

TOWARDS STRUCTURES OF HIGHER CATEGORICAL STRUCTURES

KEISUKE HOSHINO

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1. INTRODUCTION

Small categories forms a 2-category, but when profunctors are also taken into account, they form a double category $\mathbb{P}\text{rof}$. In [CS10], Cruttwell and Shulman defined a unified framework to conceive several notions of *generalised multicategories*, which are kinds of category-like structures, and showed that each of those concepts forms a *virtual double category*, which is a generalisation of double category. On the other hand, there are also some attempts to formalise category theory by means of *augmented virtual double category*, which is also an generalisation of double category [Kou22]. Thus, it is natural to assume that the collection of 1-dimensional structures such as categories forms a 2-dimensional structure. In general, the collection of n -dimensional structure is expected to have an $n + 1$ -dimensional structure. The main objective of this paper is to give an answer to this expectation.

This paper uses the theory on familial monads introduced by Shapiro in [Sha21, Sha22] as a general framework for defining higher categories. In general, for a cartesian monad T , a generalised notion of category, called T -category in this paper, is defined [Bur71, Lei04], whereas in [Sha22], it is asserted that there exists another familial monad $\text{fc}[T]$ whose algebras coincide with T -categories, whenever T is familial. When T is trivial, T -categories are ordinary categories and $\text{fc}[T]$ -categories are virtual double categories, hence $\text{fc}[T]$ -category is reworded as T -virtual double category. Moreover, for $\text{fc}^n[T]$ obtained by repeating $\text{fc}[-]$, $\text{fc}^n[T]$ -categories (=virtual $n + 1$ -tuple categories) can be considered an example of concepts of $(n + 1)$ -dimensional structure.

In Section 4.1, under some assumptions on T , we suggest a definition of the category of T -simplices, hence we obtain notions such as T -simplicial set and T -simplicial category, and show that the category of T -categories is embedded in the category of T -simplicial set.

In Section 4.2, we investigate the virtual double category of T -categories and T -profunctors, $\mathbb{P}\text{rof}(T)$, which is an example of the virtual double category of T -monoids defined in [CS10], in terms of *pseudo simplicial category*, i.e. pseudo functor from the category of simplices to 2-category of categories. We show that one can define a (2-truncated) pseudo simplicial category of T -categories and T -profunctors, $\overline{\mathbb{P}\text{rof}}_2(T)$, and $\mathbb{P}\text{rof}(T)$ is the free objects with respect to a “nerve” 2-functor from the 2-category of virtual double categories to the 2-category of pseudo simplicial categories, hence $\mathbb{P}\text{rof}(T)$ can be seen as the “realization” of $\overline{\mathbb{P}\text{rof}}_2(T)$ as a virtual double category.

Combining those observations, we suggest an definition of the fc -pseudo simplicial category of virtual double categories, and the virtual triple category of virtual double categories as its *realization*.

Moreover, we suggest a way to define the virtual $n + 2$ -tuple category of virtual $n + 1$ -tuple categories, in general.

2. TERMINOLOGY AND PRELIMINARIES

A *double category* \mathbb{X} is a category pseudo internal to the 2-category Cat of categories, hence it has the following data:

- a set of *objects* in \mathbb{X}

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- for each pair of objects X, Y , a set of *horizontal arrows*, written as slashed arrows like $X \dashrightarrow Y$
- for each pair of objects X, Y , a set of *vertical arrows*, written as $X \rightarrow Y$
- for each square formed by horizontal and vertical arrows, a set of *cells*, written as follows

$$(2.1) \quad \begin{array}{ccc} X & \xrightarrow{p} & Y \\ f \downarrow & \alpha & \downarrow g \\ A & \xrightarrow{q} & B \end{array}$$

- vertical and horizontal compositions and identities satisfying coherence conditions

We say a cell α above is *horizontal* if f and g are identities. Horizontal arrows and horizontal cells forms a bicategory, which we write $\mathcal{H}(\mathbb{X})$.

By the notations on the left side of the equations below, we mean the cells denoted on the right side;

$$(2.2) \quad \begin{array}{ccc} X & \xrightarrow{p} & Y \\ f \searrow & \gamma & \swarrow g \\ & A & \end{array} = \begin{array}{ccc} X & \xrightarrow{p} & Y \\ f \downarrow & \gamma & \downarrow g \\ & A & \xrightarrow{\text{Id}_A} & A \end{array}, \quad \begin{array}{ccc} X & & \\ f \swarrow & \gamma & \searrow g \\ & A & \xrightarrow{q} & B \end{array} = \begin{array}{ccc} X & \xrightarrow{\text{Id}_X} & X \\ f \downarrow & \gamma & \downarrow g \\ & A & \xrightarrow{q} & B \end{array}, \quad \begin{array}{ccc} X & & \\ \downarrow \gamma & & \\ & A & \xrightarrow{\text{Id}_A} & A \end{array} = \begin{array}{ccc} X & \xrightarrow{\text{Id}_X} & X \\ f \downarrow & \gamma & \downarrow g \\ & A & \xrightarrow{\text{Id}_A} & A \end{array}$$

For a vertical arrow $f : X \rightarrow A$ in a double category \mathbb{X} , a *companion* of f is a horizontal cell $f_* : X \dashrightarrow A$ such that there exists two cells, α and β , satisfying conditions below:

$$(2.3) \quad \begin{array}{ccc} & X & \\ \parallel & \searrow f & \\ X & \xrightarrow{f_*} & A \\ f \searrow & \beta & \parallel \\ & A & \end{array} = \begin{array}{ccc} X & & \\ f \downarrow & (=) & \downarrow f \\ & A & \end{array} \quad \begin{array}{ccc} X & \xrightarrow{f_*} & A \\ \parallel & \searrow f & \parallel \\ X & \xrightarrow{f_*} & A \\ \parallel & \alpha & \parallel \\ X & \xrightarrow{f_*} & A \end{array} = \begin{array}{ccc} X & \xrightarrow{f_*} & A \\ \parallel & \parallel & \parallel \\ X & \xrightarrow{f_*} & A \end{array}$$

where $=$ and \parallel are horizontal and vertical identity cells. A *conjoint* f^* of f is the horizontal dual of companion. For each vertical arrow $f : X \rightarrow Y$, its companion and conjoint are unique if exist. \mathbb{X} is called an *equipment* if every vertical arrows have companions and conjoins. A cell α (2.1) is *cartesian* if any cell on the right below uniquely factors through α :

$$(2.4) \quad \begin{array}{ccc} \cdot & \dashrightarrow & \cdot \\ k \downarrow & \bar{\beta} & \downarrow l \\ X & \xrightarrow{p} & Y \\ f \downarrow & \alpha & \downarrow g \\ A & \xrightarrow{q} & B \end{array} = \begin{array}{ccc} \cdot & \dashrightarrow & \cdot \\ f k \downarrow & \beta & \downarrow g l \\ A & \xrightarrow{q} & B \end{array}$$

If \mathbb{X} is an equipment, then α is cartesian if and only if the composite below is an isomorphism in $\mathcal{H}(\mathbb{X})$.

$$(2.5) \quad \begin{array}{ccccc} & X & \dashrightarrow & Y & \\ & \parallel & \searrow f & \alpha & \searrow g & \parallel \\ X & \xrightarrow{f_*} & A & \dashrightarrow & B & \xrightarrow{g_*} & Y \end{array}$$

Example 2.1. For each small categories \mathbf{X} and \mathbf{Y} , we define a *profunctor* $\mathbf{X} \dashrightarrow \mathbf{Y}$ as a functor $\mathbf{X}^{\text{op}} \times \mathbf{Y} \rightarrow \mathbf{Set}$. There is the double category $\mathbb{P}\text{rof}$ consisting of small categories as its objects, functors as its vertical arrows, and profunctors as its horizontal arrows. A cell in $\mathbb{P}\text{rof}$

$$\begin{array}{ccc} \mathbf{X} & \xrightarrow{p} & \mathbf{Y} \\ f \downarrow & \alpha & \downarrow g \\ \mathbf{A} & \xrightarrow{q} & \mathbf{B} \end{array}$$

is a family of functions $\alpha_{x,y} : p(x,y) \rightarrow q(f(x),g(y))$ which is natural in $x \in \mathbf{X}$ and $y \in \mathbf{Y}$. $\mathbb{P}\text{rof}$ is in fact an equipment, where for each functor $f : \mathbf{X} \rightarrow \mathbf{A}$, the companion $f_* : \mathbf{X} \dashrightarrow \mathbf{A}$ and conjoint $f^* : \mathbf{A} \dashrightarrow \mathbf{X}$ is defined as $\mathbf{A}(f(x), a)$ and $\mathbf{A}(x, f(a))$ respectively. The horizontal composition of $\mathbf{X} \xrightarrow{p} \mathbf{Y} \xrightarrow{q} \mathbf{Z}$ is given by the following coend formula:

$$(2.6) \quad q \circ p(x, z) := \int^{y \in \mathbf{Y}} p(x, y) \times q(y, z)$$

By $\mathcal{P}\text{rof}$, we mean the horizontal bicategory $\mathcal{H}(\mathbb{P}\text{rof})$. ■

Example 2.2. Let \mathbf{E} be a category with pullbacks. Then there is a double category of *spans* in \mathbf{E} , $\mathbb{S}\text{pan}(\mathbf{E})$, defined as follows: objects are those of \mathbf{E} , vertical arrows are morphisms of \mathbf{E} , and a horizontal arrow $p : X \dashrightarrow Y$ is a span $X \leftarrow |p| \rightarrow Y$ in \mathbf{E} . A cell α like (2.1) is a morphism $\alpha : |p| \rightarrow |q|$ which makes obvious diagram consisting of f, g , and the legs of p and q commutes. The horizontal composite $q \circ p$ is given by pullback of the right leg of p along the left leg of q . ■

A *virtual double category* \mathbb{A} has the following data:

- data of objects, horizontal arrows, and vertical arrows as in the definition of double category.
- for each square of the form below, a set of *cells*, written as follows

$$(2.7) \quad \begin{array}{ccccc} X_0 & \xrightarrow{p_0} & X_1 & \xrightarrow{p_1} & \cdots & \xrightarrow{\quad} & X_{n-1} & \xrightarrow{p_{n-1}} & X_n \\ f \downarrow & & & & \alpha & & & & \downarrow g \\ A_0 & \xrightarrow{\quad} & & \xrightarrow{\quad} & & \xrightarrow{\quad} & & \xrightarrow{\quad} & A_1 \end{array}$$

- vertical arrows has composition, so that objects and vertical arrows form a category.
- only cell composites that preserve the shape of cell above are allowed. In particular we do not have compositions of horizontal arrows. See [CS10] for more detail.

By *arity* of a cell, we mean the length of the sequence of horizontal arrows which is the source of the cell. In the case of (2.7), the arity is n .

Any double category can be seen as a virtual double category in an obvious way; that is, virtual double category can be seen as a generalisation of double category. In the same way as (2.4), we say a cell α of form (2.1) is *cartesian* if any cell of the form on the right below factors as the left below.

$$(2.8) \quad \begin{array}{ccccc} \cdot & \xrightarrow{\quad} & \cdots & \xrightarrow{\quad} & \cdot \\ k \downarrow & & \bar{\beta} & & \downarrow l \\ X & \xrightarrow{p} & & \xrightarrow{\quad} & Y \\ f \downarrow & & \alpha & & \downarrow g \\ A & \xrightarrow{\quad} & & \xrightarrow{\quad} & B \end{array} = \begin{array}{ccccc} \cdot & \xrightarrow{\quad} & \cdots & \xrightarrow{\quad} & \cdot \\ f k \downarrow & & \beta & & \downarrow g l \\ A & \xrightarrow{\quad} & & \xrightarrow{\quad} & B \end{array}$$

On the other hand, we say a horizontal cell on the left below is *weakly opcartesian* if any cell β on the right below factors as follows:

$$(2.9) \quad \begin{array}{ccccc} \cdot & \xrightarrow{p^0} & \cdots & \xrightarrow{p^{n-1}} & \cdot \\ \parallel & & \alpha & & \parallel \\ X & \xrightarrow{p} & & \xrightarrow{\quad} & Y \\ f \downarrow & & \bar{\beta} & & \downarrow g \\ A & \xrightarrow{\quad} & & \xrightarrow{\quad} & B \end{array} = \begin{array}{ccccc} \cdot & \xrightarrow{\quad} & \cdots & \xrightarrow{\quad} & \cdot \\ f \downarrow & & \beta & & \downarrow g \\ A & \xrightarrow{\quad} & & \xrightarrow{\quad} & B \end{array}$$

Remark 5.8 and Theorem 5.2 in [CS10] shows that if any composable string of horizontal arrows is the source of a weakly opcartesian cell, and weakly opcartesian cells are closed under vertical composite, then \mathbb{X} is a double category.

We write Δ for the category of simplices, and write \mathbf{G}_1 for the category of globes of dimension ≤ 1 , which is interpreted as a subcategory of Δ ; i.e., \mathbf{G}_1 has two objects written as $[0]$ and $[1]$, and has two non-trivial arrows written as $[0] \xrightarrow{\partial_1^1} [1]$ and $[0] \xrightarrow{\partial_0^1} [1]$.

By \emptyset , we mean both the empty set and the empty category, and we write $\mathbf{1}$ for the terminal category.

3. FAMILIAL REPRESENTATION

In this section, we briefly review some concepts surrounding *familial representation*, which is introduced in [Sha21]. Note that, since we define a profunctor $\mathbf{C} \dashv\vdash \mathbf{D}$ to be a functor $\mathbf{C}^{\text{op}} \times \mathbf{D} \rightarrow \mathbf{Set}$, we identify a presheaf with a profunctor to the terminal category $\mathbf{1}$.

Definition 3.1. Let \mathbf{C} and \mathbf{C}' be small categories. A *familial representation* $F = (S_F, E_F) : \mathbf{C}' \dashv\vdash \mathbf{C}$ is a pair of

- a presheaf $S_F : \mathbf{C} \dashv\vdash \mathbf{1}$, or equivalently, a discrete fibration $\text{ty}_F : \int S_F \rightarrow \mathbf{C}$ and
- a profunctor $E_F : \mathbf{C}' \dashv\vdash \int S_F$, or equivalently, a functor $E_F[-] : \int S_F \rightarrow \widehat{\mathbf{C}'}$.

■

Definition 3.2. Let $F : \mathbf{C}' \dashv\vdash \mathbf{C}$ and $F' : \mathbf{C}'' \dashv\vdash \mathbf{C}'$ be familial representations. We define the *composite* $FF' : \mathbf{C}'' \dashv\vdash \mathbf{C}$ as follows:

- The total category of $S_{FF'}$ is

$$(3.1) \quad \int S_{FF'} := \int \left(\widehat{\mathbf{C}'}(E_F[-], S_{F'}) \right)$$

where $\widehat{\mathbf{C}'}(E_F[-], S_{F'}) : (\int S_F)^{\text{op}} \rightarrow \mathbf{Set}$ is the presheaf induced from $E_F[-] : \int S_F \rightarrow \widehat{\mathbf{C}'}$ and $S_{F'} \in \widehat{\mathbf{C}'}$. This presheaf is the same as the right extension $\text{rex}_{E_F} S_{F'} : \int S_F \dashv\vdash \mathbf{1}$ of $S_{F'} : \mathbf{C}' \dashv\vdash \mathbf{1}$ along $E_F : \mathbf{C}' \dashv\vdash \int S_F$. We mean by $F\text{ty}_{F'} : \int S_{FF'} \rightarrow \int S_F$ the discrete fibration corresponding to $\widehat{\mathbf{C}'}(E_F[-], S_{F'})$.

- The discrete fibration $\text{ty}_{FF'} : \int S_{FF'} \rightarrow \mathbf{C}$ is the composite $\text{ty}_F \cdot F\text{ty}_{F'}$ of discrete fibrations.

- $E_{FF'} : \mathbf{C}'' \rightarrow \int S_{FF'}$ is defined as the composite

$$(3.2) \quad \mathbf{C}'' \xrightarrow{E_{F'}} \int S_{F'} \xrightarrow{E_{F\overline{S_{F'}}}} \int S_{FF'}$$

of profunctors where $E_{F\overline{S_{F'}}}(\lambda, \kappa) := \text{colim} \left(\int E_F [F\text{ty}_{F'}(\kappa)] \xrightarrow{f\kappa} \int S_{F'} \xrightarrow{\text{Hom}(\lambda, -)} \mathbf{Set} \right)$ for each $\lambda \in S_{F'}$ and $\kappa : E_F [F\text{ty}_{F'}(\kappa)] \rightarrow S_{F'}$, hence $E_{FF'}(c'', \kappa) \cong \text{colim}_{x \in [E_F [F\text{ty}_{F'}(\kappa)]} E_{F'}(c'', \kappa(x))$.

The *identity* familial representation on \mathbf{C} is defined as the pair $((!_{\mathbf{C}})_*, \text{Id}_{\mathbf{C}})$, where $\text{Id}_{\mathbf{C}}$ is the identity profunctor on \mathbf{C} and $!_{\mathbf{C}} : \mathbf{C} \rightarrow \mathbf{1}$ is the unique functor to the terminal category $\mathbf{1}$. ■

Remark 3.3. For any presheaf $S : \mathbf{C} \rightarrow \mathbf{1}$ and any profunctor $E : \mathbf{C}' \rightarrow \mathbf{C}$, we mean by $\overline{S} : \int S \rightarrow \mathbf{C}$ the familial representation $(S, \text{Id}_{\int S})$ and by $\widetilde{E} : \mathbf{C}' \rightarrow \mathbf{C}$ we mean $((!_{\mathbf{C}})_*, E)$ where $!_{\mathbf{C}}$ is the unique functor from \mathbf{C} to $\mathbf{1}$. The notation $E_{F\overline{S_{F'}}$ for the profunctor defined in Definition 3.2 is justified as the E -part of the composite $F\overline{S_{F'}}$. Any familial representation $F : \mathbf{C}' \rightarrow \mathbf{C}$ factors as $\overline{S_F}\widetilde{E_F}$. ■

Remark 3.4. On the other hand, a presheaf $X : \mathbf{C} \rightarrow \mathbf{1}$ can be identified with a familial representation $\emptyset \rightarrow \mathbf{C}$ since there is precisely one profunctor whose type is $\emptyset \rightarrow \int X$. For a familial representation $F : \mathbf{C} \rightarrow \mathbf{C}'$, the composite $FX : \emptyset \rightarrow \mathbf{C}'$ is the presheaf presented as $\coprod_{\lambda \in S_F(-)} \widehat{\mathbf{C}}(E_F[\lambda], X)$. ■

Definition 3.5. A *morphism of familial representation* $\phi : F \Rightarrow G$ between parallel familial representations is a pair (ϕ^S, ϕ^E) such that

- $\phi^S : S_F \Rightarrow S_G$ is a morphism of presheaves, and
- ϕ^E is a natural isomorphism $E_F[-] \rightarrow E_G[\int \phi^S(-)]$, or equivalently, a cartesian cell below in $\mathbb{P}\text{rof}$.

$$(3.3) \quad \begin{array}{ccc} \mathbf{C}' & \xrightarrow{E_F} & \int S_F \\ \parallel & \phi^E & \downarrow \int \phi^S \\ \mathbf{C}' & \xrightarrow{E_G} & \int S_G \end{array}$$

The composite $\psi \circ \phi : F \Rightarrow H$ of two morphisms $\phi : F \Rightarrow G$ and $\psi : G \Rightarrow H$ is defined as the pairwise composite of natural transformations $(\psi^S \circ \phi^S, \psi^E \circ \phi^E)$, where $\psi^E \circ \phi^E$ is the vertical composite of the cells (3.3). There are obvious identity morphisms, hence they form a category $\mathcal{R}\text{ep}(\mathbf{C}', \mathbf{C})$ of familial representations from \mathbf{C}' to \mathbf{C} . ■

Consider a cell in $\mathbb{P}\text{rof}$

$$(3.4) \quad \begin{array}{ccc} \mathbf{A} & \longrightarrow & \mathbf{C} \\ f \downarrow & \alpha & \downarrow g \\ \mathbf{B} & \longrightarrow & \mathbf{D} \end{array}$$

and its *conjoint* α^* ; the composite of the cells below.

$$(3.5) \quad \begin{array}{ccccc} \mathbf{B} & \xrightarrow{f^*} & \mathbf{A} & \longrightarrow & \mathbf{C} \\ & \parallel & \downarrow f & \alpha & \downarrow g \\ & & \mathbf{B} & \longrightarrow & \mathbf{D} \xrightarrow{g^*} \mathbf{C} \end{array}$$

If f is an identity, then α^* is an isomorphism if and only if α is cartesian. It is straightforward to check cells whose conjoins are isomorphisms are closed under horizontal and vertical compositions.

Definition 3.6 (Whiskerings). Let $\phi : F \Rightarrow G : \mathbf{C}' \rightarrow \mathbf{C}$ and $\phi' : F' \Rightarrow G' : \mathbf{C}'' \rightarrow \mathbf{C}'$ be morphisms of familial representations.

- A morphism $F\phi' : FF' \Rightarrow FG'$ consists of the following

- $\int S_{FF'} \xrightarrow{\int (F\phi')^S} \int S_{FG'}$ is derived from the post-composition $\widehat{\mathbf{C}}'(E_F[-], S_{F'}) \xrightarrow{\phi'^S} \widehat{\mathbf{C}}'(E_F[-], S_{G'})$.
- We write

$$(3.6) \quad \begin{array}{ccc} \int S_{F'} & \xrightarrow{E_{F\overline{S_{F'}}}} & \int S_{FF'} \\ \int \phi'^S \downarrow & (F\phi'^S)^E & \downarrow \int (F\phi')^S \\ \int S_{G'} & \xrightarrow{E_{F\overline{S_{G'}}}} & \int S_{FG'} \end{array}$$

for the cell whose component

$$\left(\overline{F\phi'^S} \right)_{t', \tau}^E : \operatorname{colim}_{x \in \int E_F [F\mathbf{ty}_{F'}(\tau)]} \int S_{F'}(t', \tau(x)) \longrightarrow \operatorname{colim}_{x \in \int E_F [F\mathbf{ty}_{F'}(\tau)]} \int S_{G'}(\phi'^S(t'), \phi'^S \circ \tau(x))$$

is given by the functor $\int \phi'^S$ on maps. $F\mathbf{ty}_{G'}(\phi'^S \circ \tau) = F\mathbf{ty}_{F'}(\tau)$, hence the right hand side is precisely $E_{F\overline{S}_{G'}}(\phi'^S(t'), \phi'^S \circ \tau)$. This conjoint of this cell is an isomorphism since those colimits commutes with the coend defining the composition $E_{F\overline{S}_{F'}} \circ \int \phi'^S$ of profunctors.

The E -part $(F\phi')^E$ is defined as the composite:

$$\begin{array}{ccc} \mathbf{C}'' & \xrightarrow{E_{F'}} \int S_{F'} & \xrightarrow{E_{F\overline{S}_{F'}}} \int S_{FF'} \\ \parallel & \phi'^E \downarrow \int \phi'^S & \left(\overline{F\phi'^S} \right)^E \downarrow \int (F\phi')^S \\ \mathbf{C}'' & \xrightarrow{E_{G'}} \int S_{G'} & \xrightarrow{E_{F\overline{S}_{G'}}} \int S_{FG'} \end{array}$$

which is cartesian.

- A morphism $\phi F' : FF' \Rightarrow GF'$ consists of the following

$$- \int S_{FF'} \xrightarrow{\int (\phi F')^S} \int S_{GF'} \text{ is derived from the pre-composition } \widehat{\mathbf{C}}'(E_F[-], S_{F'}) \xrightarrow{\cong} \widehat{\mathbf{C}}'(E_G[\phi^S-], S_{F'})$$

which can be interpreted as a cartesian morphism $\widehat{\mathbf{C}}'(E_F[-], S_{F'}) \rightarrow \widehat{\mathbf{C}}'(E_G[-], S_{F'})$ in the fibration obtained through the Grothendieck construction of the pseudo-functor $\mathbf{C} \mapsto \widehat{\mathbf{C}} : \mathbf{Cat}^{\text{op}} \rightarrow \mathcal{CAT}$.

- We write

$$(3.7) \quad \begin{array}{ccc} \int S_{F'} & \xrightarrow{E_{F\overline{S}_{F'}}} \int S_{FF'} \\ \parallel & (\phi \overline{S}_{F'})^E \downarrow \int (\phi F')^S \\ \int S_{F'} & \xrightarrow{E_{G\overline{S}_{F'}}} \int S_{GF'} \end{array}$$

for the cartesian cell whose components are the canonical isomorphisms

$$\operatorname{colim}_{x \in \int E_F [F\mathbf{ty}_{F'}(\tau)]} \int S_{F'}(t', \tau(x)) \xrightarrow{\cong} \operatorname{colim}_{y \in \int E_G [\phi^S(F\mathbf{ty}_{F'}(\tau))]} \int S_{G'}(t', \tau \circ (\phi^E)^{-1}(y))$$

induced from the isomorphisms $\int E_G[\phi^S(F\mathbf{ty}_{F'}(\tau))] \xrightarrow{\int (\phi^E)^{-1}_{F\mathbf{ty}_{F'}(\tau)}} \int E_F[F\mathbf{ty}_{F'}(\tau)]$ for each $\tau \in \int S_{FF'}$.

The E -part $(\phi F')^E$ is defined as the composite

$$\begin{array}{ccc} \mathbf{C}'' & \xrightarrow{E_{F'}} \int S_{F'} & \xrightarrow{E_{F\overline{S}_{F'}}} \int S_{FF'} \\ \parallel & \parallel & (\phi \overline{S}_{F'})^E \downarrow \int (\phi F')^S \\ \mathbf{C}'' & \xrightarrow{E_{F'}} \int S_{F'} & \xrightarrow{E_{G\overline{S}_{F'}}} \int S_{GF'} \end{array}$$

■

Proposition 3.7. *Definition 3.2, Definition 3.5, and Definition 3.6 define a bicategory \mathcal{Rep} .*

Remark 3.8. For any small category \mathbf{C} , $\mathcal{Rep}(\emptyset, \mathbf{C})$ is equivalent to the presheaf category $\widehat{\mathbf{C}}$. This is through the assignment $S \mapsto \overline{S}$ remarked in Remark 3.4. Therefore $\mathcal{Rep}(\emptyset, -)$ induces a pseudo functor $\mathcal{Rep} \rightarrow \mathcal{CAT}$ which sends small categories to their presheaf categories. Functors between presheaf categories induced by this pseudo functor are called **familial functors**, and natural transformations between familial functors are in the image of this pseudo functor precisely when they are **cartesian**, where a natural transformation is said to be cartesian if its naturality squares are pullback squares. Moreover, any familial functor is cartesian, i.e. preserving pullbacks.

We often identify a familial functor $\widehat{\mathbf{C}} \rightarrow \widehat{\mathbf{D}}$ with its familial representation $\mathbf{C} \dashrightarrow \mathbf{D}$.

■

Remark 3.9. Recall that any small set of objects in an accessible category is a set of κ -presentable objects for some regular cardinal κ (Corollary 2.3.12 of [MP90]), any presheaf category is accessible, and $\{E_F[\lambda] \mid \lambda \in \int S_F\}$ is small. Therefore, any familial functor $F? = \coprod_{\lambda \in S_F(-)} \widehat{\mathbf{C}}(E_F[\lambda], ?)$ is accessible; that is, it preserves κ -filtered colimits for some κ .

■

Remark 3.10. For any familial representation $F : \mathbf{C}' \dashrightarrow \mathbf{C}$, \mathbf{ty}_F can be seen as the unique morphism $S_F \rightarrow (!\mathbf{C})_*$ between presheaves, and the notation $F\mathbf{ty}_{F'}$ for the functor defined in Definition 3.2 is justified as the whiskering of $\mathbf{ty}_{F'} : S_{F'} \Rightarrow (!\mathbf{C}')_* : \emptyset \dashrightarrow \mathbf{C}'$ with $F : \mathbf{C}' \dashrightarrow \mathbf{C}$.

■

Definition 3.11. A monad in $\mathcal{R}ep$ is identified with the cartesian monad induced by $\mathcal{R}ep(\emptyset, -)$, which is called a *familial monad*. \blacksquare

4. CATEGORICAL STRUCTURES

In this section, we fix a familial monad T on \mathbf{C} . Moreover, we suppose the following two conditions for T

- $E_T[\lambda]$ is *connected* [CCT14] in $\widehat{\mathbf{C}}$, which means $\mathbf{Gph}(T)(E_T[\lambda], -) : \widehat{\mathbf{C}} \rightarrow \mathbf{Set}$ preserves small coproducts.
- \mathbf{C} has no non-trivial endo-morphisms.

For example, the free category monad \mathbf{fc} on \mathbf{G}_1 satisfies these condition since $E_{\mathbf{fc}}[n]$ is connected for each $n \in \mathbb{N} = S_{\mathbf{fc}}([1])$, and \mathbf{G}_1 has no non-trivial endo-morphisms. See the proof of Proposition 4.12 for more detail.

4.1. T -graphs, T -categories, and T -simplicial sets. First of all, let us extend the Grothendieck construction to normal lax functors to $\mathcal{P}rof$. This is due to Section 7 of [Ben00].

Definition 4.1. Let \mathbf{C} be a category and $X : \mathbf{C} \rightarrow \mathcal{P}rof$ be a normal lax functor. The *Grothendieck construction* of X is a category $\mathfrak{f}X$ equipped with a functor $\mathbf{ty}_X : \mathfrak{f}X \rightarrow \mathbf{C}$ defined as follows:

- $\text{Obj}(\mathfrak{f}X) := \coprod_{c \in \mathbf{C}} \text{Obj}(X_c)$.
- For each $c, c' \in \mathbf{C}$, $x \in X_c$, and $x' \in X_{c'}$, $\mathfrak{f}X(x, x') := \coprod_{f: c \rightarrow c' \text{ in } \mathbf{C}} X_f(x, x')$.
We write $t : x \rightarrow_f x'$ when $f : c \rightarrow c'$ in \mathbf{C} , $x \in X_c$, $x' \in X_{c'}$, and $t \in X_f(x, x')$.
- For each $t : x \rightarrow_f x'$ and $t' : x' \rightarrow_{f'} x''$, the composite $t't : x \rightarrow_{f'f} x''$ in $\mathfrak{f}X$ is defined by applying the lax functoriality $\mu_{f, f'} : X_{f'} \circ X_f \Rightarrow X_{f'f}$ to $[t, t'] \in \int^{\bar{x}' \in X_{c'}} X_f(x, \bar{x}') \times X_{f'}(\bar{x}', x'') =: X_{f'} \circ X_f(x, x'')$.

Example 4.2. For any presheaf $S : \mathbf{C}^{\text{op}} \rightarrow \mathbf{Set}$, the corresponding discrete fibration $\int S \rightarrow \mathbf{C}$ is the same as the Grothendieck construction on $\mathbf{C} \xrightarrow{S^{\text{op}}} \mathbf{Set}^{\text{op}} \hookrightarrow \mathcal{C}at^{\text{op}} \xrightarrow{(-)^*} \mathcal{P}rof$, where, for each $i : c \rightarrow c'$ in \mathbf{C} and $\lambda, \lambda' \in \int S$, morphism written as $\lambda \rightarrow_i \lambda'$ in $\mathfrak{f}S = \int S$ is unique if exists, which is denoted by \bar{i}_λ . \blacksquare

Let $\mathbf{2} := \{0 \rightarrow 1\}$ be the 2-element chain. Any profunctor $p : \mathbf{C} \dashv \vdash \mathbf{D}$ can be seen as a functor $\lceil p \rceil : \mathbf{2} \rightarrow \mathcal{P}rof$, hence as a category $\mathfrak{f}\lceil p \rceil$ equipped with a functor $\mathfrak{f}\lceil p \rceil \rightarrow \mathbf{2}$ whose pullback along $* \mapsto 0 : \mathbf{1} \rightarrow \mathbf{2}$ (resp. $* \mapsto 1 : \mathbf{1} \rightarrow \mathbf{2}$) is \mathbf{C} (resp. \mathbf{D}).

Definition 4.3. For each normal lax functors $X : \mathbf{C} \rightarrow \mathcal{P}rof$ and $X' : \mathbf{C}' \rightarrow \mathcal{P}rof$, a *proarrow* between normal lax functors $P : X \dashv \vdash X'$ consists of a profunctor $\text{dom}(P) : \mathbf{C} \dashv \vdash \mathbf{C}'$ and a normal lax functor $|P| : \mathfrak{f}\lceil \text{dom}(P) \rceil \rightarrow \mathcal{P}rof$ whose restriction to \mathbf{C} (resp. \mathbf{C}') is equal to X (resp. X'). One can easily check that the Grothendieck construction $\mathfrak{f}|P| \rightarrow \mathfrak{f}\lceil \text{dom}(P) \rceil \rightarrow \mathbf{2}$ defines another profunctor $\mathfrak{f}P : \mathfrak{f}X \dashv \vdash \mathfrak{f}X'$ equipped with a forgetful natural transformation

$$(4.1) \quad \begin{array}{ccc} \mathfrak{f}X & \xrightarrow{\mathfrak{f}P} & \mathfrak{f}X' \\ \mathbf{ty}_X \downarrow & \mathbf{ty}_P & \downarrow \mathbf{ty}_{X'} \\ \mathbf{C} & \xrightarrow{\text{dom}(P)} & \mathbf{C}' \end{array}$$

which is called the *Grothendieck construction for the proarrow P* . \blacksquare

In detail, a proarrow $P : X \dashv \vdash X'$ consists of

- a profunctor $\text{dom}(P) : \mathbf{C} \dashv \vdash \mathbf{C}'$
- for each $c \in \mathbf{C}$ and $c' \in \mathbf{C}'$, a functor $P_{c, c'} : \text{dom}(P)(c, c') \rightarrow \mathcal{P}rof(X(c), X'(c'))$
- for each $g \in \text{dom}(P)(c, c')$, $f : \bar{c} \rightarrow c$ in \mathbf{C} , and $f' : c' \rightarrow \bar{c}'$ in \mathbf{C}' , natural transformations

$$(4.2) \quad \begin{array}{ccc} X(\bar{c}) \xrightarrow{X(f)} X(c) \xrightarrow{P_{c, c'}(g)} X'(c') & X(c) \xrightarrow{P_{c, c'}(g)} X'(c') \xrightarrow{X'(f')} X'(\bar{c}') \\ \parallel & \lambda_{f, g}^P & \parallel & \parallel & \rho_{g, f'}^P & \parallel \\ X(\bar{c}) \xrightarrow{P_{\bar{c}, c'}(g \cdot f)} X'(c') & X(c) \xrightarrow{P_{c, \bar{c}'}(f' \cdot g)} X'(\bar{c}') \end{array}$$

satisfying suitable coherence conditions.

Definition 4.4. There exists a pseudo (hence normal lax) functor $T^{\mathbf{G}_1} : \mathbf{G}_1 \rightarrow \mathcal{P}rof$ defined as the following diagram.

$$(4.3) \quad \mathbf{C} \xrightarrow[\text{E}_T]{\text{ty}_T^*} \int S_T$$

We define the category of T -**globes** of dimension ≤ 1 , $\mathbf{G}_1(T)$, as the Grothendieck construction on $T^{\mathbf{G}_1}$ i.e. $\mathbf{G}_1(T) := \text{ff}T^{\mathbf{G}_1}$. $\mathbf{G}_1(T)$ is presented as follows:

- objects are those in \mathbf{C} and $\int S_T$.
- morphisms generated by the following maps
 - maps in \mathbf{C} : each copy of $i : c \rightarrow c'$ in \mathbf{C}
 - maps in $\int S_T$: each copy of $\bar{i}_{\lambda'} : \lambda \rightarrow \lambda'$ in $\int S_T$
 - $a : c \rightarrow \lambda$ for each $a \in E_T(c, \lambda)$
 - *target* maps: $\tau_\lambda : \text{ty}_\lambda \rightarrow \lambda$ for each $\lambda \in \int S_T$.

subject to the following

- any commutative diagram in \mathbf{C} and $\int S_T$ commutes
- for each $i : c' \rightarrow c$ in \mathbf{C} , $\bar{j}_{\lambda'} : \lambda \rightarrow \lambda'$ in $\int S_T$, and $a \in E_T(c, \lambda)$, the following diagram commutes

$$(4.4) \quad \begin{array}{ccc} c' & \xrightarrow{E_T(i, \bar{j}_{\lambda'})(a)} & \lambda' \\ i \downarrow & \circlearrowleft & j \uparrow \\ c & \xrightarrow{a} & \lambda \end{array}$$

- for each $\lambda \rightarrow_j \lambda'$ in $\int S_T$, the following diagram commutes

$$(4.5) \quad \begin{array}{ccc} \text{ty}_T(\lambda) & \xrightarrow{\tau_\lambda} & \lambda \\ j \downarrow & \circlearrowleft & j \downarrow \\ \text{ty}_T(\lambda') & \xrightarrow{\tau_{\lambda'}} & \lambda' \end{array}$$

Therefore, $\mathbf{G}_1(T)$ has no non-trivial endo-morphisms if \mathbf{C} is so. ■

Definition 4.5. A T -**graph** is a presheaf on $\mathbf{G}_1(T)$. We write $\mathbf{Gph}(T)$ for the presheaf category $\widehat{\mathbf{G}_1(T)}$. ■

Remark 4.6. For any ordinary monad T on a (possibly large) category \mathbf{E} , a T -**graph** is an endo-span $TX_0 \leftarrow X_1 \rightarrow X_0$. If T is familial then the two definitions of T -graph coincides:

Given a presheaf X on $\mathbf{G}_1(T)$, we obtain a span $TX_0 \leftarrow X_1 \rightarrow X_0$ in $\widehat{\mathbf{C}}$ as follows:

- $\text{ty}_{X_0} : \int X_0 \rightarrow \mathbf{C}$ is the pullback of $\text{ty}_X : \int X \rightarrow \mathbf{G}_1(T)$ along the canonical inclusion $\mathbf{C} \hookrightarrow \mathbf{G}_1(T)$.
- $\int X_1$ is the domain of the pullback $\text{ty}'_X : \int X_1 \rightarrow \int S_T$ of $\text{ty}_X : \int X \rightarrow \mathbf{G}_1(T)$ along the canonical inclusion $\int S_T \hookrightarrow \mathbf{G}_1(T)$, and ty_{X_1} is the composite of ty'_X and $\text{ty}_T : \int S_T \rightarrow \mathbf{C}$.
- for each $\lambda \in \int S_T$ and $\xi \in \int X_1$ over λ with respect to ty'_{X_1} , $\text{src}(\xi) \in \widehat{\mathbf{C}}(E_T[\lambda], X_0) \subset \int TX_0$ is defined as $\text{src}(\xi)_c : a \mapsto a^*\xi$ for each $c \in \mathbf{C}$, where, for each $a \in E_T(c, \lambda)$, $a^*\xi$ is the outcome of reindexing ξ along $a : c \xrightarrow{\partial_1} \lambda$ in $\mathbf{G}_1(T)$ with respect to ty_X .
- for each $\lambda \in \int S_T$ and $\xi \in \int X_1$ over λ with respect to ty'_{X_1} , $\text{tgt}(\xi) \in \int X_0$ is defined as $\tau_\lambda^*\xi$ which is the outcome of reindexing ξ along $\tau_\lambda : \text{ty}_T(\lambda) \xrightarrow{\partial_0} \lambda$ in $\mathbf{G}_1(T)$ with respect to ty_X . ■

Let Δ_a be the **augmented simplex category**, which is presented as follows:

- objects are finite ordinals $[n] := \{0, \dots, n\}$ ($n \geq -1$, $[-1] = \emptyset$)
- morphisms generated by the following two kinds of maps
 - *face maps*: $\partial_i^n : [n-1] \rightarrow [n]$ ($n \geq 0$ and $i \in [n]$)
 - *degeneracy maps*: $\sigma_i^n : [n+1] \rightarrow [n]$ ($n \geq 0$ and $i \in [n]$)

subject to the following *simplicial identities*

$$(4.6) \quad \partial_j^{n+1} \partial_i^n = \partial_i^{n+1} \partial_{j-1}^n \quad i < j$$

$$(4.7) \quad \sigma_j^n \sigma_i^{n+1} = \sigma_i^n \sigma_{j+1}^{n+1} \quad i \leq j$$

$$(4.8) \quad \sigma_j^{n-1} \partial_i^n = \begin{cases} \partial_i^{n-1} \sigma_{j-1}^{n-2} & i < j \\ \text{id}_{[n-1]} & i \in \{j, j+1\} \\ \partial_{i-1}^{n-1} \sigma_j^{n-2} & i > j+1 \end{cases}$$

Δ_a is a monoidal category with the tensor product $[n] \oplus [m] := [n + m + 1]$. By a *simplex*, we mean an object of Δ_a . The **simplex category** Δ is the full subcategory of Δ_a consisting of simplices of dimension greater than -1 . A face map $\partial_i^n : [n-1] \rightarrow [n]$ is said to be **inner** if i is neither 0 nor n . The subcategory of Δ generated from all inner faces and degeneracies is written as Δ_{inn} .

Proposition 4.7. *There exists a surjective-on-object and full functor $\Delta_{\text{inn}} \rightarrow \Delta_a^{\text{op}}$ which sends*

- $[n]$ to $[n-1]$,
- ∂_{i+1}^{n+2} to σ_i^n , and
- σ_i^n to ∂_i^n

for each $n \geq 0$ and $i \in [n]$.

The assignment $F \mapsto S_F$ extends to a functor $S : \mathcal{R}ep(\mathbf{C}', \mathbf{C}) \rightarrow \mathcal{R}ep(\emptyset, \mathbf{C})$ which can be seen as the pre-composition by $(!_{\mathbf{C}'})_* : \emptyset \rightarrow \mathbf{C}'$. On the other hand, T is a monoid in the monoidal category $\mathcal{R}ep(\mathbf{C}, \mathbf{C})$, hence induces a (strong monoidal) functor $\ulcorner T \urcorner : \Delta_a \rightarrow \mathcal{R}ep(\mathbf{C}, \mathbf{C})$. Thus we obtain a pseudo functor $T^\Delta|_{\text{inn}}$ defined as the composite of the following (pseudo) functors

$$(4.9) \quad \Delta_{\text{inn}} \longrightarrow \Delta_a^{\text{op}} \xrightarrow{\ulcorner T \urcorner^{\text{op}}} \mathcal{R}ep(\mathbf{C}, \mathbf{C})^{\text{op}} \xrightarrow{S^{\text{op}}} \mathcal{R}ep(\emptyset, \mathbf{C})^{\text{op}} \xrightarrow{f^{\text{op}}} \mathcal{C}at^{\text{op}} \xrightarrow{(-)^*} \mathcal{P}rof$$

Proposition 4.8. *By setting $T^\Delta(\partial_0^{n+1}) := (T^n \text{ty}_T)^*$ and $T^\Delta(\partial_{n+1}^{n+1}) := E_{T\overline{S_{T^n}}}$ for each $n \geq 0$, $T^\Delta|_{\text{inn}}$ extends to a pseudo functor $T^\Delta : \Delta \rightarrow \mathcal{P}rof$.*

Proof. For each $n \geq 0$ and $i \in [n]$, $T^\Delta(\partial_{i+1}^{n+2}) = (T^{n-i} \mu^S T^i)^* : \int S_{T^{n+1}} \rightarrow \int S_{T^{n+2}}$ and $T^\Delta(\sigma_i^n) = (T^{n-i} \eta^S T^i)^* : \int S_{T^{n+1}} \rightarrow \int S_{T^n}$. Of the remaining simplicial identities, those not involving ∂_{n+1}^{n+1} are trivial, considering identities such as $\text{ty}_T \circ \mu^S = \text{ty}_{T^2}$, $\text{ty}_T \circ \eta^S = \text{id}_{\mathbf{C}}$, and $\text{ty}_T \circ T \text{ty}_T = \text{ty}_{T^2}$.

- The isomorphism corresponding to $\partial_{n+1}^{n+2} \partial_{i+1}^{n+1} = \partial_{i+1}^{n+2} \partial_{n+1}^{n+1}$ ($n > i \geq 0$) is given by a cell whose conjoint is an isomorphism, written as $\left(T(\overline{T^{n-i} \mu T^i})^S \right)^E$, which appears when one defines the post-whiskering of $T^{n-i} \mu T^i$ by T , see (3.6).
- In the same way, the isomorphism corresponding to $\sigma_j^{n+1} \partial_{n+3}^{n+3} = \partial_{n+1}^{n+1} \sigma_j^n$ ($n+1 > j$) is given by $\left(T(\overline{T^{n-j} \eta T^j})^S \right)^E$.
- The isomorphism corresponding to $\sigma_n^n \partial_{n+1}^{n+1} = \text{id}_{[n]}$ is given by the cell appearing when one defines pre-whiskering, $(\eta \overline{S_{T^n}})^E$, see (3.7).
- $T^n \text{ty}_T$ induces a morphism of familial representations $\phi : \overline{S_{T^{n+1}}} \Rightarrow \overline{S_{T^n}}$ (see Remark 3.3). Therefore the post-whiskering $\left(T\overline{\phi^S} \right)^E$ gives rise to an isomorphism for $\partial_{n+1}^{n+2} \partial_0^{n+1} = \partial_0^{n+2} \partial_{n+1}^{n+1}$.

□

Definition 4.9. We define the category of *T-simplices*, $\Delta(T)$, as the total category of the Grothendieck construction on T^Δ . A presheaf on $\Delta(T)$ is called a *T-simplicial set*, and we write $\mathbf{S}Set(T)$ for the presheaf category $\widehat{\Delta(T)}$. ■

Notation 4.10. Let X be a T -graph. For each $n \in \mathbb{N}$, $\lambda_{n+1} \in \int S_{T^{n+1}}$, $x_n \in \int T^n X_0$, $x_{n+1} \in \int T^{n+1} X_0$, and $\xi_n \in \int T^n X_1$, we mean by $x_{n+1} \xrightarrow[\lambda_{n+1}]{\xi_n} x_n$ that there exists $\lambda_n \in \int S_{T^n}$ such that the following diagram commutes in $\widehat{\mathbf{C}}$.

$$(4.10) \quad \begin{array}{ccccc} & & \lambda_{n+1} & & \\ & & \curvearrowright & & \\ & & E_{T^n}[\lambda_n] & & \\ & & \downarrow \xi_n & & \\ & & X_1 & \xrightarrow{\text{tgt}} & X_0 \\ & \swarrow & \leftarrow \text{src} & \searrow & \\ S_T & \xleftarrow{T \text{ty}_{X_0}} & TX_0 & \xleftarrow{\text{src}} & X_1 & \xrightarrow{\text{tgt}} & X_0 \end{array}$$

When $n = 0$, we say $\xi_0 \in \int X_1$ is a λ_1 -**arrow**. A T -graph is completely determined by λ_1 -arrows equipped with their types for all $\lambda_1 \in \int S_T$.

A **path** of shape (λ_{n+m}, n) , or λ_{n+m} -**path** of length $n > 0$, written as $p : x_{n+m} \xrightarrow[\lambda_{n+m}]{\xi_{n+m}} x_m$, is a sequence

$$(4.11) \quad p = (x_{n+m}, \xi_{n+m-1}, x_{n+m-1}, \dots, x_{m+1}, \xi_m, x_m)$$

such that $x_{i+m+1} \xrightarrow[\lambda_{i+m+1}]{\xi_{i+m}} x_{i+m}$ holds for each $i \in [n-1]$. If $\lambda_m = T^m \text{ty}_{X_0}(x_m)$ and $m > 0$, we admits unique λ_m -path of length 0 for each $x_m \in \int T^m X_0$, which is written as $(\)_{x_m} : x_m \xrightarrow[\lambda_m]{\xi_m} x_m$. For each $c \in \mathbf{C}$ and $x_0 \in X_0(c)$, we admits unique c -path $x_0 \xrightarrow{\xi_0} x_0$.

In short, a λ_{n+m} -path is an element of $\int (T^{n+m-1}X \circ \cdots \circ T^m X)_1$ (over λ_{n+m} in a sense), where $T^{n+m-1}X \circ \cdots \circ T^m X$ is the composite of n spans $T^{i+m}X : T^{i+m+1}X_0 \dashrightarrow T^{i+m}X_0$ ($i < n$), and $(\cdots)_1$ means the root of the span. \blacksquare

For any λ_{n+m} -path $p : x_{n+m} \xrightarrow{\lambda_{n+m}} x_m$ as above, $\lambda_m := T^m \mathbf{ty}_{T^n X_0}(x_{n+m})$, $c \in \mathbf{C}$, and $a \in E_{T^m}(c, \lambda_m)$, we can define a $\lambda_{m+n}(a)$ -path $x_{n+m}(a) \xrightarrow{\lambda_{m+n}(a)} x_m(a)$, whose component is written as follows:

$$(x_{n+m}(a), \xi_{n+m-1}(a), x_{n+m-1}(a), \dots, x_{m+1}(a), \xi_m(a), x_m(a))$$

where each of x_{i+m} and ξ_{i+m} is interpreted as a morphism $E_{T^m}[\lambda_m] \rightarrow T^i X_j$.

Definition 4.11. A T -category is a T -graph X equipped with *compositions*, written as \mathbf{comp} : for each λ_n -path $p : x_n \xrightarrow{\lambda_n} x_0$, \mathbf{comp}_n assigns an $\mu^n(\lambda_n)$ -arrow $\mathbf{comp}_n(p) : \mu_{X_0}^n(x_n) \xrightarrow{\mu(\lambda_n)} x_0$, where $\mu^n : T^n \rightarrow T$ is the n -ary multiplication for T . We suppose that \mathbf{comp} satisfies the following coherence condition:

$$(4.12) \quad \mathbf{comp}_2 \left(T\mu^n(\mu^m(x_{n+m})) \xrightarrow[\tau_{\mu^n \cdot \mu^m(\lambda_{n+m})}]{\mathbf{comp}_n(\mu^m(p))} \mu^m(x_m) \xrightarrow[\mu^m(\lambda_m)]{\mathbf{comp}_m(q)} x_0 \right) = \mathbf{comp}_{n+m}(q \circ p)$$

for each composable paths $x_{n+m} \xrightarrow{\lambda_{n+m}} x_m \xrightarrow{\lambda_m} x_0$, where

- $\mathbf{comp}_n(p) : T^m \mu^n(x_{n+m}) \xrightarrow[\tau_{\mu^n(\lambda_{m+n})}]{\mathbf{comp}_n(p)} x_m$ is defined as an arrow obtained by applying \mathbf{comp}_n to all $p(a) : x_{n+m}(a) \xrightarrow{\lambda_{n+m}(a)} x_m(a)$ for each $a \in E_{T^m}[\lambda_{n+m-1}]$.
- $\mu^m(p) : \mu^m(x_{n+m}) \xrightarrow[\mu^m(\lambda_{m+n})]{\mu^m(p)} \mu^m(x_m)$ is defined as a path obtained by μ^m to each component of p as a sequence.

The naturality of μ^m guarantees that

$$\mu^m(T^m \mu^n(x_{n+m})) \xrightarrow[\mu^m \cdot T^m \mu^n(\lambda_{n+m})]{\mu^m(\mathbf{comp}_n(p))} \mu^m(x_m) \text{ is equivalent to } T\mu^n(\mu^m(x_{n+m})) \xrightarrow[\tau_{\mu^n \cdot \mu^m(\lambda_{n+m})}]{\mathbf{comp}_n(\mu^m(p))} \mu^m(x_m). \quad \blacksquare$$

In [Lei99], it is shown that S -categories are cartesian monadic over $\mathbf{Gph}(S)$ provided that S is what is called a *suitable monad*, see [Lei99] or Appendix D of [Lei04]. Although it is not clear that our T is suitable, the proof of this fact in [Lei04] is still valid for monad S on a presheaf category which preserves coproducts but is not necessarily suitable, hence there is a cartesian monad $\mathbf{fc}[T]$ on $\mathbf{Gph}(T)$ whose algebras are T -categories.

In detail, $\mathbf{fc}[T](X)_0$ is defined as X_0 and $\mathbf{fc}[T](X)_1$ is defined as the coproduct $\coprod_{n \in \mathbb{N}} (T^{n-1}X \circ \cdots \circ X)_1$, hence the discrete fibration $\mathbf{ty}'_{\mathbf{fc}[T](X)} : \mathbf{fc}[T](X)_1 \rightarrow \mathbf{fc}[T](X)_0$ (see Remark 4.6) places λ_m -paths of length m over $\mu^m(\lambda_m) \in \mathbf{fc}[T](X)_0$ for each $\lambda_m \in \mathbf{fc}[T](X)_1$.

Note that since μ^m is cartesian, a λ_{n+m} -path $p : x_{n+m} \xrightarrow{\lambda_{n+m}} x_m$ of length n corresponds to a pair

$$(4.13) \quad \left(\lambda_{n+m}, \quad \mu^m(p) : \mu^m(x_{n+m}) \xrightarrow[\mu^m(\lambda_{m+n})]{\mu^m(p)} \mu^m(x_m) \right)$$

where $\mu^m(p)$ is $\mu^m(\lambda_{m+n})$ -path obtained by applying μ^m to each component of p .

Moreover, a λ_{m+1} -arrow $x_{m+1} \xrightarrow{\lambda_{m+1}} x_m$ in $\mathbf{fc}[T](X)$ corresponds to a λ'_{m+n} -path $x_{m+n} \xrightarrow{\lambda'_{m+n}} x_m$ such that $\mu^n(\lambda'_{m+n}) = \lambda_{m+1}$ holds, since $E_{T^m}[\lambda_m]$ is connected.

Therefore an element of $\mathbf{fc}[T]^n(X)_1$ corresponds to an n -times nested path whose target is in X_0 :

- a 0-times nested path is an arrow $x_{m+1} \xrightarrow{\lambda_{m+1}} x_m$ in X and
- an $n + 1$ -times nested path is a sequence $x_{k_s} \xrightarrow[\lambda_{k_s}]{\zeta_{s-1}} x_{k_{s-1}} \cdots x_{k_1} \xrightarrow[\lambda_{k_1}]{\zeta_0} x_m$, where $k_{i+1} \geq k_i$, $k_0 := m$, and ζ_i is an n -times nested path for each $i \in [s-1]$.

The n -ary multiplication $\mathbf{fc}[T]^n(X) \rightarrow \mathbf{fc}[T](X)$ is given by the concatenation of n times nested paths in X , in particular, the unit $\mathbf{fc}[T]^0(X) \rightarrow \mathbf{fc}[T](X)$ sends an arrow to a path of length 1.

It is asserted in [Sha22] that $\mathbf{fc}[T]$ is familial if T is so.

In this paper, we directly give the familial representation of $\mathbf{fc}[T]$.

Proposition 4.12. $\mathbf{fc}[T]$ is familial.

Proof. There are two pseudo (hence normal lax) functors $T^{\mathbf{G}_1} : \mathbf{G}_1 \rightarrow \mathcal{P}rof$ and $\ulcorner \mathbf{1} \urcorner : \mathbf{1} \rightarrow \mathcal{P}rof$ which corresponds to $\mathbf{G}_1(T) \rightarrow \mathbf{G}_1$ and $\mathbf{1} \rightarrow \mathbf{1}$ respectively. We define a proarrow (Definition 4.3) $T^{S_{\mathbf{fc}}} : T^{\mathbf{G}_1} \dashrightarrow \ulcorner \mathbf{1} \urcorner$ as follows:

- The domain profunctor $\mathbf{dom}(T^{S_{\mathbf{fc}}})$ is the presheaf part $S_{\mathbf{fc}} : \mathbf{G}_1 \dashrightarrow \mathbf{1}$ of the familial representation $\mathbf{fc} : \mathbf{G}_1 \dashrightarrow \mathbf{G}_1$ of the ordinary free category monad. $S_{\mathbf{fc}}$ is defined as $S_{\mathbf{fc}}([0]) := \{0\}$ and $S_{\mathbf{fc}}([1]) := \mathbb{N}$, where $\{0\} \cong [0]$ is the terminal set. Note that for each $n \in \mathbb{N}$, the two morphisms $0 \rightrightarrows n$ in \mathbf{fc} can be interpreted as monotone functions $[0] \xrightarrow[\ulcorner n \urcorner]{\ulcorner 0 \urcorner} [n]$ which send 0 to 0 and n respectively.

- For each $[\varepsilon] \in \mathbf{G}_1$, define a function $T_{[\varepsilon]}^{S_{\mathbb{k}}} : S_{\mathbb{k}}([\varepsilon]) \rightarrow \mathcal{P}rof(T^{\mathbf{G}_1}([\varepsilon]), \mathbf{1})$ as follows:
 - $T_{[0]}^{S_{\mathbb{k}}}(\underline{0}) := (!_{\mathbf{C}})_* : \mathbf{C} \rightarrow \mathbf{1}$
 - $T_{[1]}^{S_{\mathbb{k}}}(n) : \int S_T \rightarrow \mathbf{1}$ is the presheaf corresponding to the discrete fibration $\int (\mu^n)^S : \int S_{T^n} \rightarrow \int S_T$, where $(\mu^n)^S$ is the presheaf part of n -th composition $\mu^n : T^n \rightarrow T$.
- The left actions (4.2) are uniquely determined through the universality of the terminal presheaf $T_{[0]}^{S_{\mathbb{k}}}(\underline{0}) = (!_{\mathbf{C}})_* : \mathbf{C} \rightarrow \mathbf{1}$.

$S_{\mathbb{k}[T]} := \mathbb{f}T^{S_{\mathbb{k}}} : \mathbf{G}_1(T) \rightarrow \mathbf{1}$ is the Grothendieck construction on this proarrow, see Definition 4.3. Since the E -part of any morphism of familial representations is cartesian, the canonical forgetful natural transformation (4.1) induces a functor $\int S_{\mathbb{k}[T]} \rightarrow \int S_{\mathbb{k}}$ which corresponds to a normal lax functor $fT^{S_{\mathbb{k}}} : \int S_{\mathbb{k}} \rightarrow \mathcal{P}rof$ defined as follows:

$$(4.14) \quad [0] \xrightarrow[\ulcorner n \urcorner]{\lceil 0 \rceil} [n] \quad \mapsto \quad \mathbf{C} \xrightarrow[\ulcorner E_{T^n} \urcorner]{\text{ty}_{T^n}^*} \int S_{T^n}$$

For $E_{\mathbb{k}[T]} : \mathbf{G}_1(T) \rightarrow \int S_{\mathbb{k}[T]}$, we define a proarrow $T^{E_{\mathbb{k}}} : T^{\mathbf{G}_1} \rightarrow fT^{S_{\mathbb{k}}}$ as follows

- The domain profunctor $\text{dom}(T^{E_{\mathbb{k}}})$ is the E -part $E_{\mathbb{k}} : \mathbf{G}_1 \rightarrow \int S_{\mathbb{k}}$ of the familial representation fc . $E_{\mathbb{k}}$ is defined as $E_{\mathbb{k}}([\varepsilon], a) := \{i \mid i + \varepsilon \in [a]\}$ where $[0] := [0]$. $i \in E_{\mathbb{k}}([\varepsilon], a)$ is interpreted as a monotone function $i + - : [\varepsilon] \rightarrow [a]$ and the left and right actions for $E_{\mathbb{k}}$ are given by pre- and post-compositions.
- For each $[\varepsilon] \in \mathbf{G}_1$ and $a \in \int S_{\mathbb{k}}$, define a function $T_{[\varepsilon], a}^{E_{\mathbb{k}}} : E_{\mathbb{k}}([\varepsilon], a) \rightarrow \mathcal{P}rof(T^{\mathbf{G}_1}([\varepsilon]), fT^{S_{\mathbb{k}}}(a))$ as $T_{[\varepsilon], a}^{E_{\mathbb{k}}}(i) := T^\Delta(i + -)$, hence in detail,
 - $T_{[0], 0}^{E_{\mathbb{k}}}(0) = T_{[0], \underline{0}}^{E_{\mathbb{k}}}(0) := \text{Id}_{\mathbf{C}} : \mathbf{C} \rightarrow \mathbf{C}$ is the identity profunctor on \mathbf{C} .
 - $T_{[1], 1}^{E_{\mathbb{k}}}(0) := \text{Id}_{\int S_T} : \int S_T \rightarrow \int S_T$ is the identity profunctor on $\int S_T$.
 - $T_{[0], n}^{E_{\mathbb{k}}}(i) : \int S_{T^0} = \mathbf{C} \rightarrow \int S_{T^n}$ is the composite of profunctors

$$(4.15) \quad \mathbf{C} \xrightarrow[\ulcorner \quad \urcorner]{E_{T^{n-i}}} \int S_{T^{n-i}} \xrightarrow[\ulcorner \quad \urcorner]{(T^{n-i} \text{ty}_{T^i})^*} \int S_{T^n}$$

- $T_{[1], n}^{E_{\mathbb{k}}}(i) : \int S_T \rightarrow \int S_{T^n}$ is the composite of profunctors

$$(4.16) \quad \int S_T \xrightarrow[\ulcorner \quad \urcorner]{E_{T^{n-i-1} \overline{S_T}}} \int S_{T^{n-i}} \xrightarrow[\ulcorner \quad \urcorner]{(T^{n-i} \text{ty}_{T^i})^*} \int S_{T^n}$$

Moreover, the pseudo-functoriality of T^Δ defines the (isomorphic) actions since $T^{\mathbf{G}_1}(\partial_{1-i}^1) = T^\Delta(i + -) : \mathbf{C} \rightarrow \int S_T$ and $fT^{S_{\mathbb{k}}}(\ulcorner j \urcorner) = T^\Delta(j + -) : \mathbf{C} \rightarrow \int S_{T^n}$ for each $i \in \{0, 1\}$ and $j \in \{0, n\}$.

$E_{\mathbb{k}[T]} := \mathbb{f}T^{E_{\mathbb{k}}} : \mathbf{G}_1(T) \rightarrow \int S_{\mathbb{k}[T]}$ is its Grothendieck construction. In detail, $E_{\mathbb{k}[T]}$ is generated by the following components:

- $\ulcorner i \urcorner \lambda_{n+1, a} \in E_{\mathbb{k}[T]}(T \text{ty}_{T^i} \circ \lambda_{n+1}(a), \lambda_{n+1})$ for each $n \geq 1$, $i \in [n-1]$, $\lambda_{n-i} \in \int S_{T^{n-i}}$, $a \in \int E_{T^{n-i}}[\lambda_{n-i}]$, and $\lambda_{n+1} \in \int S_{T^{n+1}} \cong \int S_{T^{n-i} T^{i+1}}$ which can be interpreted as $\lambda_{n+1} : E_{T^{n-i}}[\lambda_{n-i}] \rightarrow S_{T^{i+1}}$
- $a \in E_{\mathbb{k}[T]}(c, \lambda)$ for each $c \in \mathbf{C}$, $\lambda \in \int S_T$, and $a \in E_T(c, \lambda)$
- $\tau_\lambda \in E_{\mathbb{k}[T]}(\text{ty}_T(\lambda), \lambda)$ for each $\lambda \in \int S_T$
- $\ulcorner n - 1 \urcorner \lambda_n \in E_{\mathbb{k}[T]}(T \text{ty}_{T^{n-1}}(\lambda_n), \lambda_n)$ for each $n > 0$ and $\lambda_n \in \int S_{T^n}$
- $j \in E_{\mathbb{k}[T]}(c', c)$ for each $c' \xrightarrow{j} c$ in $\mathbf{C} = T^{\mathbf{G}_1}([0]) = fT^{S_{\mathbb{k}}}(\underline{0})$, where c is in $fT^{S_{\mathbb{k}}}(\underline{0})$.
- $j \in E_{\mathbb{k}[T]}(c', \lambda_0)$ for each $c' \xrightarrow{j} \lambda_0$ in $\int S_{T^0} = T^{\mathbf{G}_1}([0]) = fT^{S_{\mathbb{k}}}(0)$ where λ_0 is in $fT^{S_{\mathbb{k}}}(0)$.

which are subject to the following conditions:

- the restriction of $E_{\mathbb{k}[T]}$ to $\mathbf{G}_1(T) \rightarrow \mathbf{G}_1(T)$ is identity profunctor, hence morphisms in ii), iii), v), and $\ulcorner 0 \urcorner \lambda_1$ in iv) satisfy obvious commutativity.
- morphisms in vi) defines identity profunctor $T^{\mathbf{G}_1} = \mathbf{C} \rightarrow \mathbf{C} = fT^{S_{\mathbb{k}}}(0)$
- $E_{\mathbb{k}[T]}(T \text{ty}_{T^i} \circ \lambda_{n+1}(\bar{j}_a), \lambda_{n+1}) (\ulcorner i \urcorner \lambda_{n+1, a}) = \ulcorner i \urcorner \lambda_{n+1, a'}$
for each $a' \xrightarrow{\bar{j}_a} a$ in $\int E_{T^{n-i}}[\lambda_{n-i}]$ and $\lambda_{n+1} : E_{T^{n-i}}[\lambda_{n-i}] \rightarrow S_{T^{i+1}}$
- $E_{\mathbb{k}[T]}(T \text{ty}_{T^i} \circ \lambda_{n+1}(a), \bar{j}_{\lambda_{n+1}}) (\ulcorner i \urcorner \lambda'_{n+1, a'}) = \ulcorner i \urcorner \lambda_{n+1, a}$
for each $\lambda'_{n+1} \xrightarrow{\bar{j}_{\lambda_{n+1}}} \lambda_{n+1}$ in $\int S_{T^{n+1}}$, $a' \in \int E_{T^{n-i}}[\lambda'_{n-i}]$, $a := E_{T^{n-1}}[\bar{j}_{\lambda_{n+1}}](a') \in \int E_{T^{n-i}}[\lambda_{n-i}]$, $\lambda_{n+1} : E_{T^{n-i}}[\lambda_{n-i}] \rightarrow S_{T^{i+1}}$, and $\lambda'_{n+1} : E_{T^{n-i}}[\lambda'_{n-i}] \rightarrow S_{T^{i+1}}$. Note that $\lambda_{n+1}(a) = \lambda'_{n+1}(a')$ holds by definition of $S_{T^{n-i} T^{i+1}}$.

Now it is straightforward to check that, for any T -graph X and $\lambda_n \in \int S_{T^n}$, a morphism $\alpha : E_{\mathbb{k}[T]}[\lambda_n] \rightarrow X$ corresponds to a λ_n -path $x_n \xrightarrow[\lambda_n]{\xi_{n-1}} x_{n-1} \rightarrow \dots \rightarrow x_1 \xrightarrow[\lambda_1]{\xi_0} x_0$ in X as follows:

- If $n \geq 1$,
 - $x_0 := \alpha(\ulcorner n - \Gamma_{\lambda_n} \cdot \tau_{\lambda_1} \urcorner)$, where τ_{λ_1} is seen as a morphism in $\mathbf{G}_1(T)$.
 - $x_1 : E_{T^1}[\lambda_1] \rightarrow X_0$ sends $u \in \int E_{T^1}[\lambda_1]$ to $\alpha(\ulcorner n - \Gamma_{\lambda_n} \cdot u \urcorner)$ for $i < n - 1$.
 - $\xi_0 := \alpha(\ulcorner n - \Gamma_{\lambda_n} \urcorner)$.
 - $x_{i+1} : E_{T^{i+1}}[\lambda_{i+1}] \rightarrow X$ sends $\kappa_{\bar{a}}(u) \in E_{T^{i+1}}(c, \lambda_{i+1})$ to $\alpha(\ulcorner n - i - \Gamma_{\lambda_n, \bar{a}} \cdot u \urcorner)$ for $0 < i \leq n - 1$, where $\bar{a} \in \int E_{T^i}[\lambda_i]$, $u \in E_T(c, \lambda_{i+1}(\bar{a}))$, and $\kappa_{\bar{a}} : E_T(c, \lambda_{i+1}(\bar{a})) \rightarrow E_{T^{i+1}}(c, \lambda_{i+1})$ is the coprojection of the colimit defining $E_{TT^i}(c, \lambda_{i+1})$.
 - $\xi_{i+1} : E_{T^{i+1}}[\lambda_{i+1}] \rightarrow X$ sends $a \in \int E_{T^{i+1}}[\lambda_{i+1}]$ to $\alpha(\ulcorner n - i - 2\Gamma_{\lambda_n, a} \urcorner)$ for $i < n - 1$.
- If $n = 0$, $x_0 := \alpha(\text{id}_{\lambda_0})$, where id_{λ_0} is what is introduced in **vi** above.

Therefore, $(S_{\text{fc}[T]}, E_{\text{fc}[T]})$ actually gives a familial representation of the functor $\text{fc}[T]$. Since $\text{fc}[T]$ is a cartesian monad and cartesian natural transformations between familial functors coincide with morphisms of familial representations (**Remark 3.8**), this finishes the proof. \square

Note that a path in a coproduct $\coprod_{i \in I} X^i$ of T -graphs is contained in X^i for some $i \in I$, since each component x_n or $\xi_n : E_{T^n}[\lambda_n] \rightarrow \coprod_{i \in I} X_j^i$ ($j = 0, 1$) is contained in X_i for some $i \in I$ for $E_{T^n}[\lambda_n]$ is connected.

Therefore, since $E_{\text{fc}[T]}[\lambda_m]$ represents λ_m -paths, it is connected for each $\lambda_m \in \int S_{T^m} \subset \int S_{\text{fc}[T]}$, hence $\text{fc}[T]$ satisfies conditions we imposed on T . $\text{fc}^n[T]$ is defined as n -times iteration of $T \mapsto \text{fc}[T]$.

Definition 4.13. The category of T -categories, $\mathbf{Cat}(T)$, is the Eilenberg-Moore category of $\text{fc}[T]$, and morphisms in this category are called T -functors. A T -virtual n -tuple category is defined as an algebra of $\text{fc}^n[T]$, and we write $\mathbf{V}\text{-}n\text{-tplCat}(T)$ for the Eilenberg-Moore category of $\text{fc}^n[T]$. In particular, we write $\mathbf{VDbICat}(T)$ if $n = 2$, which is the category of T -virtual double categories. \blacksquare

Remark 4.14. When $T = \text{id}_1$, then $\mathbf{Cat}(T) = \mathbf{Cat}$ and T -virtual double categories are virtual double categories defined in **[CS10]**, for example.

On the other hand, for $n > 2$, our notion of virtual n -tuple category ($:= \text{id}_1$ -virtual n -tuple category) is *not* consistent with the notion of “virtual widget” proposed in Section 8 of **[CS10]**. \blacksquare

Let $F_{\text{fc}[T]} : \mathbf{Gph}(T) \rightarrow \mathbf{Cat}(T)$ be the free functor. Recall that for ordinary category, the ordinal $[n]$ as a category is the free category of $E_{\text{fc}[n]} := (0 \rightarrow 1 \rightarrow \dots \rightarrow n) \in \mathbf{Gph}$, and Δ is a full subcategory of \mathbf{Cat} . An analogy of this fact for T -category holds:

Theorem 4.15. $\lambda_n \mapsto F_{\text{fc}[T]}(E_{\text{fc}[T]}[\lambda_n])$ induces a fully faithful, dense functor $\Delta(T) \hookrightarrow \mathbf{Cat}(T)$.

To show this theorem, we firstly show the following lemma:

Lemma 4.16. Let $n \leq m$, $\lambda'_m \in S_{T^m}$, and $\lambda_n \in S_{T^n}$. A λ_n -path of length n in the T -graph $E_{\text{fc}[T]}[\lambda'_m]$ precisely corresponds to a pair (l, ζ) of $l \in [m - n]$ and $\zeta \in \text{colim}_{a \in \int E_{T^l}[\lambda'_l]} \int S_{T^n}(\lambda_n, T^{n+l} \text{ty}_{T^{m-n-l}}(\lambda'_m)(a))$, where $\lambda'_l := T^l \text{ty}_{T^{m-l}}(\lambda'_m)$.

Proof. We write $\lambda_i := T^i \text{ty}_{T^{n-i}}(\lambda_n)$ and $\lambda'_j := T^j \text{ty}_{T^{m-j}}(\lambda'_m)$, so that λ_n and λ'_m are interpreted as maps $\lambda_n : E_{T^i}[\lambda_i] \rightarrow S_{T^{n-i}}$ and $\lambda'_m : E_{T^j}[\lambda'_j] \rightarrow S_{T^{m-j}}$ for each $i \in [n]$ and $j \in [m]$. For each $l \in [m - n]$, we show the correspondence between the following data:

- λ_n -path $x_n \xrightarrow{\xi_{n-1}} x_{n-1} \dots x_1 \xrightarrow{\xi_0} x_0$ such that x_0 is in $T_{[0],m}^{E_{\text{fc}}} (m - l)(\lambda_0, \lambda'_m)$
- $\zeta \in \text{colim}_{a \in \int E_{T^l}[\lambda'_l]} \int S_{T^n}(\lambda_n, \lambda'_{n+l}(a))$

Suppose $\zeta \in \text{colim}_{a \in \int E_{T^l}[\lambda'_l]} \int S_{T^n}(\lambda_n, \lambda'_{n+l}(a))$. Note that $\text{colim}_{a \in \int E_{T^l}[\lambda'_l]} \int S_{T^n}(\lambda_n, \lambda'_{n+l}(a))$ is the same as $T^\Delta((m - l - n) + -)(\lambda_n, \lambda'_m)$, where $(m - l - n) + - : [n] \rightarrow [m]$ is a map in Δ . Now we obtain a morphism of T -graphs $E_{\text{fc}[T]}[\lambda_n] \rightarrow E_{\text{fc}[T]}[\lambda'_m]$ by the post-composition of $\zeta : \lambda_n \rightarrow \lambda'_m$ in $\Delta(T)$, but we have already checked that such a morphism corresponds to a λ_n -path in $E_{\text{fc}[T]}[\lambda'_m]$ (see the proof of **Proposition 4.12**).

Moreover, if $n > 0$, since $x_0 := \zeta \cdot \ulcorner n - \Gamma_{\lambda_n} \cdot \tau_{\lambda_1} \urcorner$ and the maps $\lambda_0 \xrightarrow{\tau_{\lambda_1}} \lambda_1 \xrightarrow{\ulcorner n-1 \urcorner \lambda_n} \lambda_n \xrightarrow{\zeta} \lambda'_m$ in $\Delta(T)$ are over $[0] \xrightarrow{1+} [1] \xrightarrow{(n-1)+} [n] \xrightarrow{(m-n-l)+} [m]$, x_0 is in $T_{[0],m}^{E_{\text{fc}}} (m - l)(\lambda_0, \lambda'_m)$. $n = 0$ case is trivial.

It suffices to prove this assignment $\zeta \mapsto (x_n \xrightarrow{\xi_{n-1}} x_{n-1} \dots x_1 \xrightarrow{\xi_0} x_0)$ defines a bijection between i) and ii), by induction on n . It is trivial ζ itself gives x_0 , hence the assignment is a bijection when $n = 0$. In the same way, when $n = 1$, $\zeta = \xi_0$ gives rise to a bijection. Suppose the assignment is a bijection for $n > 0$ and let

$x_{n+1} \xrightarrow{\xi_n} x_n \xrightarrow{\xi_{n-1}} x_{n-1} \dots x_1 \xrightarrow{\xi_0} x_0$ be a λ_{n+1} -path. We suppose there exists $\zeta' \in \text{colim}_{a \in \int E_{T^l}[\lambda'_l]} \int S_{T^n}(\lambda_n, \lambda'_{n+l}(a))$

which corresponds to $x_n \xrightarrow{\xi_{n-1}} x_{n-1} \dots x_1 \xrightarrow{\xi_0} x_0$. We see ζ' as a morphism $\zeta' : \lambda_n \xrightarrow{0+} \lambda'_{n+l}$ in $\Delta(T)$, which factors

as $\lambda_n \xrightarrow{[n]} \lambda'_{n+l}(x) \xrightarrow{0+} \lambda'_{n+l}$ for some $x \in \int E_{T^l}[\lambda'_l]$. ζ' is also considered the composite $\lambda_n \xrightarrow{[n]} \lambda'_{n+l} \xrightarrow{\tau_{(m-n-l)+}} \lambda'_m$,

where $\lambda'_{n+l} \xrightarrow[(m-n-l)+-]{\tau} \lambda'_m$ is the identity morphism on λ'_{n+l} in $\int S_{T^{n+l}}$, hence is one of the generating elements of

$$T^\Delta((m-n-l)+- : [n+l] \rightarrow [m]) = \int S_{T^{n+l}}(\text{id}, T^{n+l} \mathbf{ty}_{T^{m-n-l}})$$

which is sometimes omitted.

Let $b : c_b \xrightarrow{0+-} \lambda_n$ in $\Delta(T)$, or equivalently, $b \in E_{T^n}(c_b, \lambda_n)$ for some $c_b \in \mathbf{C}$. There exists $\bar{x}_n(b) : c_b \xrightarrow{0+-} \lambda'_{n+l}(x)$

which factors $x_n(b)$ as $c_b \xrightarrow[\text{0+-}]{\bar{x}_n(b)} \lambda_{n+l}(x) \xrightarrow{\text{0+-}} \lambda_{n+l}$, which is defined as follows:

Since we have assumed that x_n is induced from post-composition of ζ' , if $n = 1$, $x_n(b) = x_1(b) = \zeta' \cdot b$ holds. Even if $n > 1$, b factors as $c_b \xrightarrow[\text{0+-}]{e'} \lambda_n(e) \xrightarrow[\text{0+-}]{\Gamma^{\text{0}} \lambda_n, e} \lambda_n$ for some $e \in \int E_{T^{n-1}}[\lambda_{n-1}]$, hence $x_n(b) = \zeta' \cdot (\Gamma^{\text{0}} \lambda_n, e \cdot e')$ holds, but moreover $x_n(b) = \zeta' \cdot b$ since the actions for $E_{\text{fc}[T]}$ is defined by restricting composites in $\Delta(T)$. Let $\bar{x}_n(b)$ be the composite $\zeta' \cdot b$.

On the other hand, for each $b : c_b \xrightarrow{0+-} \lambda_n$, there exist $a_b \in \int E_{T^{n+l}}[\lambda'_{n+l}]$ and $\tilde{\xi}_n(b) : \lambda_{n+l}(b) \xrightarrow{[1]} \lambda'_{n+l+1}(a_b)$ such that $\xi_n(b)$ is the composite

$$\lambda_{n+l}(b) \xrightarrow[\text{[1]}]{\tilde{\xi}_n(b)} \lambda'_{n+l+1}(x) \xrightarrow[\text{0+-}]{\Gamma^{\text{0}} \lambda'_{n+l+1}, a_b} \lambda'_{n+l+1}$$

T^Δ sends the commutative square

$$(4.17) \quad \begin{array}{ccc} [0] & \xrightarrow{0+-} & [m] \\ k+- \downarrow & & \downarrow k+- \\ [k] & \xrightarrow{0+-} & [m+k] \end{array}$$

to an isomorphism which makes the following diagram commutes in $\Delta(T)$ (see the proof of [Proposition 4.8](#)):

$$(4.18) \quad \begin{array}{ccc} c & \xrightarrow{a} & \tilde{\lambda}_m \\ \tau \downarrow & & \downarrow \tau \\ \tilde{\lambda}_{m+k}(a) & \xrightarrow[\text{0+-}]{\Gamma^{\text{0}} \tilde{\lambda}_{k+m}, a} & \tilde{\lambda}_{k+m} \end{array}$$

for arbitrary $c \in \mathbf{C}$, $\tilde{\lambda}_{m+k} \in \int S_{T^{m+k}}$, and $\tilde{\lambda}_m := T^m \mathbf{ty}_{T^k}(\tilde{\lambda}_{m+k})$.

In summary, we have the following commutative diagrams

$$(4.19) \quad \begin{array}{ccccc} c_b & \xrightarrow[\text{[0]}]{\mathbf{ty}_T(\tilde{\xi}_n(b))} & c_{a_b} & \xrightarrow{a_b} & \lambda'_{n+l} \\ \tau \downarrow & & \downarrow \tau & & \downarrow \tau \\ \lambda_{n+l}(b) & \xrightarrow[\text{[1]}]{\tilde{\xi}_n(b)} & \lambda_{n+l+1}(a) & \xrightarrow[\text{0+-}]{\Gamma^{\text{0}}} & \lambda'_{n+l+1} \\ & \searrow \xi_n(b) & \nearrow & & \end{array} \quad \begin{array}{ccc} & \lambda'_{n+l}(x) & \\ \bar{x}_n(b) \nearrow & & \searrow \\ c_b & \xrightarrow{x_n(b)} & \lambda'_{n+l} \\ \tau \downarrow & & \downarrow \tau \\ \lambda_{n+l}(b) & \xrightarrow{\xi_n(b)} & \lambda'_{n+l+1} \end{array}$$

where the square on the right hand side commutes since x_n is the codomain of the λ_n -arrow ξ_n . Since any map over $1+- : [0] \rightarrow [n+l+1]$ uniquely factors through $\tau : \lambda'_{n+l} \rightarrow \lambda'_{n+l+1}$, the following square commutes:

$$(4.20) \quad \begin{array}{ccc} c_b & \xrightarrow{\bar{x}_n(b)} & \lambda'_{n+l}(x) \\ \mathbf{ty}_T(\tilde{\xi}_n(b)) \downarrow & & \downarrow \\ c_{a_b} & \xrightarrow{a_b} & \lambda'_{n+l} \end{array}$$

Therefore, $\lambda'_{n+l+1}(x)(\bar{\zeta}' \cdot b) = \lambda'_{n+l+1}(x)(\bar{x}_n(b)) = \lambda'_{n+l+1}(a_b \cdot \mathbf{ty}_T(\tilde{\xi}_n(b))) = S_{T^{n+l}}(\mathbf{ty}_T(\tilde{\xi}_n(b)))(\lambda'_{n+l+1}(a_b)) = \lambda_{n+l+1}(b)$, where λ_{n+l+1} is interpreted as a morphism $E_{T^l}[\lambda'_l] \rightarrow S_{T^{n+l}}$ in the first and second terms and as a morphism $E_{T^{n+l}}[\lambda'_{n+l}] \rightarrow S_T$ in the third term. This means the commutativity of the following:

$$(4.21) \quad \begin{array}{ccc} E_{T^n}[\lambda_n] & \xrightarrow{\lambda_{n+1}} & S_T \\ E_{T^n}[\zeta'] \downarrow & \nearrow & \\ E_{T^n}[\lambda'_{n+l+1}(x)] & & \end{array}$$

i.e. $\bar{\zeta}'$ extends to a morphism $\bar{\zeta} : \lambda_{n+1} \rightarrow \lambda'_{n+l+1}(x)$ in $\int S_{T^{n+l+1}}$.

Let $\zeta : \lambda_{n+1} \xrightarrow{0+-} \lambda'_{n+l+1} \in \text{colim}_{x \in \int E_{T^l}[\lambda'_l]} \int S_{T^{n+l+1}}(\lambda_{n+1}, \lambda'_{n+l+1}(x))$ be what is represented by $\bar{\zeta}$. It is straightforward to check this precisely induces $\xi_n(b)$ through the post-composition $\zeta \cdot b$. \square

proof of Theorem 4.15. For $\bar{\lambda}_n \in \int S_{T^n}$ and $\lambda'_m \in \int S_{T^m}$, it suffices to show the correspondence between the following data:

- i) $\bar{\lambda}_n$ -path $z_n \xrightarrow{\zeta_n} z_{n-1} \dots z_1 \xrightarrow{\zeta_0} z_0$ in $\mathbf{fc}[T](E_{T^n}[\lambda'_m])$
- ii) $p : \bar{\lambda}_n \rightarrow \lambda'_m$ in $\Delta(T)$

The discussion preceding Proposition 4.12 shows that a $\bar{\lambda}_n$ -path $z_n \xrightarrow{\zeta_n} z_{n-1} \dots z_1 \xrightarrow{\zeta_0} z_0$ in $\mathbf{fc}[T]E_{\mathbf{fc}[T]}[\lambda'_m]$ corresponds to a 2-times nested path $x_{k_n} \xrightarrow{\zeta_n} x_{k_{n-1}} \dots x_{k_1} \xrightarrow{\zeta_0} x_0$ in $E_{\mathbf{fc}[T]}[\lambda'_m]$ such that $\mu^{k_n-n+1}(\lambda_{k_n}) = \lambda_n$. We write $\bar{\zeta}$ for the concatenation $x_{k_n} \xrightarrow{\zeta_n} x_0$, which can be seen as a morphism $\lambda_{k_n} \xrightarrow{0+} \lambda'_{k_n+l} \xrightarrow{\tau} \lambda'_m$ for some l .

For each sequence $k_n \geq k_{n-1} \geq \dots k_1 \geq k_0 = 0$, define a monotone function $u_k : [n] \rightarrow [k_n]$ by $i \mapsto k_n - k_{n-i}$, of which the image by T^Δ , $T^\Delta(u_k) : \int S_{T^n} \rightarrow \int S_{T^{k_n}}$, is isomorphic to $(\mu^{r_n} \cdot \mu^{r_{n-1}} \dots \mu^{r_1})^* : \int S_{T^n} \rightarrow \int S_{T^{r_n} T^{r_{n-1}} \dots T^{r_1}}$, where $r_{n-i} := k_{i+1} - k_i$ for each $i < n$. Therefore for arbitrary $\lambda_{k_n} \in \int S_{T^{k_n}}$ and $\bar{\lambda}_s \in \int S_{T^s}$, a map $\bar{\lambda}_s \xrightarrow{u_k} \lambda_{k_n}$ over u_k uniquely factors through a map $\mu^{r_n, \dots, r_1}(\lambda_{k_n}) \xrightarrow{\nu_k} \lambda_{k_n}$.

Let $p : \lambda_n \rightarrow \lambda'_m$ be a morphism in $\Delta(T)$ over $f : [n] \rightarrow [m]$, which sends $i \in [n]$ to $f_i \in [m]$. Let l be $m - f_n$ and k_i be $f_n - f_{n-i}$ for each $i \in [n]$.

We define a 2-times nested path $x_{k_n} \xrightarrow{\zeta_n} x_{k_{n-1}} \dots x_{k_1} \xrightarrow{\zeta_0} x_0$ satisfying $p = \bar{\zeta} \cdot \nu_k$.

Since $l = m - f_n$, p factors through $\lambda'_{k_n+l} \xrightarrow{\tau} \lambda'_m$. Moreover, since there is a commutative diagram

$$(4.22) \quad \begin{array}{ccc} [k_n] & \xrightarrow{u_k} & [n] \\ 0+- \downarrow & & 0+- \downarrow \\ [k_n+l] & \xrightarrow{\text{id}_l \oplus u_k} & [n+l] \end{array}$$

p uniquely factors through $T^l \mu^{r_n, \dots, r_1}(\lambda'_{k_n+l}) \xrightarrow{\nu_k} \lambda'_{k_n+l}$. Therefore, there exists a $x \in \int E_{T^l}[\lambda'_l]$ such that p factors through $\mu^{r_1, \dots, r_n} \cdot \lambda'_{k_n+l}(x) \xrightarrow{0+-} \lambda'_{k_n+l}$, and since $\mu^{r_n, \dots, r_1} : \int S_{T^{k_n}} \rightarrow \int S_{T^n}$ is a discrete fibration, there exists a λ_{k_n} such that $\mu^{r_n, \dots, r_1}(\lambda_{k_n}) = \bar{\lambda}_n$ and p factors through $\bar{\lambda}_n \xrightarrow{\nu_k} \lambda_{k_n}$. This λ_{k_n} does not depend on the choice of x since we have assumed that \mathbf{C} has no non-trivial endo-morphisms and hence so is $\int S_{T^n}$ and all possible candidates of λ_{k_n} is connected through maps over a fixed element $\bar{\lambda}_n$ with respect to μ^{r_n, \dots, r_1} , which is a discrete fibration.

Now we have a map $\bar{\zeta} : \lambda_{k_n} \xrightarrow{0+-} \lambda'_{k_n+l}$, which yields a path $x_{k_n} \xrightarrow{0+-} 0$, which is uniquely decomposed to a 2-times nested path since μ^\bullet are cartesian through the sequence $k_n \geq \dots k_0 = 0$.

Thus we obtain a fully faithful functor $[-]_T : \Delta(T) \hookrightarrow \mathbf{Cat}(T)$. The induced nerve functor $\mathbf{Cat}(T)([-]_T, ?) : \mathbf{Cat}(T) \rightarrow \mathbf{SSet}(T)$ is faithful, since maps of T -categories are completely determined by assignments on paths. It is also full since reindexing by inner faces and degeneracies in $\Delta(T)$ means compositions of paths, and a T -graph morphism commutes with those if and only if it is a T -functor. \square

On the other hand, since

- any familial monad is accessible as an endo-functor on $\widehat{\mathbf{C}}$ (Remark 3.9),
- $\mathbf{fc}[T]$ is familial by Proposition 4.12,
- and the Eilenberg-Moore category for an accessible monad on any locally presentable category is locally presentable (Theorem 5.5.9. of [Bor94]),

$\mathbf{Cat}(T)$ is locally presentable, hence cocomplete as a category. Therefore, $\mathbf{Cat}(T)$ can be seen as a reflective full subcategory of $\mathbf{SSet}(T)$.

For each category X , $(\mathbf{ty}_{T^\Delta}(X))_n$ is $\coprod_{a \in X_n} S_{T^n}$, whose elements are pairs of paths in X and $\lambda_n \in \int S_{T^n}$. On the other hand, a λ_{n+1} -arrow $x_{n+1} \xrightarrow{\xi_n} x_n$ in the underlying T -graph of $\mathbf{ty}_{T^\Delta}(X)$ is a pair $(|\xi_n|, \lambda_{n+1})$, where $|\xi_n| \in X_1$ is an arrow in X such that for each $a \in \int E_{T^n}[\lambda_n]$ and $a' \in \int E_{T^{n+1}}[\lambda_{n+1}]$, $x_n(a) = \mathbf{src}(|\xi_n|)$ and $x_{n+1}(a') = \mathbf{tgt}(|\xi_n|)$, since $E_{T^n}[\lambda_n]$ is connected. Therefore a path of length n whose target is in X_0 in this T -graph precisely corresponds to an element of $(\mathbf{ty}_{T^\Delta}(X))_n$. This shows that $\mathbf{ty}_{T^\Delta}(X)$ is a T -category, hence we have a functor $\nabla_T(X) : \mathbf{Cat} \rightarrow \mathbf{Cat}(T)$.

4.2. Structures of T -categories. In the framework of [CS10], for a cartesian monad \mathbf{S} on \mathbf{E} , an \mathbf{S} -category is an example of what they call a $\mathbf{Span}(\mathbf{S})$ -monoid, where $\mathbf{Span}(\mathbf{S})$ is the outcome of extending \mathbf{S} to a monad on the double category of spans in \mathbf{E} , $\mathbf{Span}(\mathbf{E})$. In general, for any virtual double category \mathbb{X} and monad S on \mathbb{X} , they define another virtual double category of S -monoids, $\mathbb{KMod}(\mathbb{X}, S)$.

Definition 4.17. We write $\mathbb{P}\text{rof}(T)$ for $\mathbb{K}\text{Mod}(\mathbb{S}\text{pan}(\widehat{\mathbf{C}}), \mathbb{S}\text{pan}(T))$. A T -**profunctor** is a horizontal cell in $\mathbb{P}\text{rof}(T)$.

$\mathbb{P}\text{rof}(T)$ has units (Proposition 5.5 of [CS10]), hence one obtains its vertical 2-category (Proposition 6.1 of [CS10]), $\text{Cat}(T)$, of $\mathbb{P}\text{rof}(T)$. We write $\mathcal{V}\text{DblCat}(T)$ for $\text{Cat}(\text{fc}[T])$ and $\mathcal{V}\text{-}n\text{-tplCat}(T)$ for $\text{Cat}(\text{fc}^n[T])$. In detail, $\mathbb{P}\text{rof}(T)$ consists of the following data

- objects are T -categories, and vertical arrows are morphisms in $\mathbf{Cat}(T)$, T -functors.
- for each T -categories X and Y , a horizontal arrow $p : X \dashrightarrow Y$, a T -profunctor, consists of
 - a span $TX_0 \leftarrow |p| \rightarrow Y_0$, i.e., a horizontal arrow $TX_0 \xrightarrow{p} Y_0$ in $\mathbb{S}\text{pan}(\widehat{\mathbf{C}})$, and
 - left and right actions: cells in the double category $\mathbb{S}\text{pan}(\widehat{\mathbf{C}})$

$$(4.23) \quad \begin{array}{ccccc} T^2X_0 & \xrightarrow{TX} & TX_0 & \xrightarrow{p} & Y_0 & & T^2X_0 & \xrightarrow{Tp} & TY_0 & \xrightarrow{Y} & Y_0 \\ \downarrow \mu_{X_0} & & \lambda^p & & \parallel & & \downarrow \mu_{X_0} & & \rho^p & & \parallel \\ TX_0 & \xrightarrow{p} & Y_0 & & TX_0 & \xrightarrow{p} & Y_0 \end{array}$$

where $X : TX_0 \dashrightarrow X_0$ and $Y : TY_0 \dashrightarrow Y_0$ are spans defining underlying T -graphs of X and Y . The compositions of X and Y can be seen as cells in $\mathbb{S}\text{pan}(\widehat{\mathbf{C}})$: e.g.

$$(4.24) \quad \begin{array}{ccc} T^2X_0 & \xrightarrow{TX} & TX_0 & \xrightarrow{X} & X_0 \\ \downarrow \mu_{X_0} & & \text{comp}^2 & & \parallel \\ TX_0 & \xrightarrow{X} & X_0 \end{array}$$

We suppose that λ^p and ρ^p are compatible with those compositions.

- a cell

$$(4.25) \quad \begin{array}{ccccc} X^0 & \xrightarrow{p^0} & X^1 & \xrightarrow{p^1} & \dots & \xrightarrow{p^{n-1}} & X^n \\ \downarrow f & & \alpha & & & & \downarrow g \\ Y^0 & \xrightarrow{q} & Y^1 \end{array}$$

in $\mathbb{P}\text{rof}(T)$ is a cell

$$(4.26) \quad \begin{array}{ccccc} T^n X_0^0 & \xrightarrow{T^{n-1}p^0} & T^{n-1} X_0^1 & \xrightarrow{T^{n-2}p^1} & \dots & \xrightarrow{p^{n-1}} & X_0^n \\ \downarrow Tf_0 \cdot \mu^n & & \alpha & & & & \downarrow g_0 \\ TY_0^0 & \xrightarrow{q} & Y_0^1 \end{array}$$

in $\mathbb{S}\text{pan}(\widehat{\mathbf{C}})$ which is compatible with the ‘‘arrow’’ part $f_1 : X_1^0 \rightarrow Y_1^0$ and $g_1 : X_1^n \rightarrow Y_1^1$ of f and g , and the left and right actions of p^i ($i \in [n-1]$) and q .

We write $p^{n-1} \vee \dots \vee p^1 \vee p^0 : TX_0^0 \dashrightarrow X_0^n$ for the composite

$$(4.27) \quad TX_0^0 \xrightarrow{\mu^*} T^n X_0^0 \xrightarrow{T^{n-1}p^0} \dots \xrightarrow{p^{n-1}} X_0^n$$

in $\mathbb{S}\text{pan}(\widehat{\mathbf{C}})$. It is straightforward to show this span extends to a T -profunctor $X^0 \dashrightarrow X^n$ by equipping the left action of X^0 and the right action of X^n . Right and left actions define 2^{n-1} cells of the form

$$(4.28) \quad \begin{array}{ccc} X^0 & \xrightarrow{p^{n-1} \vee \text{Id}_{X^{n-1}} \vee \dots \vee \text{Id}_{X^2} \vee p^1 \vee \text{Id}_{X^1} \vee p^0} & X^n \\ \parallel & & \parallel \\ X^0 & \xrightarrow{p^{n-1} \vee \dots \vee p^1 \vee p^0} & X^n \end{array}$$

and any cell α can be seen as a cell

$$(4.29) \quad \begin{array}{ccc} X^0 & \xrightarrow{p^{n-1} \vee \dots \vee p^1 \vee p^0} & X^n \\ \downarrow & \bar{\alpha} & \downarrow \\ Y^0 & \xrightarrow{q} & Y^1 \end{array}$$

equalizing those cells, where Id_Y , the *identity T -profunctor*, is the span defining underlying graph of Y , $TY_0 \leftarrow Y_1 \rightarrow Y_0$, equipped with 2-ary compositions of Y as its left and right actions, for each T -category Y . Since μ^n is cartesian, \vee is associative up to isomorphism: $(p^{n-1} \vee \dots \vee p^1 \vee p^0 : TX_0^0 \dashrightarrow X_0^n) \vee (q^{m-1} \vee \dots \vee q^1 \vee q^0 : TY_0^0 \dashrightarrow Y_0^m)$ is invertible to $(p^{n-1} \vee \dots \vee p^1 \vee p^0 \vee q^{m-1} \vee \dots \vee q^1 \vee q^0 : TX_0^0 \dashrightarrow Y_0^m)$ both in $\mathbb{S}\text{pan}(T)$ and $\mathbb{P}\text{rof}(T)$, where $Y^0 = X^n$.

$Cat(T)$ consists of the following data:

- The underlying category is $\mathbf{Cat}(T)$.
- A 2-cell $\sigma : f \Rightarrow g : X \rightarrow Y$ is a cell

$$(4.30) \quad \begin{array}{ccc} & X & \\ f \swarrow & & \searrow g \\ & \sigma & \\ Y \xrightarrow{\text{Id}_Y} & & Y \end{array}$$

The identity T -profunctor satisfies some suitable property and called *unit* in [CS10], and this property implies that σ can be seen as a cell

$$(4.31) \quad \begin{array}{ccc} X & \xrightarrow{\text{Id}_X} & X \\ f \downarrow & \sigma & \downarrow g \\ Y & \xrightarrow{\text{Id}_Y} & Y \end{array}$$

■

A T -profunctor $p : X \dashv\dashv Y$ can be seen as a T -category $\lceil p \rceil$ defined as follows

- $\lceil p \rceil_0$ is the coproduct $X_0 \sqcup Y_0$
- $\lceil p \rceil_1$ is the coproduct $X_1 \sqcup |p| \sqcup Y_1$, and a λ -arrow $\xi_1 : a_1 \rightarrow_{\lambda} a_0$ in $\lceil p \rceil$ is one of the following
 - a λ -arrow $\xi_1 : x_1 \rightarrow_{\lambda} x_0$ in X
 - a λ -arrow $\xi_1 : y_1 \rightarrow_{\lambda} y_0$ in Y
 - ξ_1 is an element of $\int |p|$, $a_1 = \mathbf{src}_{|p|}(\xi_1)$, and $a_0 = \mathbf{tgt}_{|p|}(\xi_1)$, where $\mathbf{src}_{|p|}$ and $\mathbf{tgt}_{|p|}$ are the left and right legs of the underlying span of p .
- since $E_{T^n}[\lambda_n]$ is connected, a λ_n -path is either contained in X or Y , or of the form

$$(4.32) \quad x_n \xrightarrow[\lambda_n]{\xi_{n-1}} x_{n-1} \xrightarrow[\lambda_{n-1}]{} \cdots \xrightarrow{} x_{m+1} \xrightarrow[\lambda_{m+1}]{\bar{\xi}_m} y_m \xrightarrow[\lambda_m]{} \cdots \xrightarrow[\lambda_1]{\xi_0} y_0$$

where all but $\bar{\xi}_m$ are either contained in X or Y . The composite of such a λ_n -path is defined by applying composition of X and Y for its parts contained in X and Y respectively, and applying left and right actions to those composites with $\bar{\xi}_m$.

On the other hand, $\nabla_T[1]$ consists of the following data:

- $(\nabla_T[1])_0$ is the coproduct of two terminal presheaves $1_{\mathbf{C}} \sqcup 1_{\mathbf{C}}$, whose elements are written as pairs $(c, [i])$ where $i = 0, 1$ and $c \in \mathbf{C}$. Since T preserves coproducts, elements in $T(\nabla_T[1])_0$ is also written as (λ, i) .
- $(\nabla_T[1])_1$ is the coproduct $S_T \sqcup S_T \sqcup S_T$, whose elements are written as pairs $(\lambda, 00)$, $(\lambda, 01)$, and $(\lambda, 11)$. For each $(\lambda, ij) \in (\nabla_T[1])_1$, $\mathbf{src}(\lambda) = (\lambda, i)$ and $\mathbf{tgt}(\lambda) = (\mathbf{ty}_T(\lambda), j)$.

which is the same data as the identity T -profunctor on the terminal T -category, $\nabla_T[0]$.

Thus, $\nabla_T[1]$ classifies T -profunctors, i.e., a T -profunctor can be seen as a T -functor $\lceil p \rceil \rightarrow \nabla_T([1])$ whose pullbacks along $\nabla_T(\partial_1^1)$ and $\nabla_T(\partial_0^1)$ are the unique maps $X \rightarrow 1$ and $Y \rightarrow 1$ respectively. Moreover, identity T -profunctors are images of the pullback along $\nabla_T(\sigma^0) : \nabla_T[1] \rightarrow \nabla_T[0]$.

A 2-ary horizontal cell

$$(4.33) \quad \begin{array}{ccccc} X^0 & \xrightarrow{p_{01}} & X^1 & \xrightarrow{p_{12}} & X^2 \\ \parallel & & \alpha & & \parallel \\ X^0 & \xrightarrow[p_{02}]{} & & \xrightarrow{} & X^2 \end{array}$$

can be seen as a T -functor $\lceil p \rceil \rightarrow \nabla_T[2]$ as follows

- $\lceil p \rceil_0$ is the coproduct $X_0^1 \sqcup X_0^1 \sqcup X_0^2$
- $\lceil p \rceil_1$ is the coproduct $X_1^0 \sqcup X_1^1 \sqcup X_1^2 \sqcup |p_{01}| \sqcup |p_{12}| \sqcup |p_{02}|$, and a λ -arrow $\xi_1 : x_1 \rightarrow_{\lambda} x_0$ in $\lceil p \rceil$ is either of the following
 - a λ -arrow $\xi_1 : x_1 \rightarrow_{\lambda} x_0$ in X^i for some $i \in [2]$
 - ξ_1 is an element of $\int |p_{ij}|$, $x_1 = \mathbf{src}_{|p_{ij}|}(\xi_1)$, and $x_0 = \mathbf{tgt}_{|p_{ij}|}(\xi_1)$.
- since $E_{T^n}[\lambda_n]$ is connected, a λ_n -path is either of the following:
 - a path contained in X^i for some $i \in [2]$, which is composed in each T -categories.
 - a path of the form

$$(4.34) \quad x_n \xrightarrow[\lambda_n]{\xi_{n-1}} \cdots \xrightarrow{} x_{m+1} \xrightarrow[\lambda_{m+1}]{\bar{\xi}_m} y_m \xrightarrow[\lambda_m]{} \cdots \xrightarrow[\lambda_1]{\xi_0} y_0$$

where, $\bar{\xi}_m$ is in $|p_{ik}|$, x_j and ξ_j are in X^i for each $j > m$, and y_{l+1} and ξ_l are in X^k for each $l < m$, for some $i < k \in [2]$. Such a path is composed through the composition of X^i and X^k and left and right actions of p_{ik} .

– a path of the form

$$(4.35) \quad x_n \xrightarrow{\xi_{n-1}}_{\lambda_n} \cdots \longrightarrow x_{m+1} \xrightarrow{\bar{\xi}_m}_{\lambda_{m+1}} y_m \xrightarrow{\lambda_m} \cdots \xrightarrow{\xi_{s+1}}_{\lambda_{s+1}} y_{s+1} \xrightarrow{\bar{\xi}_s}_{\lambda_{s+1}} z_s \xrightarrow{\lambda_s} \cdots \xrightarrow{\xi_0}_{\lambda_1} z_0$$

where, x_j and ξ_j are in X^0 for each $j > m$, $\bar{\xi}_m$ is in $|p_{01}|$, y_{l+1} and ξ_l are in X^1 for each $s < l < m$, $\bar{\xi}_s$ is in $|p_{12}|$, and y_{t+1} and ξ_t are in X^1 for each $t < s$. Such a path composed to a path of the form

$$(4.36) \quad \mu^{n-m} \mu^m(x_n) \xrightarrow{\bar{\xi}'}_{\mu^{n-m} \mu^m \lambda_n} \mu^m(y_m) \xrightarrow{\bar{\xi}''}_{\mu^m \lambda_m} z_0$$

by applying compositions of X^i ($i \in [2]$) and actions of p_{01} and p_{12} , and then composed to an arrow $\mu^n(x_n) \xrightarrow{\mu^n \lambda_n} z_0$ contained in $|p_{02}|$ by applying α .

In the same way, one can check that a functor $\Gamma p^{-1} \longrightarrow \nabla_T[n]$ corresponds to $n(n+1)/2$ T -profunctors p_{ij} ($i < j \in [n]$) equipped with horizontal cells

$$(4.37) \quad \begin{array}{ccc} \xrightarrow{p_{ij}} & p_{jk} & \xrightarrow{\phantom{p_{ij}}} \\ \parallel & \alpha_{ijk} & \parallel \\ \xrightarrow{p_{ik}} & & \end{array}$$

for each $i < j < k \in [n]$ which are coherent in obvious way. The pullback along the unique map of the form $\nabla_T(![n]) : \nabla_T[n] \longrightarrow \nabla_T[0]$ sends a T -category X to $n(n+1)/2$ copies of the identity $\text{Id}_X : X \rightrightarrows X$ equipped with the canonical cells induced from the 2-ary composition of X .

Example 7.7 of [CS10] shows that $\mathbb{P}\text{rof}(T)$ is in fact a *virtual equipment*; i.e., for each horizontal and vertical arrows p , f and g in the diagram below, there exists a horizontal arrow $q(f, g) : X \rightrightarrows Y$, which is called the *restriction of q along f and g* , equipped with an *cartesian cell* (see (2.8))

$$(4.38) \quad \begin{array}{ccc} X & \xrightarrow{q(f,g)} & Y \\ f \downarrow & \text{cart} & \downarrow g \\ A & \xrightarrow{q} & B \end{array}$$

Proposition 4.18. *Any path of horizontal arrows in $\mathbb{P}\text{rof}(T)$ is a source of an weakly opcartesian cell. Moreover, the composite of cells below is weakly opcartesian if α and β are so:*

$$(4.39) \quad \begin{array}{ccc} Y & \xrightarrow{p} & X^0 \xrightarrow{p^0} \cdots \xrightarrow{p^{n-1}} & X^n \\ \parallel & \parallel & \parallel & \parallel \\ Y & \xrightarrow{} & \xrightarrow{} \cdots \xrightarrow{\phantom{p^{n-1}}} & \xrightarrow{} \\ \parallel & & \parallel & \parallel \\ & & \xrightarrow{\beta} & \\ & & (p^{n-1} \circ \cdots \circ p^0) \circ p & \end{array}$$

Proof. (See (2.9) and proceeding discussions.)

For each path of T -profunctors p^i ($i = 0, \dots, n-1$), define a span $p^{n-1} \circ \cdots \circ p^1 \circ p^0$ as the coequalizer of the 2^{n-1} arrows explained in (4.28), interpreted as parallel morphisms in $\mathcal{H}(\widehat{\mathbf{C}})(TX_0^0, X_0^n)$ of the form

$$(4.40) \quad p^{n-1} \vee \cdots \vee \text{Id}_{X^2} \vee p^1 \vee \text{Id}_{X^1} \vee p^0 \xrightarrow{\quad \quad \quad \text{; } 2^{n-1} \quad \quad \quad} p^{n-1} \vee \cdots \vee p^1 \vee p^0$$

Since pullbacks in $\widehat{\mathbf{C}}$ preserves coequalizers, for each spans $p : TA_0 \rightrightarrows X_0^0$ and $q : TX_0^n \rightrightarrows B_0$, the above coequalizer is preserved by $q \vee - \vee p$. Therefore, by taking $p := \text{Id}_{X^0}$ and $q := \text{Id}_{X^n}$, the left and right actions of p^{n-1} and p^0 induces those for $p^{n-1} \circ \cdots \circ p^1 \circ p^0$ and it becomes a T -profunctor. It is straightforward to check the canonical cell defining $p^{n-1} \circ \cdots \circ p^1 \circ p^0$ is weakly opcartesian.

Since T preserves pullbacks and pullbacks in $\widehat{\mathbf{C}}$ preserves coequalizers, the universality of the coequalizer defining $p^{n-1} \circ \cdots \circ p^1 \circ p^0 \circ p$ is the same as that defining $(p^{n-1} \circ \cdots \circ p^1 \circ p^0) \circ p$, which means that the cell (4.39) is weakly opcartesian. \square

Remark 4.19. $\mathbb{P}\text{rof}(T)$ is not a double category in general, which means some composite of weakly opcartesian cells may not be weakly opcartesian. This is because T does not have to preserve coequalizers defining those weakly opcartesian cells. \blacksquare

For any two 2-categories \mathcal{A} and \mathcal{B} , we write $\mathbb{2}[\mathcal{A}, \mathcal{B}]$ for the 2-category of 2-functors, 2-natural transformations, and modifications, and $\mathbb{P}[\mathcal{A}, \mathcal{B}]$ for the 2-category of pseudo functors, pseudo natural transformations, and modifications. If we take $\mathcal{B} := \mathbf{Cat}$ and \mathcal{A} to be small, the canonical inclusion $\mathbb{2}[\mathcal{A}, \mathbf{Cat}] \hookrightarrow \mathbb{P}[\mathcal{A}, \mathbf{Cat}]$ is the right adjoint part of a 2-adjoint whose unit is an equivalence; see 4.2 of [Pow89] and [Lac02].

We write $\Delta_2(T)$ for the subcategory of $\Delta(T)$ obtained by taking the pullback of $\Delta(T) \rightarrow \Delta$ along $\Delta_2 \hookrightarrow \Delta$, where Δ_2 is the full subcategory of Δ which consists of simplices of dimension lower than or equal to 2.

Definition 4.20. We write $\mathcal{SCat}(T)$, $\mathcal{SCat}_2(T)$, $\mathcal{PSCat}(T)$, and $\mathcal{PSCat}_2(T)$ for $\mathbb{2}[\Delta(T)^{\text{op}}, \mathbf{Cat}]$, $\mathbb{2}[\Delta_2(T), \mathbf{Cat}]$, $\mathbb{P}[\Delta(T)^{\text{op}}, \mathbf{Cat}]$, and $\mathbb{P}[\Delta_2(T)^{\text{op}}, \mathbf{Cat}]$ respectively. Objects in $\mathcal{SCat}(T)$ ($\mathcal{SCat}_2(T)$) are called (**2-truncated**) *T-simplicial categories*, while those in $\mathcal{PSCat}(T)$ ($\mathcal{PSCat}_2(T)$) are (**2-truncated**) *pseudo T-simplicial categories*. We omit the preposition “T-” when $T = \text{id}_1$. ■

The classical Grothendieck construction shows that a pseudo *T*-simplicial category can be seen as a fibration over $\Delta(T)$.

The discussion about *T*-profunctors above suggests that profunctors may be treated in an ordinary pseudo simplicial category, i.e. a fibration over Δ .

Note that since $\mathbf{Cat}(T)$ is finitely complete, its codomain functor $\mathbf{Cat}(T)^{[1]} \rightarrow \mathbf{Cat}(T)$ is a fibration.

Definition 4.21. The large *pseudo simplicial category of T-categories*, $\overline{\mathbf{Prof}}(T)$, is the pullback of the codomain fibration $\mathbf{Cat}(T)^{[1]} \rightarrow \mathbf{Cat}(T)$ along the functor $\Delta \hookrightarrow \mathbf{Cat} \xrightarrow{\nabla_T} \mathbf{Cat}(T)$. The large **2-truncated pseudo simplicial category of T-categories**, $\overline{\mathbf{Prof}}_2(T)$, is defined as the restriction of $\overline{\mathbf{Prof}}(T)$ to $\Delta_2(T)$. ■

The reflection $\mathbf{SSET} \rightarrow \mathbf{CAT}$ preserves finite products, see for example Lemma 3.3.13 of [Cis19], hence the 2-category of locally large large 2-categories, $\mathbf{CAT-CAT}$, is a full sub 2-category of the 2-category of large simplicially enriched categories, $\mathbf{SSET-CAT}$. A simplicially enriched category is precisely a simplicial category whose structure maps are identity-on-object functors, and one can check that 2-functors and 2-natural transformations are precisely 1-cells and 2-cells in $\mathcal{SCAT} := \mathbb{2}[\Delta^{\text{op}}, \mathbf{CAT}]$, i.e. $\mathbf{SSET-CAT}$ is a full sub 2-category of \mathcal{SCAT} . We write $\mathcal{N} : \mathbf{CAT-CAT} \hookrightarrow \mathcal{SCAT}$ for the composite of those embeddings.

Since a vertically composable *n*-tuple of natural transformations $(\beta_1, \dots, \beta_n)$ in $\mathbf{Cat}(T)$ is precisely a map in $\mathbf{Cat}(T)/\nabla_T[n]$ between identities, $\mathcal{N}(\mathbf{Cat}(T))_n$ is given by the (bijective-on-objects, fully faithful)-factorization of the functor $\mathbf{Cat}(T)/\nabla_T[0] \rightarrow \mathbf{Cat}(T)/\nabla_T[n]$ induced by pullback along the unique map $\nabla_T[n] \rightarrow \nabla_T[0]$. This means that the 2-category structure of *T*-categories is induced from $\overline{\mathbf{Prof}}(T)$.

Let us denote by *T-Alg* the Eilenberg-Moore category of *T*. A *T-algebra* is an element of *T-Alg*. Given a *T*-algebra *X*, one obtains a *T*-category X^* as follows:

- $(X^*)_0$ is the underlying object $|X| \in \widehat{\mathbf{C}}$.
- $(X^*)_1$ is $T|X|$. $\text{src} := \text{id}_{T|X|} : T|X| \rightarrow T|X|$, and $\text{tgt} := h_X : T|X| \rightarrow |X|$ is the structure map of *X*.
- Now that we obtain a *T*-graph, a path $x_n \xrightarrow{x_n} x_0$ makes sense and corresponds to a sequence $(\lambda_n; x_n, \dots, x_0)$ such that $T^i h_X(x_{i+1}) = x_i$ and $T^n \text{ty}_X(x_n) = \lambda_n$ for each $i < n$. One can easily check that $\text{comp}_n(\lambda_n; x_n, \dots, x_0) := (\mu^n(\lambda_n); \mu^n(x_n), x_0)$ is well defined; i.e. the right hand side is $\mu^n(\lambda_n)$ -arrow, which is exactly the same as the condition for h_X to be an algebra.

This construction induces a functor $(-)^* : T\text{-Alg} \rightarrow \mathbf{Cat}(T)$.

Therefore, we obtain a 2-functor $\overline{\mathfrak{M}}_T^s : \mathcal{V}\mathcal{D}\mathcal{b}\mathcal{l}\mathcal{C}\mathcal{a}\mathcal{t}(T) \rightarrow \mathcal{SCat}_2(T)$ as the \mathbf{Cat} -enriched left Kan extension of the 2-Yoneda embedding $\mathfrak{y} : \Delta_2(T) \rightarrow \mathcal{SCat}_2(T)$ along $[-]_T^*$, which is a 2-functor from locally discrete 2-category $\Delta_2(T)$, i.e., $\overline{\mathfrak{M}}_T^s(X)$ is $\mathcal{V}\mathcal{D}\mathcal{b}\mathcal{l}\mathcal{C}\mathcal{a}\mathcal{t}(T)([-]_T^*, X)$.

We write $\overline{\mathfrak{M}}_T$ for the composite $\mathcal{V}\mathcal{D}\mathcal{b}\mathcal{l}\mathcal{C}\mathcal{a}\mathcal{t}(T) \xrightarrow{\overline{\mathfrak{M}}_T^s} \mathcal{SCat}_2(T) \rightarrow \mathcal{PSCat}_2(T)$.

Remark 4.22. If the 2-category $\mathcal{V}\mathcal{D}\mathcal{b}\mathcal{l}\mathcal{C}\mathcal{a}\mathcal{t}(T)$ is cocomplete as a 2-category, those 2-functors are the right parts of 2-adjoints. This follows from Theorem 4.51 of [Kel82]. ■

When $T = \text{id}_1$, for each $[n] \in \Delta$, $[n]^*$ is a virtual double category defined as follows

- the set of objects is the underlying set of $[n]$.
- there is no non-trivial vertical arrows i.e. the vertical category is discrete.
- for each pair $i \leq j$ in $[n]$, there is a unique horizontal arrow $i \xrightarrow{ij} j$.
- any possible squares are filled in with a unique cell.

A map $[0]^* \rightarrow \mathbb{X}$ in $\mathbf{V}\mathcal{D}\mathcal{b}\mathcal{l}\mathcal{C}\mathcal{a}\mathcal{t}$ is precisely a *monoid* in \mathbb{X} (in the sense of [CS10]) and a map $[1]^* \rightarrow \mathbb{X}$ is what is called a *module* between monoids.

Theorem 4.23. $\overline{\mathbf{Prof}}(T)$ is equivalent to a free object of $\overline{\mathbf{Prof}}_2(T)$ with respect to the 2-functor $\overline{\mathfrak{M}}_{\text{id}_1} : \mathcal{V}\mathcal{D}\mathcal{b}\mathcal{l}\mathcal{C}\mathcal{a}\mathcal{t} \rightarrow \mathcal{PSCAT}_2$.

Proof. The discussion preceding Definition 4.17 suggests $\overline{\mathbf{Prof}}(T)$ can be seen as a strict simplicial category $\overline{\mathbf{Prof}}(T)'$ up-to-equivalence defined as follows: $\overline{\mathbf{Prof}}(T)'_n$ is the full subcategory of $\overline{\mathfrak{M}}_{\text{id}_1}(\overline{\mathbf{Prof}}(T))_n =$

$\mathcal{VDblCat}([n]^*, \mathbb{P}\text{rof}(T))$ consisting of *strictly normal functors*; i.e., maps $[n]^* \rightarrow \mathbb{P}\text{rof}(T)$ in $\mathbf{VDblCAT}$ which send $i \xrightarrow{ii} i$ to identity profunctors, and cells

$$(4.41) \quad \begin{array}{ccc} \xrightarrow{ii} & \xrightarrow{ij} & \\ \parallel & \cdot & \parallel \\ \xrightarrow{ij} & & \end{array}, \quad \begin{array}{ccc} \xrightarrow{ij} & \xrightarrow{jj} & \\ \parallel & \cdot & \parallel \\ \xrightarrow{ij} & & \end{array}$$

to cells in $\mathbb{P}\text{rof}(T)$ induced from left and right actions:

$$(4.42) \quad \begin{array}{ccc} X_i \xrightarrow{\text{Id}_{X_i}} X_i & \xrightarrow{p_{ij}} & X_j \\ \parallel & \lambda^{p_{ij}} & \parallel \\ X_i & \xrightarrow{p_{ij}} & X_j \end{array}, \quad \begin{array}{ccc} X_i & \xrightarrow{p_{ij}} & X_j \xrightarrow{\text{Id}_{X_j}} X_j \\ \parallel & \rho^{p_{ij}} & \parallel \\ X_i & \xrightarrow{p_{ij}} & X_j \end{array}$$

This induces a strict simplicial category since images of maps in Δ sends each $i \xrightarrow{ii} i$ to $j \xrightarrow{jj} j$ for some j .

Let $\overline{\mathbb{P}\text{rof}}_2(T)'$ be the restriction of $\overline{\mathbb{P}\text{rof}}(T)'$ to Δ_2 . Now we show that $\mathbb{P}\text{rof}(T)$ is a free object of $\overline{\mathbb{P}\text{rof}}_2(T)'$ with respect to $\overline{\mathfrak{M}}_{\text{id}_1}^s$. Let \mathbb{X} be a virtual double category. For each map $\mathbb{P}\text{rof}(T) \rightarrow \mathbb{X}$ in $\mathcal{VDblCat}$, the post composition gives rise to a map $\overline{\mathbb{P}\text{rof}}_2(T)' \rightarrow \overline{\mathfrak{M}}_{\text{id}_1}^s(\mathbb{X})$, and this induces a functor

$$\mathcal{VDblCat}(\mathbb{P}\text{rof}(T), \mathbb{X}) \rightarrow \mathcal{SCAT}(\overline{\mathbb{P}\text{rof}}(T)', \overline{\mathfrak{M}}_{\text{id}_1}^s(\mathbb{X}))$$

which is 2-natural in \mathbb{X} . In fact, this functor is faithful since any 2-cell in $\mathcal{VDblCat}$ is completely determined by its whiskerings with strictly normal functors $[1]^* \rightarrow \mathbb{P}\text{rof}(T)$. On the other hand, for each map $F : \overline{\mathbb{P}\text{rof}}_2(T)' \rightarrow \overline{\mathfrak{M}}_{\text{id}_1}^s(\mathbb{X})$, we can construct $\bar{F} : \mathbb{P}\text{rof}(T) \rightarrow \mathbb{X}$ as follows:

- \bar{F} sends a T -category X to $F_0(\ulcorner X \urcorner)(0)$, where $\ulcorner X \urcorner$ is the strictly normal functor $[0]^* \rightarrow \mathbb{P}\text{rof}(T)$ representing X , and $F_0(\ulcorner X \urcorner)$ is a map $[0]^* \rightarrow \mathbb{X}$.
- \bar{F} sends a T -profunctor $X \xrightarrow{p} Y$ to $F_1(\ulcorner p \urcorner)(01)$, where $\ulcorner p \urcorner$ is the strictly normal functor $[1]^* \rightarrow \mathbb{P}\text{rof}(T)$ representing p , and $F_1(\ulcorner p \urcorner)$ is a map $[1]^* \rightarrow \mathbb{X}$.
- \bar{F} sends a T -functor $X \xrightarrow{f} Y$ to $F_0(\ulcorner f \urcorner)_0$, where $\ulcorner f \urcorner$ is the 2-cell $\ulcorner f \urcorner : \ulcorner X \urcorner \Rightarrow \ulcorner Y \urcorner : [0]^* \rightarrow \mathbb{P}\text{rof}(T)$ representing f .
- In the same way, \bar{F} sends a unary cell

$$(4.43) \quad \begin{array}{ccc} X & \xrightarrow{p} & Y \\ f \downarrow & \alpha & \downarrow g \\ A & \xrightarrow{q} & B \end{array}$$

to $F_1(\ulcorner \alpha \urcorner)_{01}$, where $\ulcorner \alpha \urcorner$ is the 2-cell $\ulcorner \alpha \urcorner : \ulcorner p \urcorner \Rightarrow \ulcorner q \urcorner : [1]^* \rightarrow \mathbb{P}\text{rof}(T)$ representing α .

- Let α be an arbitrary cell in $\mathbb{P}\text{rof}(T)$ of the form

$$(4.44) \quad \begin{array}{ccccccc} X^0 & \xrightarrow{p^0} & X^1 & \xrightarrow{p^1} & \dots & \xrightarrow{p^{n-1}} & X^n \\ \downarrow f & & & \alpha & & & \downarrow g \\ Y^0 & \xrightarrow{q} & & & & & Y^1 \end{array}$$

[Proposition 4.18](#) shows α uniquely factors through the weakly opcartesian cell written as $p^{n-1} \circ \dots \circ p^0$. We write $\bar{\alpha}$ for the result of this factorization; i.e., α factors as follows:

$$(4.45) \quad \begin{array}{ccccccc} X^0 & \xrightarrow{p^0} & X^1 & \xrightarrow{p^1} & \dots & \xrightarrow{p^{n-1}} & X^n \\ \downarrow f & & & \alpha & & & \downarrow g \\ Y^0 & \xrightarrow{q} & & & & & Y^1 \end{array} = \begin{array}{ccccccc} X^0 & \xrightarrow{p^0} & X^1 & \xrightarrow{p^1} & \dots & \xrightarrow{p^{n-1}} & X^n \\ \parallel & & & \tilde{p} & & & \parallel \\ X^0 & \xrightarrow{p^{n-1} \circ \dots \circ p^0} & & & & & X^n \\ \downarrow f & & & \bar{\alpha} & & & \downarrow g \\ Y^0 & \xrightarrow{q} & & & & & Y^1 \end{array}$$

Since we have already defined where α is sent, we define the cell $\bar{F}(\tilde{p})$ in \mathbb{X} when $n \neq 1$, and $\bar{F}(\alpha)$ is defined as the composite in \mathbb{X} .

- If $n = 0$, then the identity T -profuctor is the 0-ary composition i.e., $p^{n-1} \circ \dots \circ p^0 = \text{Id}_{X^0}$. \tilde{p} is sent to what the following cell in $[0]^*$ is sent to by $F_0(\ulcorner X^{0\urcorner})$:

$$(4.46) \quad \begin{array}{ccc} & 0 & \\ \swarrow & & \searrow \\ & \cdot & \\ 0 & \xrightarrow{00} & 0 \end{array}$$

- If $n > 1$, again by Proposition 4.18, $p^{n-1} \circ \dots \circ p^0$ is isomorphic to $(\dots (p^{n-1} \circ p^{n-2}) \circ p^{n-3}) \dots \circ p^0$, and \tilde{p} is the composite of $n - 1$ opcartesian cells whose sources are of length 2. Therefore what to define is $n = 2$ case. In this case, $\bar{F}(\tilde{p})$ is defined as the image of the unique cell filling the square below in $[2]^*$ by $F_2(\ulcorner \tilde{p} \urcorner) : [2]^* \rightarrow \mathbb{X}$, where $\ulcorner \tilde{p} \urcorner$ is the strictly normal functor $[2]^* \rightarrow \mathbb{P}\text{rof}(T)$ representing \tilde{p} .

$$(4.47) \quad \begin{array}{ccccc} 0 & \xrightarrow{01} & 1 & \xrightarrow{12} & 2 \\ \parallel & & & & \parallel \\ 0 & \xrightarrow{02} & & & 2 \end{array}$$

It is straightforward to check this \bar{F} is a morphism in $\mathcal{V}\text{Db}l\text{Cat}(T)$ and prove surjectiveness of the functor $\mathcal{V}\text{Db}l\text{Cat}(\mathbb{P}\text{rof}(T), \mathbb{X}) \rightarrow \mathcal{S}\text{Cat}(\overline{\mathbb{P}\text{rof}(T)}, \overline{\mathfrak{M}}_{\text{id}_1}^s(\mathbb{X}))$.

Moreover, the construction of \bar{F} suggests that this functor is full since naturality with α above of a transformation between morphisms \bar{F} and \bar{G} in $\mathcal{V}\text{Db}l\text{Cat}$ is determined by

- naturality with $\bar{\alpha}$, which follows from naturality of $F \Rightarrow G$ in $\mathcal{S}\text{Cat}_2$ on $[1]$.
- naturality with \tilde{p} , which follows from naturality of $F \Rightarrow G$ in $\mathcal{S}\text{Cat}_2$ on $[2]$.

Thus this functor is an isomorphism, which means freeness of $\mathbb{P}\text{rof}(T)$. \square

Finally, we suggest a way to define a $n + 2$ -dimensional structure of virtual $n + 1$ -tuple categories.

In the Section 3 of [CS10], for each virtual double category \mathbb{X} , the virtual double category of monoids in \mathbb{X} , $\text{Mod}(\mathbb{X})$, is defined and it is proved that Mod extends to an endo-functor on $\mathbf{V}\text{Db}l\text{Cat}$. Moreover, in the Section 5, it is proved that this endo-functor is induced from the pseudo-adjunction between the 2-category of virtual double categories and the category of *unital* virtual double categories.

On the other hand, there are elements written as $0 \in \int S_{\text{fc}} \subset \mathbf{G}_1(\text{fc})$ and $[1] \in \mathbf{G}_1 \subset \mathbf{G}_1(\text{fc})$, and the virtual double category $[[1]]_{\text{fc}}$ is the smallest virtual double category which contains a horizontal arrow $0 \rightarrow 1$, while $[0]_{\text{fc}}$ is the smallest virtual double category which contains a vertical arrow $0 \rightarrow 1$. In particular $\text{Mod}([0]_{\text{fc}})$ is the same as $\nabla_{\text{fc}}([1])$, hence a fc -functor $\ulcorner \mathbb{p} \urcorner : \ulcorner \mathbb{p} \urcorner \rightarrow \text{Mod}([0]_{\text{fc}})$ is the same as a fc -profuctor $\mathbb{X} \rightarrow \mathbb{Y}$, which can be seen as a **vertical profuctor**, since if the restrictions of \mathbb{p} on 0 and 1 are \mathbb{X} and \mathbb{Y} respectively, then it can be seen as a virtual double category which consists of each copy of \mathbb{X} and \mathbb{Y} , additional vertical arrows from elements in \mathbb{X} to those in \mathbb{Y} , and additional cells containing those vertical arrows. In the same way, we can define a **horizontal profuctor** $\ulcorner \mathbb{u} \urcorner : \mathbb{X} \rightarrow \mathbb{Y}$: a virtual double category which consists of each copy of \mathbb{X} and \mathbb{Y} , additional horizontal arrows from elements in \mathbb{X} to those in \mathbb{Y} , and additional cells containing those horizontal arrows. One can easily check that a horizontal profuctor can be seen as a fc -functor $\ulcorner \mathbb{u} \urcorner \rightarrow \text{Mod}([1]_{\text{fc}})$ in the same way as vertical profunctors.

For each $n \in \int S_{\text{fc}} \subset \mathbf{G}_1(\text{fc})$, the virtual double category $[n]_{\text{fc}}$ classifies n -cells in arbitrary virtual double category, i.e., a map $[n]_{\text{fc}} \rightarrow \mathbb{X}$ is the same as a n -ary cell in \mathbb{X} . Therefore, $\text{Mod}([n]_{\text{fc}})$ has the same objects as $[n]_{\text{fc}}$, but each object has its unit. Hence a fc -functor $\ulcorner \alpha \urcorner \rightarrow \text{Mod}([n]_{\text{fc}})$ is a virtual double category consisting of $n + 1$ horizontal profunctors and 2 vertical profunctors which makes the frame of the square below and additional cells connecting those profunctors.

$$(4.48) \quad \begin{array}{ccccccc} & \dashrightarrow & \dashrightarrow & \dots & \dashrightarrow & \dashrightarrow & \\ \downarrow & & & \alpha & & & \downarrow \\ & \dashrightarrow & \dashrightarrow & \dashrightarrow & \dashrightarrow & \dashrightarrow & \end{array}$$

Those define a fc -graph of virtual double categories and vertical and horizontal profunctors.

By considering the analogy of Theorem 4.23, we obtain a definition of the virtual triple category of virtual double categories, as follows:

Definition 4.24. The large fc -pseudo simplicial category of virtual double categories, $\overline{2\mathfrak{P}\text{rof}}$, is defined as the pullback of the codomain fibration $\mathbf{V}\text{Db}l\text{Cat}^{[1]} \rightarrow \mathbf{V}\text{Db}l\text{Cat}$ along the composite

$$(4.49) \quad \Delta(\text{fc}) \xrightarrow{[-]_{\text{fc}}} \mathbf{V}\text{Db}l\text{Cat} \xrightarrow{\text{Mod}} \mathbf{V}\text{Db}l\text{Cat}$$

The large **2-truncated fc -pseudo simplicial category of virtual double categories**, $\overline{2\mathfrak{P}\text{rof}}_2$, is defined as the restriction of $\overline{2\mathfrak{P}\text{rof}}$ to $\Delta_2(\text{fc})$. The **virtual triple category of virtual double categories**, $2\mathfrak{P}\text{rof}$, is the free object of $\overline{2\mathfrak{P}\text{rof}}_2$ with respect to the 2-functor $\overline{\mathfrak{M}}_{\text{fc}} : \mathcal{V}\text{-}3\text{-tp}l\text{CAT} \rightarrow \mathcal{P}\text{SCAT}_2(\text{fc})$. \blacksquare

This clearly indicates a way to define the virtual $n + 2$ -tuple category of virtual $n + 1$ -tuple category: Firstly we define endo-functors $M_n : \mathbf{V}\text{-}n + 1\text{-}\mathbf{tpICat} \rightarrow \mathbf{V}\text{-}n + 1\text{-}\mathbf{tpICat}$ which makes virtual $n + 1$ -tuple categories “*unital*” in some way. Then we can define structures of virtual $n + 1$ -tuple categories as follows:

- The large \mathbf{fc}^n -pseudo simplicial category, $\overline{n + 1\mathfrak{Prof}}$, is the pullback of the codomain fibration on $\mathbf{V}\text{-}n + 1\text{-}\mathbf{tpICat}$ along

$$(4.50) \quad \Delta(\mathbf{fc}^n) \xrightarrow{[-]_{\mathbf{fc}^n}} \mathbf{V}\text{-}n + 1\text{-}\mathbf{tpICat} \xrightarrow{M_n} \mathbf{V}\text{-}n + 1\text{-}\mathbf{tpICat}$$

and let $\overline{n + 1\mathfrak{Prof}_2}$ be its 2-truncated version.

- The large virtual $n + 2$ -tuple category, $n + 1\mathfrak{Prof}$, is the free object of $\overline{n + 1\mathfrak{Prof}_2}$ with respect to the 2-functor $\mathfrak{M}_{\mathbf{fc}^n} : \mathcal{V}\text{-}n + 2\text{-}\mathit{tplCAT} \rightarrow \mathcal{PSCAT}_2(\mathbf{fc}^n)$.

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Email address: hoshinok@kurims.kyoto-u.ac.jp