \parallel = \parallel = \parallel = \parallel = \parallel \\
 X_0 \xrightarrow{\Gamma_1} X_1 \dots X_{k-1} \xrightarrow{\Gamma_k} X_k = A \xrightarrow{B(1, f)} B \xrightarrow{g} D \xrightarrow{D(g, \lambda)} C = Y_0 \xrightarrow{s_1} Y_1 \dots Y_{e-1} \xrightarrow{s_p} Y_e \\
 \downarrow h \qquad \qquad \qquad \downarrow \alpha \qquad \qquad \qquad \downarrow k \\
 Z_0 \xrightarrow{t} Z_1
 \end{array}$$

and

$$\begin{array}{c}
 X_0 \xrightarrow{\Gamma_1} X_1 \dots X_{k-1} \xrightarrow{\Gamma_k} X_k = A \xrightarrow{B(\lambda, \rho)} B \xrightarrow{q} D \xrightarrow{D(g, \lambda)} C = Y_0 \xrightarrow{s_1} Y_1 \dots Y_{e-1} \xrightarrow{s_e} Y_e \\
 \parallel \qquad \qquad \qquad \parallel \qquad \qquad \qquad \parallel \qquad \qquad \qquad \parallel \qquad \qquad \qquad \parallel \\
 A \xrightarrow{B(\lambda, \rho)} B \xrightarrow{q} D \xrightarrow{D(g, \lambda)} C \\
 \parallel \qquad \qquad \qquad \parallel \qquad \qquad \qquad \parallel \qquad \qquad \qquad \parallel \\
 A \xrightarrow{B(\lambda, \rho)} B \xrightarrow{e_B} B = D \xrightarrow{D(g, \lambda)} C \\
 \parallel \qquad \qquad \qquad \parallel \qquad \qquad \qquad \parallel \\
 X_0 \xrightarrow{\Gamma_1} X_1 \dots X_{k-1} \xrightarrow{\Gamma_k} X_k = A \xrightarrow{B(\lambda, \rho)} B \xrightarrow{q} D \xrightarrow{D(g, \lambda)} C = Y_0 \xrightarrow{s_1} Y_1 \dots Y_{e-1} \xrightarrow{s_e} Y_e \\
 \downarrow h \qquad \qquad \qquad \alpha \qquad \qquad \qquad \downarrow k \\
 Z_0 \xrightarrow{t} Z_1
 \end{array}$$

$$\begin{array}{c}
 X_0 \xrightarrow{\Gamma_1} X_1 \dots X_{k-1} \xrightarrow{\Gamma_k} X_k = A \xrightarrow{B(\lambda, \rho)} B \xrightarrow{q} D \xrightarrow{D(g, \lambda)} C = Y_0 \xrightarrow{s_1} Y_1 \dots Y_{e-1} \xrightarrow{s_e} Y_e \\
 \parallel \qquad \qquad \qquad \parallel \qquad \qquad \qquad \parallel \qquad \qquad \qquad \parallel \\
 A \xrightarrow{B(\lambda, \rho)} B \xrightarrow{q} D \xrightarrow{D(g, \lambda)} C \\
 \parallel \qquad \qquad \qquad \parallel \qquad \qquad \qquad \parallel \\
 X_0 \xrightarrow{\Gamma_1} X_1 \dots X_{k-1} \xrightarrow{\Gamma_k} X_k = A \xrightarrow{B(\lambda, \rho)} B \xrightarrow{q} D \xrightarrow{D(g, \lambda)} C = Y_0 \xrightarrow{s_1} Y_1 \dots Y_{e-1} \xrightarrow{s_e} Y_e \\
 \downarrow h \qquad \qquad \qquad \alpha \qquad \qquad \qquad \downarrow k \\
 Z_0 \xrightarrow{t} Z_1
 \end{array}$$

constructions of

$\rho_{D(g, \lambda)}$ and $\lambda_{B(\lambda, \rho)}$ $= \alpha$.

Thus, we obtain $\xi = \alpha = \zeta$, which proves (*).

Uniqueness of $\check{\alpha}$:

Let β be any 2-cell in X with the same frame as $\check{\alpha}$ satisfying

$$\alpha = \beta \circ (\lambda_{\Gamma_1}, \dots, \lambda_{\Gamma_k}, \hat{\gamma}, \lambda_{s_1}, \dots, \lambda_{s_e}).$$

We have:

$$\begin{array}{c}
 X_0 \xrightarrow{\Gamma_1} X_1 \dots X_{k-1} \xrightarrow{\Gamma_k} X_k = A \xrightarrow{q(g, f)} C = C = Y_0 \xrightarrow{s_1} Y_1 \dots Y_{e-1} \xrightarrow{s_e} Y_e \\
 \parallel \qquad \qquad \qquad \parallel \qquad \qquad \qquad \parallel \qquad \qquad \qquad \parallel \\
 A \xrightarrow{e_A} A \xrightarrow{\bar{q}(g, f)} C \xrightarrow{e_C} C \\
 \parallel \qquad \qquad \qquad \parallel \qquad \qquad \qquad \parallel \\
 X_0 \xrightarrow{\Gamma_1} X_1 \dots X_{k-1} \xrightarrow{\Gamma_k} X_k = A \xrightarrow{B(\lambda, \rho)} B \xrightarrow{q} D \xrightarrow{D(g, \lambda)} C = Y_0 \xrightarrow{s_1} Y_1 \dots Y_{e-1} \xrightarrow{s_e} Y_e \\
 \downarrow h \qquad \qquad \qquad \alpha \qquad \qquad \qquad \downarrow k \\
 Z_0 \xrightarrow{t} Z_1
 \end{array}$$

$$\begin{array}{c}
 X_0 \xrightarrow{\tau_1} X_1 \dots X_{k-1} \xrightarrow{\tau_k} X_k = A \xrightarrow{q(g,f)} C = C = Y_0 \xrightarrow{s_1} Y_1 \dots Y_{e-1} \xrightarrow{s_e} Y_e \\
 \parallel \quad \quad \quad \parallel \quad \quad \quad \parallel \quad \quad \quad \parallel \\
 A \xrightarrow{e_A} A \xrightarrow{\bar{q}(g,f)} C \xrightarrow{e_C} C \\
 \parallel \quad \quad \quad \downarrow f \quad \quad \quad \downarrow g \quad \quad \quad \parallel \\
 A \xrightarrow{B(\tau, f)} B \xrightarrow{q} D \xrightarrow{D(g, \tau)} C \\
 \parallel \quad \quad \quad \hat{y} \quad \quad \quad \parallel \\
 X_0 \xrightarrow{\tau_1} X_1 \dots X_{k-1} \xrightarrow{\tau_k} X_k = A \xrightarrow{q(g,f)} C = Y_0 \xrightarrow{s_1} Y_1 \dots Y_{e-1} \xrightarrow{s_e} Y_e \\
 \downarrow h \quad \quad \quad \beta \quad \quad \quad \downarrow k \\
 Z_0 \xrightarrow{t} Z_1
 \end{array}$$

$\alpha = \beta \circ (\tau_1, \dots, \tau_k, \hat{y}, s_1, \dots, s_e)$,
 by assumption on β

$$\begin{array}{c}
 X_0 \xrightarrow{\tau_1} X_1 \dots X_{k-1} \xrightarrow{\tau_k} X_k = A \xrightarrow{q(g,f)} C = C = Y_0 \xrightarrow{s_1} Y_1 \dots Y_{e-1} \xrightarrow{s_e} Y_e \\
 \parallel \quad \quad \quad \parallel \quad \quad \quad \parallel \quad \quad \quad \parallel \\
 X_0 \xrightarrow{\tau_1} X_1 \dots X_{k-1} \xrightarrow{\tau_k} X_k = A \xrightarrow{q(g,f)} C = Y_0 \xrightarrow{s_1} Y_1 \dots Y_{e-1} \xrightarrow{s_e} Y_e \\
 \downarrow h \quad \quad \quad \beta \quad \quad \quad \downarrow k \\
 Z_0 \xrightarrow{t} Z_1
 \end{array}$$

$= \beta$

so $\check{\alpha} = \beta$.

Proof of (*):

It suffices to show that the composite 2-cell

$$\begin{array}{c}
 A \xrightarrow{q(g,f)} C = C \\
 \parallel \quad \quad \quad \parallel \\
 A \xrightarrow{e_A} A \xrightarrow{\bar{q}(g,f)} C \xrightarrow{e_C} C \\
 \parallel \quad \quad \quad \downarrow f \quad \quad \quad \downarrow g \quad \quad \quad \parallel \\
 A \xrightarrow{B(\tau, f)} B \xrightarrow{q} D \xrightarrow{D(g, \tau)} C \\
 \parallel \quad \quad \quad \hat{y} \quad \quad \quad \parallel \\
 A \xrightarrow{q(g,f)} C
 \end{array}$$

is equal to the identity 2-cell $\parallel \tau_{q(g,f)} \parallel$ of $q(g,f): A \rightarrow C$.

By the universal property of the cartesian 2-cell $f \downarrow \tau_{q(g,f)} \text{ cart} \downarrow g$ of $B \xrightarrow{q} D$.

it is sufficient to show that the post-compositions with $\bar{q}(g, f)$ agree, i.e., that

$$\begin{array}{ccc}
 A & \xrightarrow{q(g, f)} & C \\
 \parallel e_A \parallel & & \parallel e_C \parallel \\
 A & \xrightarrow{e_A} & A \xrightarrow{\bar{q}(g, f)} & C \xrightarrow{e_C} & C \\
 \parallel v(1, f) \parallel & \downarrow f & & \downarrow v(g, 1) & \parallel \\
 A & \xrightarrow{B(1, f)} & B \xrightarrow{q} & D \xrightarrow{D(g, 1)} & C \\
 \parallel & & \hat{\gamma} & & \parallel \\
 A & \xrightarrow{q(g, f)} & C \\
 \downarrow f & & \downarrow g \\
 B & \xrightarrow{\bar{q}(g, f)} & D
 \end{array}
 =
 \begin{array}{ccc}
 A & \xrightarrow{q(g, f)} & C \\
 \downarrow f & & \downarrow g \\
 B & \xrightarrow{q} & D
 \end{array}$$

We compute:

$$\begin{array}{ccc}
 A & \xrightarrow{q(g, f)} & C \\
 \parallel e_A \parallel & & \parallel e_C \parallel \\
 A & \xrightarrow{e_A} & A \xrightarrow{\bar{q}(g, f)} & C \xrightarrow{e_C} & C \\
 \parallel v(1, f) \parallel & \downarrow f & & \downarrow v(g, 1) & \parallel \\
 A & \xrightarrow{B(1, f)} & B \xrightarrow{q} & D \xrightarrow{D(g, 1)} & C \\
 \parallel & & \hat{\gamma} & & \parallel \\
 A & \xrightarrow{q(g, f)} & C \\
 \downarrow f & & \downarrow g \\
 B & \xrightarrow{\bar{q}(g, f)} & D
 \end{array}
 \begin{array}{c}
 = \\
 \uparrow \\
 \text{construction} \\
 \text{of } \hat{\gamma}
 \end{array}$$

$$\begin{array}{ccc}
 A & \xrightarrow{q(g, f)} & C \\
 \parallel e_A \parallel & & \parallel e_C \parallel \\
 A & \xrightarrow{e_A} & A \xrightarrow{\bar{q}(g, f)} & C \xrightarrow{e_C} & C \\
 \parallel v(1, f) \parallel & \downarrow f & & \downarrow v(g, 1) & \parallel \\
 A & \xrightarrow{B(1, f)} & B \xrightarrow{q} & D \xrightarrow{D(g, 1)} & C \\
 \downarrow f & \parallel e_B \parallel & \parallel = \parallel & \parallel \lambda_q \parallel & \parallel e_D \parallel \\
 B & \xrightarrow{e_B} & B \xrightarrow{q} & D \xrightarrow{e_D} & D \\
 \parallel & & \rho_q & & \parallel \\
 B & \xrightarrow{q} & D
 \end{array}$$

$$\begin{array}{ccc}
 A & \xrightarrow{q(g, f)} & C \\
 \parallel e_A \parallel & & \parallel e_C \parallel \\
 A & \xrightarrow{e_A} & A \xrightarrow{\bar{q}(g, f)} & C \xrightarrow{e_C} & C \\
 \downarrow f & \parallel e_f \parallel & \downarrow f & \parallel e_g \parallel & \downarrow g \\
 B & \xrightarrow{e_B} & B \xrightarrow{q} & D \xrightarrow{e_D} & D \\
 \parallel & & \lambda_q & & \parallel \\
 B & \xrightarrow{e_B} & B \xrightarrow{q} & D \\
 \parallel & & \rho_q & & \parallel \\
 B & \xrightarrow{q} & D
 \end{array}$$

$$\begin{array}{ccc}
 A & \xrightarrow{q(g, f)} & C \\
 \downarrow f & \parallel e_B \parallel & \parallel = \parallel & \parallel e_D \parallel \\
 B & \xrightarrow{e_B} & B \xrightarrow{q} & D \xrightarrow{e_D} & D \\
 \parallel & & \lambda_q & & \parallel \\
 B & \xrightarrow{e_B} & B \xrightarrow{q} & D \\
 \parallel & & \rho_q & & \parallel \\
 B & \xrightarrow{q} & D
 \end{array}$$

constructions of $v(1, f)$ and $v(g, 1)$

constructions of e_f and e_g

$$\begin{array}{ccc}
 A & \xrightarrow{q(g, f)} & C \\
 \downarrow f & \parallel e_B \parallel & \parallel = \parallel \\
 B & \xrightarrow{e_B} & B \xrightarrow{q} & D \\
 \parallel & & \rho_q & & \parallel \\
 B & \xrightarrow{q} & D
 \end{array}
 \begin{array}{c}
 = \\
 \uparrow \\
 \text{construction} \\
 \text{of } \lambda_q
 \end{array}$$

$$\begin{array}{ccc}
 A & \xrightarrow{q(g, f)} & C \\
 \downarrow f & \parallel e_B \parallel & \parallel = \parallel \\
 B & \xrightarrow{e_B} & B \xrightarrow{q} & D \\
 \parallel & & \rho_q & & \parallel \\
 B & \xrightarrow{q} & D
 \end{array}
 \begin{array}{c}
 = \\
 \uparrow \\
 \text{construction} \\
 \text{of } \rho_q
 \end{array}$$

This proves (*).

We may conclude that $\check{\alpha}$ is the unique 2-cell in \mathbb{X} of its form which satisfies $\alpha = \check{\alpha} \circ (1_{r_1}, \dots, 1_{r_n}, \hat{\gamma}, 1_{s_1}, \dots, 1_{s_e})$.

This shows that the 2-cell $\hat{\gamma}$ is opcartesian,

$$\begin{array}{ccccc} A & \xrightarrow{B(1,f)} & B & \xrightarrow{g} & D & \xrightarrow{D(g,1)} & C \\ & & & \hat{\gamma} & & & \\ \parallel & & & \text{opcart} & & & \parallel \\ A & \xrightarrow{g(g,f)} & & & & & C. \end{array}$$

Hence, it exhibits the loose arrow $g(g,f): A \rightrightarrows C$ as a loose composite of the triple $(B(1,f), g, D(g,1))$.

Corollary.

Let \mathbb{X} be a virtual equipment.

Let $f: A \rightarrow B$ and $g: B \rightarrow C$ be tight arrows in \mathbb{X} .

1) Consider the distinguished loose companion $C(1, gf): A \rightrightarrows C$ of the composite tight arrow $gf: A \rightarrow C$ in \mathbb{X} .

$$\begin{array}{ccc} A & \xrightarrow{C(1,gf)} & C \\ gf \downarrow & \bar{C}(1,gf) \text{ cart} & \parallel \\ C & \xrightarrow{e_c} & C \end{array}$$

The loose arrow $C(1, gf): A \rightrightarrows C$ is also a loose composite of the pair $(B(1, f), C(1, g))$ of the distinguished loose companions $B(1, f): A \rightrightarrows B$ of f and $C(1, g): B \rightrightarrows C$ of g , i.e., \exists opcartesian 2-cell in \mathbb{X} of shape

$$\begin{array}{ccccc} A & \xrightarrow{B(1,f)} & B & \xrightarrow{C(1,g)} & C \\ & & & \text{opcart} & \\ \parallel & & & & \parallel \\ A & \xrightarrow{C(1,gf)} & & & C. \end{array}$$

In particular, the pair $(B(1, f), C(1, g))$ admits a loose composite in \mathbb{X} .

Given any loose composite $B(1, f) \circ C(1, g)$ of $(B(1, f), C(1, g))$, then

$$C(1, gf) \cong B(1, f) \circ C(1, g) \quad \text{in } \mathbb{X}[A, C].$$

2) Similarly, consider the distinguished loose conjoint $C(g, f, \lambda): C \rightrightarrows A$ of the composite tight arrow $gf: A \rightarrow C$ in \mathcal{X} .

$$\begin{array}{ccc} C & \xrightarrow{C(g, f, \lambda)} & A \\ \parallel \bar{C}(g, f, \lambda) \downarrow gf & & \\ C & \xrightarrow[e_c]{} & C \end{array}$$

The loose arrow $C(g, f, \lambda): C \rightrightarrows A$ is also a loose composite of the pair $(C(g, \lambda), B(f, \lambda))$ of the distinguished loose conjoints $C(g, \lambda): C \rightrightarrows B$ of g and $B(f, \lambda): B \rightrightarrows A$ of f , i.e., \exists opcartesian 2-cell in \mathcal{X} of shape

$$\begin{array}{ccccc} C & \xrightarrow{C(g, \lambda)} & B & \xrightarrow{B(f, \lambda)} & A \\ \parallel & & \text{opcart} & & \parallel \\ C & \xrightarrow{C(g, f, \lambda)} & & & A. \end{array}$$

In particular, the pair $(C(g, \lambda), B(f, \lambda))$ admits a loose composite in \mathcal{X} .

Given any loose composite $C(g, \lambda) \circ B(f, \lambda)$ of $(C(g, \lambda), B(f, \lambda))$, then the loose arrows $C(g, f, \lambda), C(g, \lambda) \circ B(f, \lambda): C \rightrightarrows A$ in \mathcal{X} are isomorphic as objects in $\mathcal{X} \llbracket C, A \rrbracket$,

$$C(g, f, \lambda) \cong C(g, \lambda) \circ B(f, \lambda).$$

Corollary.

Let \mathcal{X} be a virtual equipment.

Let $f: A \rightarrow B$ and $g: C \rightarrow B$ be tight arrows in \mathcal{X} .

We consider the niche configuration $\begin{array}{ccc} A & & C \\ f \downarrow & & \downarrow g \\ B & \xrightarrow[e_B]{} & B \end{array}$ and the corresponding distinguished

cartesian 2-cell $\begin{array}{ccc} A & \xrightarrow{B(g, f)} & C \\ f \downarrow \bar{B}(g, f) \text{ cart} \downarrow g & & \\ B & \xrightarrow[e_B]{} & B. \end{array}$

The loose arrow $B(g, f): A \rightrightarrows C$ is also a loose composite of the pair $(B(\lambda, f), B(g, \lambda))$ of the companion $B(\lambda, f): A \rightrightarrows B$ of f and the conjoint $B(g, \lambda): B \rightrightarrows C$ of g , i.e., \exists opcartesian 2-cell in \mathcal{X} of shape

$$\begin{array}{ccccc} A & \xrightarrow{B(\lambda, f)} & B & \xrightarrow{B(g, \lambda)} & C \\ \parallel & & \text{opcart} & & \parallel \\ A & \xrightarrow{B(g, f)} & & & C. \end{array}$$

In particular, the pair $(B(\lambda, f), B(g, \lambda))$ admits a loose composite in \mathcal{X} .

\forall loose composite $B(\lambda, f) \circ B(g, \lambda)$ of $(B(\lambda, f), B(g, \lambda))$, we have

$$B(g, f) \cong B(\lambda, f) \circ B(g, \lambda) \text{ in } \mathcal{X} \llbracket A, C \rrbracket.$$

Theorem.

Let X be a virtual equipment.

Suppose we are given tight and loose arrows in X as depicted in the following configuration:

$$\begin{array}{ccc} A & \xrightarrow{P} & C \\ f \downarrow & & \downarrow g \\ B & \xrightarrow{q} & D. \end{array}$$

Then there exists a bijection between 2-cells in X with frame

$$\begin{array}{ccc} \begin{array}{ccc} A & \xrightarrow{P} & C \\ f \downarrow & & \downarrow g \\ B & \xrightarrow{q} & D \end{array} & \text{and} & \begin{array}{ccccc} B & \xrightarrow{B(f,1)} & A & \xrightarrow{P} & C & \xrightarrow{D(1,g)} & D \\ \parallel & & & & & & \parallel \\ B & \xrightarrow{\quad\quad\quad} & & \xrightarrow{\quad\quad\quad} & & \xrightarrow{\quad\quad\quad} & D. \end{array} \end{array}$$

Proof.

We consider the following 2-cells in X :

$$\begin{array}{ccc} \begin{array}{ccc} B & \xrightarrow{B(f,1)} & A \\ \parallel \bar{B}(f,1) \downarrow f & & \\ B & \xrightarrow{e_B} & B, \end{array} & & \begin{array}{ccc} A & \xrightarrow{e_A} & A \\ \downarrow v(f,1) \parallel & & \\ B & \xrightarrow{B(f,1)} & A, \end{array} \\ \\ \begin{array}{ccc} C & \xrightarrow{D(1,g)} & D \\ g \downarrow \bar{D}(1,g) \parallel & & \\ D & \xrightarrow{e_D} & D, \end{array} & & \begin{array}{ccc} C & \xrightarrow{e_C} & C \\ \parallel v(1,g) \downarrow g & & \\ C & \xrightarrow{D(1,g)} & D. \end{array} \end{array}$$

1) Suppose we are given a 2-cell α in X of the form $\begin{array}{ccc} A & \xrightarrow{P} & C \\ f \downarrow & \alpha & \downarrow g \\ B & \xrightarrow{q} & D. \end{array}$

We form the following composite 2-cell in X :

$$S(\alpha) := \begin{array}{ccccc} B & \xrightarrow{B(f,1)} & A & \xrightarrow{P} & C & \xrightarrow{D(1,g)} & D \\ \parallel \bar{B}(f,1) \downarrow f & & \alpha & & g \downarrow \bar{D}(1,g) \parallel & & \\ B & \xrightarrow{e_B} & B & \xrightarrow{q} & D & \xrightarrow{e_D} & D \\ \parallel & & = & & \parallel \lambda_q & & \parallel \\ B & \xrightarrow{e_B} & B & \xrightarrow{\quad\quad\quad} & & \xrightarrow{\quad\quad\quad} & D \\ \parallel & & \rho_q & & & & \parallel \\ B & \xrightarrow{\quad\quad\quad} & & \xrightarrow{\quad\quad\quad} & & \xrightarrow{\quad\quad\quad} & D. \end{array}$$

This provides a well-defined map

$$S: \left\{ \begin{array}{c} \text{2-cells in } X \text{ with frame} \\ \begin{array}{ccc} A & \xrightarrow{P} & C \\ f \downarrow & & \downarrow g \\ B & \xrightarrow{q} & D \end{array} \end{array} \right\} \rightarrow \left\{ \begin{array}{c} \text{2-cells in } X \text{ with frame} \\ \begin{array}{ccccc} B & \xrightarrow{B(f,1)} & A & \xrightarrow{P} & C & \xrightarrow{D(1,g)} & D \\ \parallel & & & & & & \parallel \\ B & \xrightarrow{q} & & & & & D \end{array} \end{array} \right\}$$

$$\alpha \mapsto S(\alpha).$$

2) Suppose we are given a 2-cell in X of the form

$$\begin{array}{ccccc} B & \xrightarrow{B(f,1)} & A & \xrightarrow{P} & C & \xrightarrow{D(1,g)} & D \\ \parallel & & & & & & \parallel \\ B & \xrightarrow{q} & & & & & D \end{array}$$

We form the following composite 2-cell in X :

$$T(\beta) := \begin{array}{c} \begin{array}{c} \begin{array}{ccc} A & \xrightarrow{P} & C \\ \varepsilon_A \parallel = \parallel \varepsilon_C \\ A & \xrightarrow{e_A} & A & \xrightarrow{P} & C & \xrightarrow{e_C} & C \\ f \downarrow & \nu(f,1) \parallel = \parallel \nu(1,g) \downarrow g \\ B & \xrightarrow{B(f,1)} & A & \xrightarrow{P} & C & \xrightarrow{D(1,g)} & D \\ \parallel & & \beta & & \parallel \\ B & \xrightarrow{q} & & & & & D \end{array} \end{array} \end{array}$$

This yields a well-defined map

$$T: \left\{ \begin{array}{c} \text{2-cells in } X \text{ with frame} \\ \begin{array}{ccccc} B & \xrightarrow{B(f,1)} & A & \xrightarrow{P} & C & \xrightarrow{D(1,g)} & D \\ \parallel & & & & & & \parallel \\ B & \xrightarrow{q} & & & & & D \end{array} \end{array} \right\} \rightarrow \left\{ \begin{array}{c} \text{2-cells in } X \text{ with frame} \\ \begin{array}{ccc} A & \xrightarrow{P} & C \\ f \downarrow & & \downarrow g \\ B & \xrightarrow{q} & D \end{array} \end{array} \right\}$$

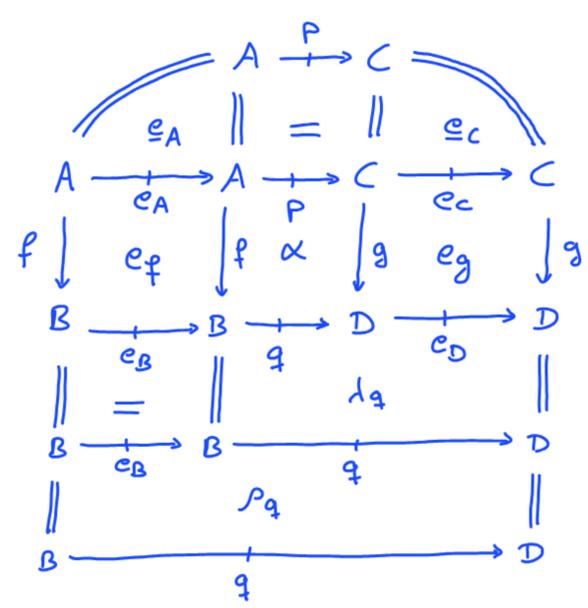
$$\beta \mapsto T(\beta).$$

3) We want to show that the maps S and T are inverse to each other.

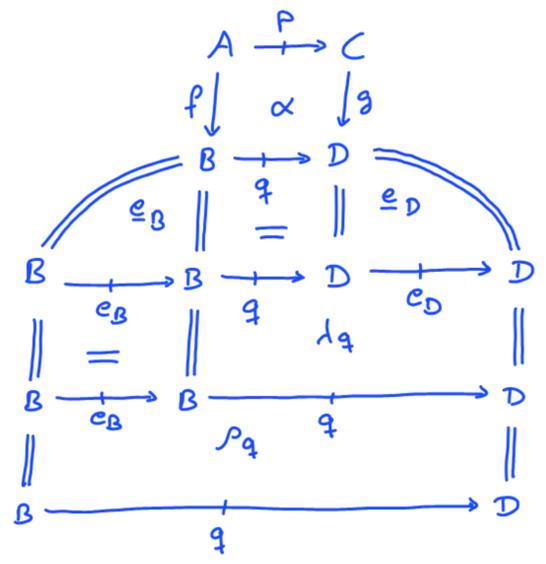
3.1) \forall 2-cell α as in 1), we have $T(S(\alpha)) = \alpha$:

$$T(S(\alpha)) \stackrel{\text{definition of } T}{=} \begin{array}{c} \begin{array}{c} \begin{array}{ccc} A & \xrightarrow{P} & C \\ \varepsilon_A \parallel = \parallel \varepsilon_C \\ A & \xrightarrow{e_A} & A & \xrightarrow{P} & C & \xrightarrow{e_C} & C \\ f \downarrow & \nu(f,1) \parallel = \parallel \nu(1,g) \downarrow g \\ B & \xrightarrow{B(f,1)} & A & \xrightarrow{P} & C & \xrightarrow{D(1,g)} & D \\ \parallel & & S(\alpha) & & \parallel \\ B & \xrightarrow{q} & & & & & D \end{array} \end{array} \end{array} \stackrel{\text{definition of } S}{=} \begin{array}{c} \begin{array}{c} \begin{array}{ccc} A & \xrightarrow{P} & C \\ \varepsilon_A \parallel = \parallel \varepsilon_C \\ A & \xrightarrow{e_A} & A & \xrightarrow{P} & C & \xrightarrow{e_C} & C \\ f \downarrow & \nu(f,1) \parallel = \parallel \nu(1,g) \downarrow g \\ B & \xrightarrow{B(f,1)} & A & \xrightarrow{P} & C & \xrightarrow{D(1,g)} & D \\ \parallel & & \alpha & & \parallel \\ B & \xrightarrow{e_B} & B & \xrightarrow{q} & D & \xrightarrow{e_D} & D \\ \parallel & & \parallel & & \parallel \\ B & \xrightarrow{e_B} & B & \xrightarrow{q} & D & \xrightarrow{e_D} & D \\ \parallel & & \rho_q & & \parallel \\ B & \xrightarrow{e_B} & B & \xrightarrow{q} & D & \xrightarrow{e_D} & D \\ \parallel & & \rho_q & & \parallel \\ B & \xrightarrow{q} & & & & & D \end{array} \end{array} \end{array}$$

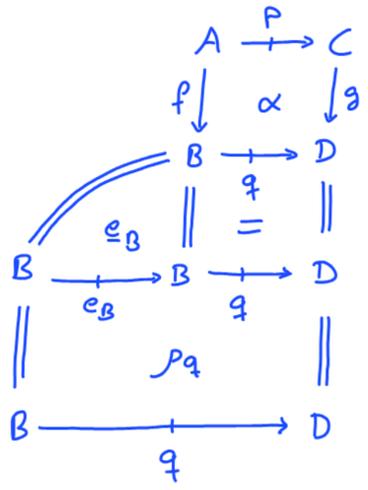
↑
= constructions of $v(f, \lambda)$ and $v(\lambda, g)$



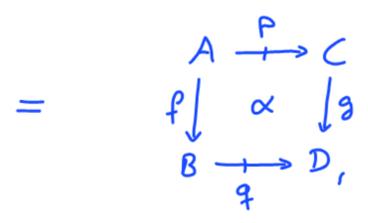
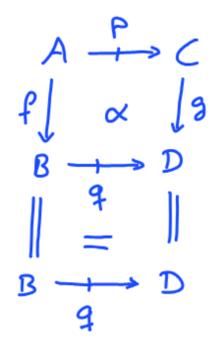
↑
= constructions of e_f and e_g



↑
= construction of λ_q



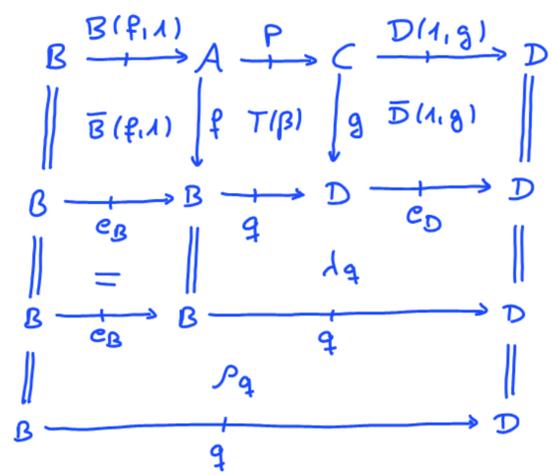
↑
= construction of ρ_q



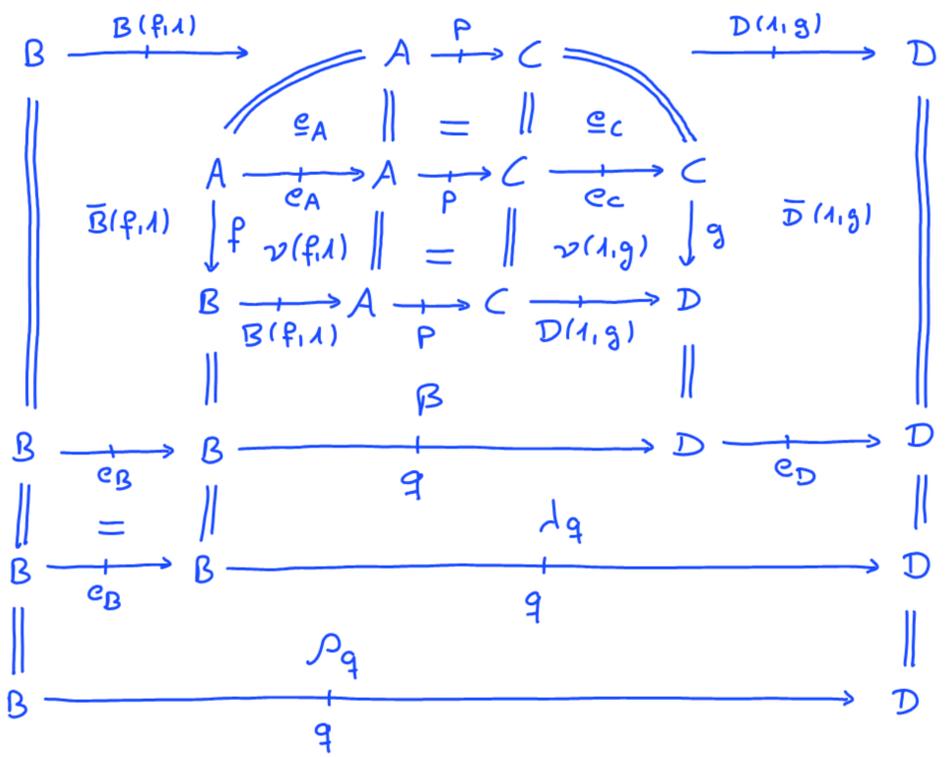
so $T(S(\alpha)) = \alpha$.

3.2) \forall 2-cell β as in 2), we have $S(T(\beta)) = \beta$:

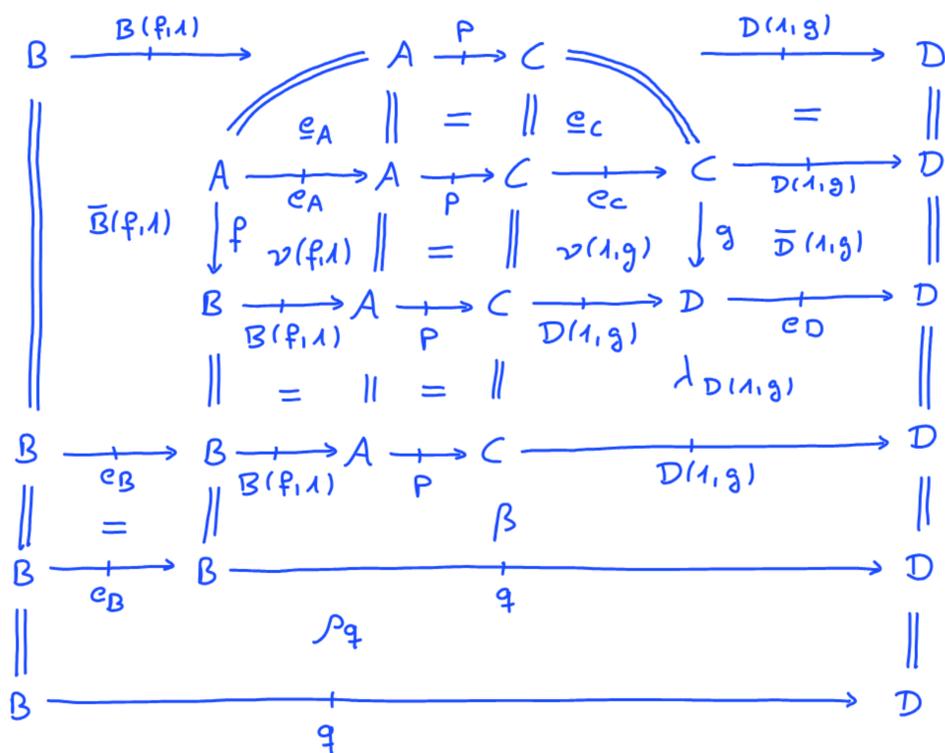
↑
= definition of S



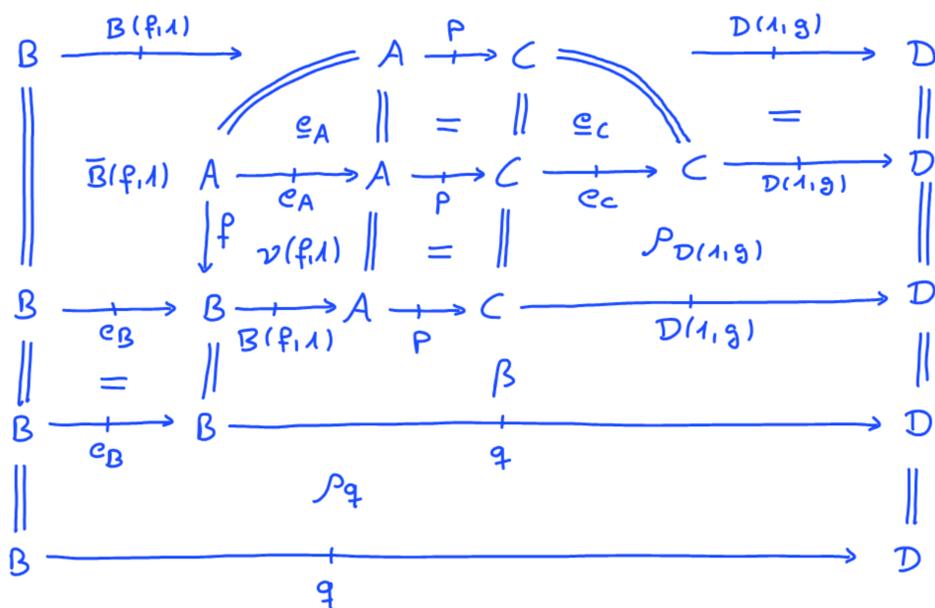
↑
= definition of T



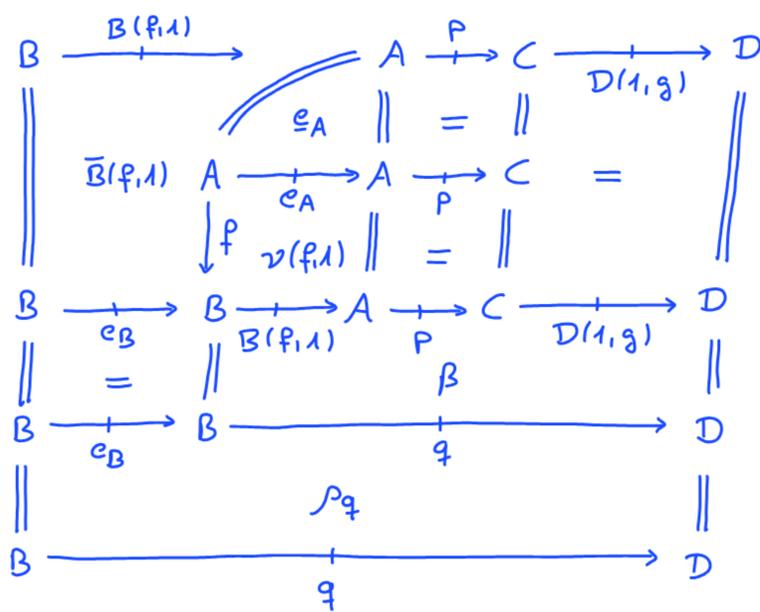
\uparrow
 =
 naturality
 of λ wrt β
 (and $\lambda e_D = e_{1_D}$)



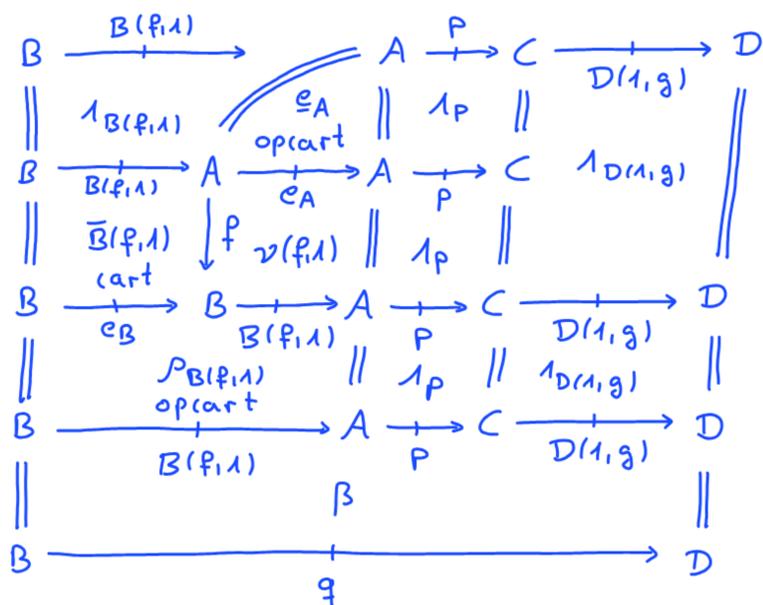
\uparrow
 =
 by an earlier remark
 concerning $\nu(1,g)$



\uparrow
 =
 construction
 of $\rho_{D(1,g)}$

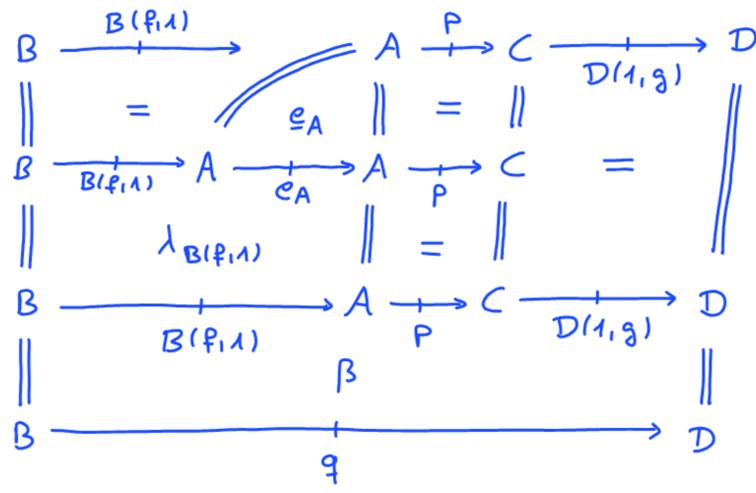


\uparrow
 =
 naturality
 of ρ wrt β
 (and $\lambda e_B = e_{1_B}$)



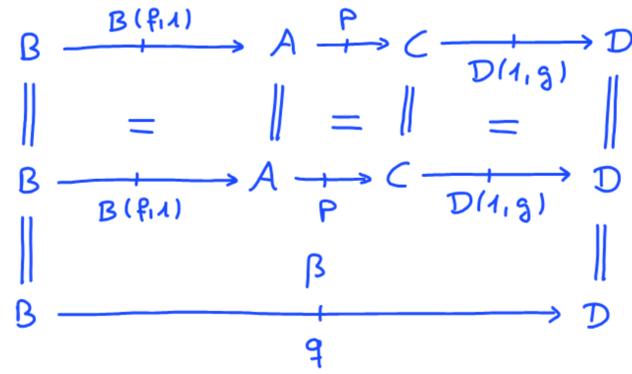
by an earlier remark concerning $\gamma(f, 1)$

\Rightarrow

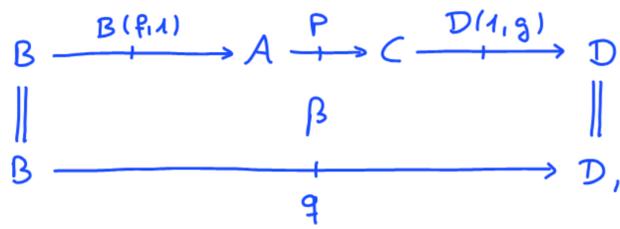


construction of $\lambda_{B(f, 1)}$

\Rightarrow



=



so $S(T(\beta)) = \beta$.

By 3.1) and 3.2), the maps S and T are inverse to each other and thus mutually inverse bijections.

====