

The Grothendieck Construction for Double Categories

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Double Categories, the concise and the expanded definition

- A **double category** is an internal category in **Cat**,

$$C_1 \times_{C_0} C_1 \xrightarrow{\circ} C_1 \begin{array}{c} \xrightarrow{s} \\ \xleftarrow[t]{} \end{array} C_0 .$$

- It has

- objects (objects of C_0),
- vertical arrows (arrows of C_0), denoted $A \xrightarrow{u} A'$,
- horizontal arrows (objects of C_1), denoted $A \xrightarrow{f} B$,
- double cells (arrows of C_1), denoted

$$\begin{array}{ccc} A & \xrightarrow{f} & B \\ u \downarrow & \alpha & \downarrow v \\ A' & \xrightarrow[f']{} & B' \end{array}$$

Compositions and identities

- Since C_0 is a category, we have vertical compositions

$u' \bullet u : A \xrightarrow{u} A' \xrightarrow{u'} A''$ and identities $A \xrightarrow{id_A} A$.

- Since C_1 is a category too, we have vertical compositions (pastings)

$$\alpha' \bullet \alpha : \begin{array}{ccc} A & \xrightarrow{f} & B \\ u \bullet & \downarrow & \downarrow v \\ A' & \xrightarrow{f'} & B' \\ u' \bullet & \downarrow & \downarrow v' \\ A'' & \xrightarrow{f''} & B'' \end{array} \quad \text{and identities} \quad \begin{array}{ccc} A & \xrightarrow{f} & B \\ id_A \bullet & \downarrow & \downarrow id_B \\ A & \xrightarrow{f} & B \end{array}$$

Compositions and identities

Since $C_1 \times_{C_0} C_1 \xrightarrow{\circ} C_1 \xleftarrow{1} C_0$ are functors:

- horizontal arrows form a category too, we have thus $g \circ f : A \xrightarrow{f} B \xrightarrow{g} C$ and identities $A \xrightarrow{1_A} A$.
- double cells can also be pasted horizontally

$$\begin{array}{ccccc} & A & \xrightarrow{f} & B & \xrightarrow{g} \\ \beta \circ \alpha : & u \bullet & \downarrow \alpha & v \bullet & \downarrow \beta \\ & A' & \xrightarrow{f'} & B' & \xrightarrow{g'} \\ & & & & \end{array} \quad \text{with identities} \quad \begin{array}{ccc} A & \xrightarrow{1_A} & A \\ u \bullet & \downarrow & 1_u & \downarrow u \\ A' & \xrightarrow{1_{A'}} & A' \end{array}$$

Middle four interchange

Since $C_1 \times_{C_0} C_1 \xrightarrow{\circ} C_1$ is a functor, it commutes with composition.
That means that given two arrows of $C_1 \times_{C_0} C_1$

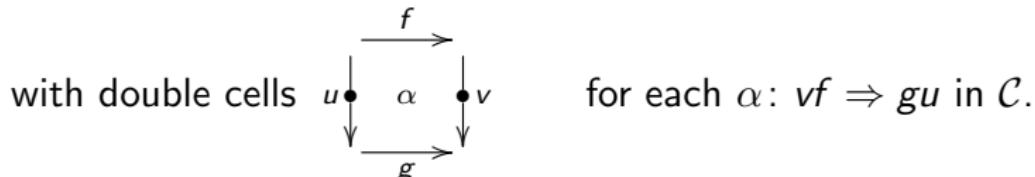
$$\begin{array}{c} (g, f) \\ \downarrow (\beta, \alpha) \\ (g', f') \\ \downarrow (\beta', \alpha') \\ (g'', f'') \end{array} \quad \text{i.e. a configuration}$$

$$\begin{array}{ccccc} A & \xrightarrow{f} & B & \xrightarrow{g} & C \\ u \bullet & \downarrow & v \bullet & \downarrow & w \bullet \\ A' & \xrightarrow{f'} & B' & \xrightarrow{g'} & C' \\ u' \bullet & \downarrow & v' \bullet & \downarrow & w' \bullet \\ A'' & \xrightarrow{f''} & B'' & \xrightarrow{g''} & C'' \end{array}$$

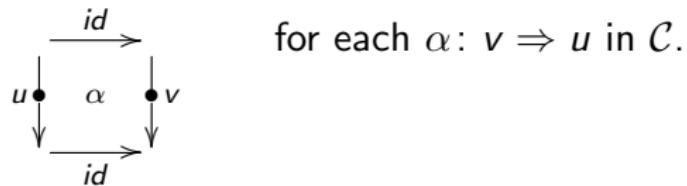
$$\text{we have } (\beta' \circ \alpha') \bullet (\beta \circ \alpha) = (\beta' \bullet \beta) \circ (\alpha' \bullet \alpha)$$

Examples

- ① A 2-category can be defined as a double category in which every horizontal arrow is an identity. It has 2-cells $\alpha : v \Rightarrow u$.
- ② For any 2-category \mathcal{C} , $\mathbb{Q}(\mathcal{C})$ is the double category of quintets in \mathcal{C} ,



- ③ $\mathbb{V}(\mathcal{C})$ is the double category with double cells



- ④ More generally, if Σ is a 1-subcategory of \mathcal{C} , in $\mathbb{Q}^\Sigma(\mathcal{C}) \subseteq \mathbb{Q}(\mathcal{C})$ we require the horizontal arrows to be in Σ . ($\mathbb{Q}^{\{id\}}(\mathcal{C}) = \mathbb{V}(\mathcal{C})$)
- ⑤ The double category $\mathbb{H}(\mathcal{C})$ is defined analogously.

The category **DblCat** - Definition

The category **DblCat** of double categories has:

- **objects:** double categories $\mathbb{C}, \mathbb{D}, \dots$;
- **arrows:** double functors F, G, \dots ;
- **transformations:** these come in two *flavors*:
 - a horizontal transformation $\gamma: F \Rightarrow G$ is given by

$$\begin{array}{ccc} FA & \xrightarrow{\gamma_A} & GA \\ \downarrow Fv \bullet & \gamma_v & \downarrow \bullet Gv \\ FB & \xrightarrow{\gamma_B} & GB \end{array} \quad \text{for each } A \text{ in } \text{dom}(F)$$

functorial in the vertical direction and natural in the horizontal direction.

- **vertical transformations** $v: F \Rightarrow G$ are defined dually;
- **modifications** given by a family of double cells.

The category **DblCat** - Properties

- **DblCat** is not a double category;
- a double category has two types of arrows, and **DblCat** has only one;
- a double category has one type of 2-cell, and **DblCat** has two;
- **DblCat** is enriched in double categories: $\mathbf{DblCat}(\mathbb{C}, \mathbb{D})$ is a double category.

Review of the Grothendieck construction \rightsquigarrow Questions

Whiteboard

- Can we do this for $F: \mathbb{D} \rightarrow \mathbf{DblCat}$?
- What sort of colimit do we get?
- What's the relation to other pre-existing constructions?

Double Index Functors

- We would like to have a double functor $F: \mathbb{D} \rightarrow \mathbf{DblCat}$.
- So we need to build a double category out of \mathbf{DblCat} .
- We first make the 2-category \mathbf{DblCat}_v of double categories, double functors and **vertical** transformations.
- And then apply the quintet construction to get the double category $\mathbb{Q}\mathbf{DblCat}_v$.
- So a **double index functor** is a double functor $F: \mathbb{D} \rightarrow \mathbb{Q}\mathbf{DblCat}_v$.
- We will also call this a **vertical double functor** $\mathbb{D} \rightarrow \mathbf{DblCat}$.
- It looks as if at this point we have lost most of the horizontal data.

Double Transformations

- We regain use of some of the horizontal structure in the definition of **double transformation** between double index functors

$$F, G : \mathbb{D} \rightarrow \mathbb{Q}\mathbf{DblCat}_v.$$

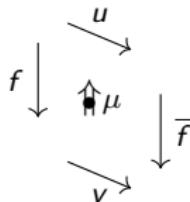
- Analogous to the set-up for bicategories, we will define these transformations in terms of a double category of *cylinders*, $\text{Cyl}_v(\mathbf{DblCat})$.

Whiteboard

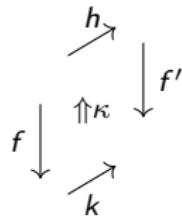
The Double Category of (Vertical) Cylinders

The double category $\text{Cyl}_v(\mathbf{DblCat})$ of vertical cylinders has

- Objects f are arrows of \mathbf{DblCat} , $\downarrow f$.
- Vertical arrows $f \xrightarrow{(u,\mu,v)} \bar{f}$ are given by vertical transformations,



- Horizontal arrows are given by horizontal transformations,



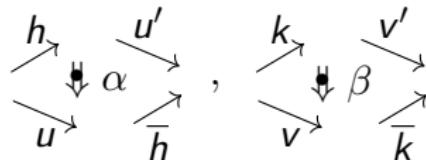
Double Cylinders

$$f \xrightarrow{(h,\kappa,k)} f'$$

A **double cell**, $(u, \mu, v) \downarrow (\alpha, \Sigma, \beta) \downarrow (u', \mu', v')$ is given by two vertical

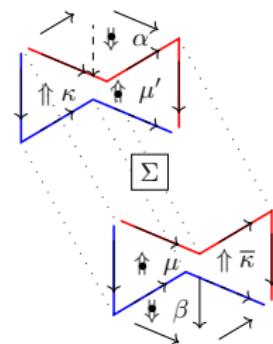
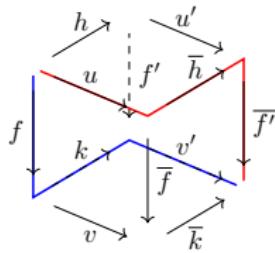
$$\overline{f} \xrightarrow{(\bar{h}, \bar{\kappa}, \bar{k})} \overline{f'}$$

transformations



and a modification Σ ,

$$\begin{array}{c} v'kf \xrightarrow{v'\kappa} v'f'h \\ \Downarrow \beta f \qquad \Downarrow \mu'h \\ \bar{k}vf \quad \Sigma \quad \bar{f}'u'h \\ \Downarrow \bar{k}\mu \qquad \Downarrow \bar{f}'u \\ \bar{k}fu \xrightarrow{\bar{\kappa}u} \bar{f}'\bar{h}u \end{array}$$



Cylinders lead to double lax transformations

- There are vertical double functors $\pi_0, \pi_1: \text{Cyl}_v(\mathbf{DblCat}) \rightarrow \mathbf{DblCat}$, projecting to the top and the bottom of each cylinder respectively;
- A **double lax transformation** $\theta: F \Rightarrow G$ between vertical double functors $F, G: \mathbb{D} \rightarrow \mathbf{DblCat}$ is given by a double functor

$$\theta: \mathbb{D} \rightarrow \text{Cyl}_v(\mathbf{DblCat}), \quad \text{with } \pi_0\theta = F, \pi_1\theta = G.$$

- For each A :

$$\begin{array}{ccc} FA & & \\ \theta_A \downarrow & & \\ GA & & \end{array} .$$

For each u :

$$\begin{array}{ccccc} FA & \xrightarrow{Fu} & FA' & & \\ \theta_A \downarrow & \uparrow \theta_u & \downarrow \theta_{A'} & & \\ GA & & GA' & & \\ & \searrow & \swarrow & & \\ & & Gu & & GA' \end{array}$$

For each f :

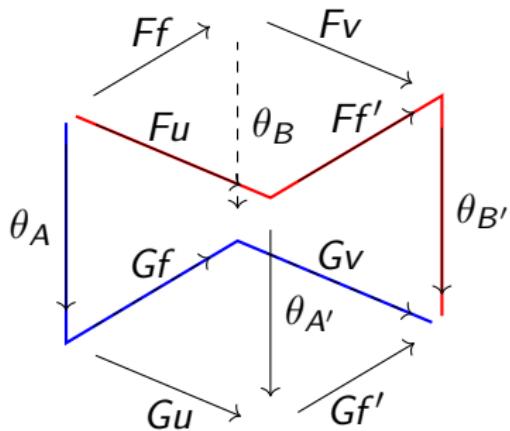
$$\begin{array}{ccccc} & FB & & & \\ & \nearrow Ff & & & \\ FA & & \downarrow \theta_B & & \\ \uparrow \theta_f & & & & \\ GA & & GB & & \\ & \nearrow & \searrow & & \\ & & Gf & & \end{array}, \text{ and...}$$

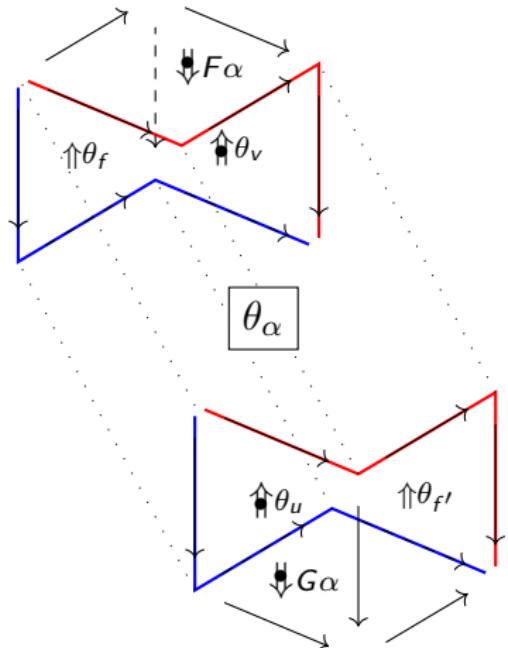
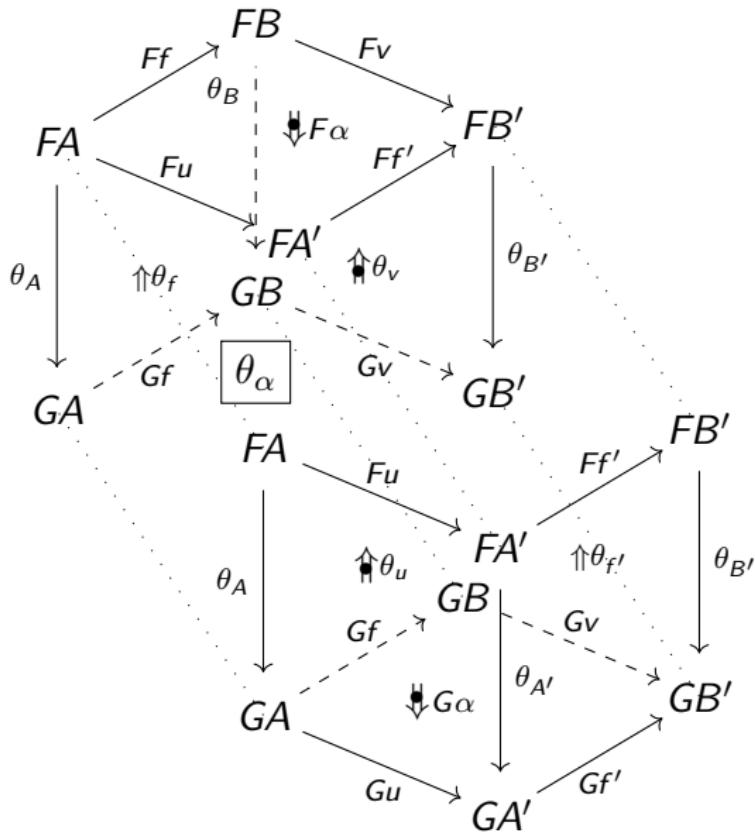
A Double Lax Transformation $\theta: F \Rightarrow G$

For each double cell

$$\begin{array}{ccc} A & \xrightarrow{f} & B \\ u \downarrow & \alpha & \downarrow v \\ A' & \xrightarrow{f'} & B' \end{array}$$

$$\begin{array}{c} GvGf\theta_A \xrightarrow{Gv\theta_f} Gv\theta_B Ff \\ \parallel \qquad \qquad \qquad \parallel \\ G\alpha\theta_A \qquad \qquad \qquad \theta_v Ff \\ \parallel \qquad \qquad \qquad \parallel \\ Gf'Gu\theta_A \qquad \theta_\alpha \qquad \theta_{B'}FvFf \\ \parallel \qquad \qquad \qquad \parallel \\ Gf'\theta_u \qquad \qquad \qquad \theta_{B'}F\alpha \\ \parallel \qquad \qquad \qquad \parallel \\ Gf'\theta_{A'}Fu \xrightarrow{\theta_{f'}Fu} \theta_{B'}Ff'Fu \end{array}$$





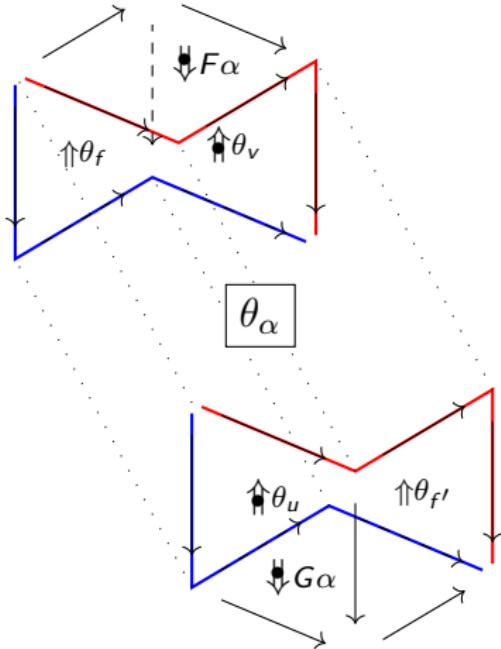
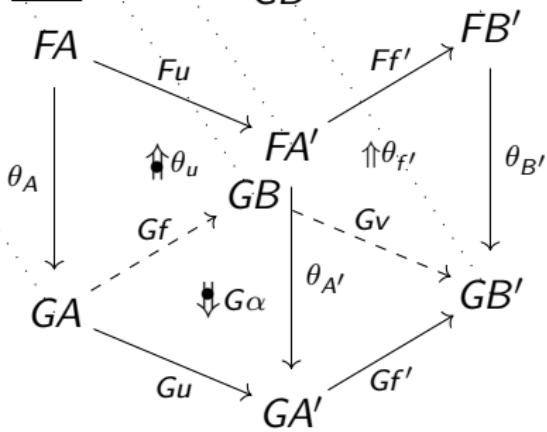
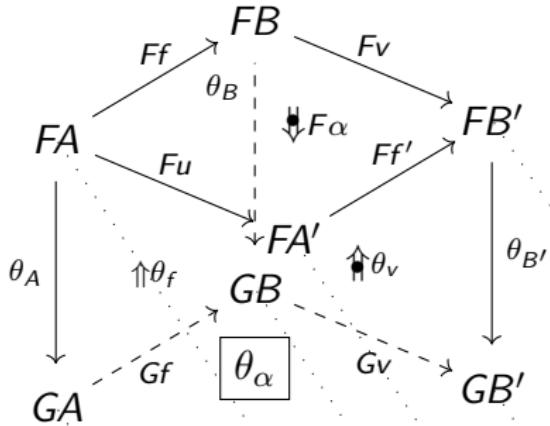
Double Lax Cones

Let $F : \mathbb{D} \rightarrow \mathbf{DblCat}$ be a vertical double functor and $\mathbb{E} \in \mathbf{DblCat}$. A **double lax (co)cone** for F , with vertex \mathbb{E} , is a double lax transformation

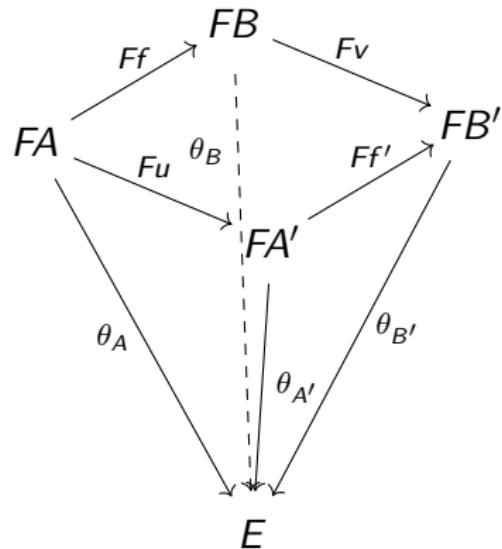
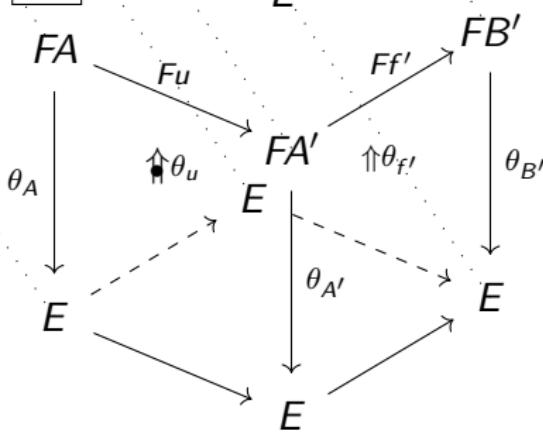
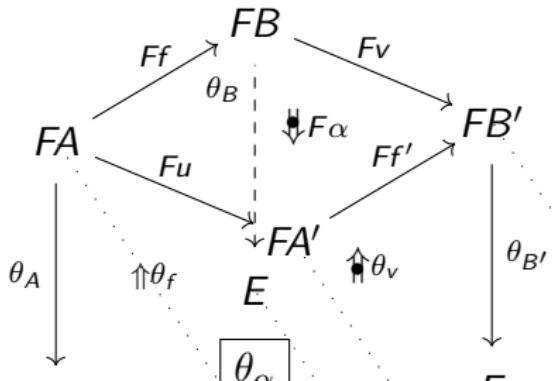
$$F \xrightarrow{\theta} \Delta\mathbb{E}$$

A **double lax colimit** of F is a *universal* double lax cone $F \xrightarrow{\lambda} \Delta\mathbb{L}$.

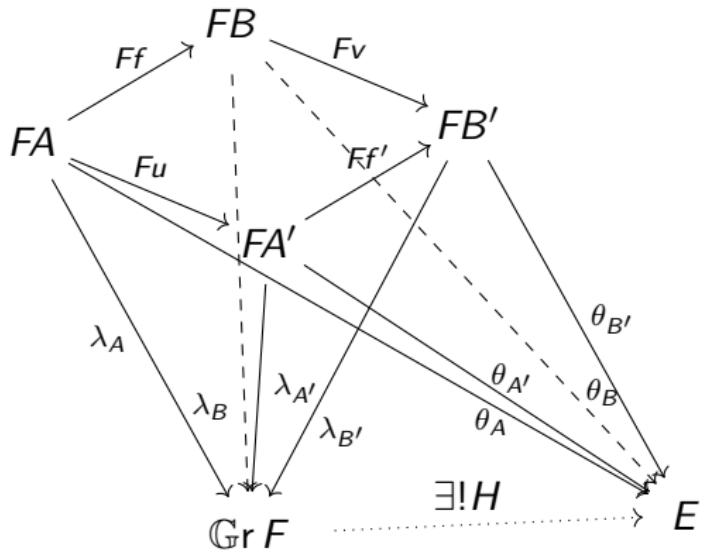
Recall the general double lax natural (6 double cells, one triple cell):



A double lax cone (5 double cells, one triple cell):



The double lax colimit of F :



$$H\lambda_A = \theta_A, \quad H\lambda_u = \theta_u, \quad H\lambda_f = \theta_f, \quad H\lambda_\alpha = \theta_\alpha$$

(for all $A, u, f, \alpha!$)

Examples “It matters how we index”

- If \mathcal{A} is a 2-category, then a 2-functor $\mathcal{A} \xrightarrow{F} \mathbf{DblCat}_v$ can be seen as a vertical double functor $\mathbb{V}\mathcal{A} \xrightarrow{F} \mathbf{DblCat}$ and as $\mathbb{H}\mathcal{A} \xrightarrow{F} \mathbf{DblCat}$.
- For $\mathbb{H}\mathcal{A} \xrightarrow{F} \mathbf{DblCat}$, we see that a lax cone is given by:
 - double functors $FA \xrightarrow{\theta_A} E$,
 - horizontal transformations $\theta_A \xrightarrow{\theta_f} \theta_B Ff$,
 - for each 2-cell $f \xrightarrow{\alpha} f'$ of \mathcal{A} , a modification θ_α ,

$$\begin{array}{ccc} \theta_A & \xrightarrow{\theta_f} & \theta_B Ff \\ id \downarrow & \theta_\alpha & \downarrow \theta_B F\alpha \\ \theta_A & \xrightarrow[\theta_{f'}]{} & \theta_B Ff' \end{array}$$

- For $\mathbb{V}\mathcal{A} \xrightarrow{F} \mathbf{DblCat}$, the 2-cells θ_f are required to be vertical, thus the triple cells θ_α are triple cells of \mathbf{DblCat}_v . Now everything is vertical! This is a *lax triccolimit* in the 3-category \mathbf{DblCat}_v .

The Double Grothendieck Construction: Objects and Arrows

Let \mathbb{D} be a double category, and let $\mathbb{D} \xrightarrow{F} \mathbb{Q}\mathbf{DblCat}_v$ be a double functor. The **double category of elements**, $\mathbb{G}r F = \mathbb{E}l F = \int_{\mathbb{D}} F$, is defined by:

- Objects: (C, X) with C in \mathbb{D} and X in FC ,
- Vertical arrows:

$$(C, X) \xrightarrow{(u, \rho)} (C', X'),$$

where $C \xrightarrow{u} C'$ in \mathbb{D} and $FuX \xrightarrow{\rho} X'$ in FC' .

- Horizontal arrows:

$$(C, X) \xrightarrow{(f, \varphi)} (D, Y),$$

where $C \xrightarrow{f} D$ in \mathbb{D} , and $FfX \xrightarrow{\varphi} Y$ in FD .

The Double Grothendieck Construction: Double Cells

$$(C, X) \xrightarrow{(f, \varphi)} (D, Y)$$

- Double cells: $(u, \rho) \bullet \downarrow \quad (\alpha, \Phi) \quad \bullet (v, \lambda) \downarrow$, where $\alpha : (u \underset{f'}{\sim} v)$ is a double cell in \mathbb{D} and Φ is a double cell in FD' :
- $$(C', X') \xrightarrow{(f', \varphi')} (D', Y')$$

$$\begin{array}{ccc} FvFfX & \xrightarrow{Fv\varphi} & FvY \\ (F\alpha)_X \bullet \downarrow & & \downarrow \bullet \lambda \\ Ff'FuX & \xrightarrow{\Phi} & \\ Ff'\rho \bullet \downarrow & & \downarrow \\ Ff'X' & \xrightarrow{\varphi'} & Y' \end{array}$$

Examples

Let A be a 2-category and $F: A \rightarrow \mathbf{2-Cat}$ a 2-functor. There are several ways to construct a double index functor: “first compose, then apply \mathbb{Q} , and then (optional) restrict”

$$① A \xrightarrow{F} \mathbf{2-Cat} \xrightarrow{\mathbb{V}, \mathbb{Q}, \mathbb{Q}^{op}} \mathbf{DblCat}_v$$

$$② \mathbb{Q}(A) \rightarrow \mathbb{Q}(\mathbf{DblCat}_v)$$

$$③ \text{Restrict to } \mathbb{H}(A) \text{ or } \mathbb{V}(A)$$

- $\int_{\mathbb{V}A} \mathbb{Q}(\mathbb{V} \circ F) = \int_{\mathbb{V}A} \mathbb{V}(\mathbb{V} \circ F) = \mathbb{V} \int_A F$ (Bakovic, Buckley)
- $\int_{\mathbb{Q}A} \mathbb{Q}(\mathbb{Q} \circ F) = \mathbb{Q} \int_A F$
- $\int_{\mathbb{Q}A} \mathbb{Q}(\mathbb{V} \circ F) = \mathbb{Q}^{\{cart\}} \int_A F$ (only the cartesian arrows horizontally)
- $\int_{\mathbb{H}A} \mathbb{Q}(\mathbb{V} \circ F) = \mathbb{E}I(F)$ (Pare)
- $\int_{\mathbb{Q}A} \mathbb{Q}(\mathbb{Q}^{op} \circ F) = F \wr F^{op}$ (Myers)

Factorization

- Any horizontal arrow (f, φ) can be factored as $(A, x) \xrightarrow{(f, id)} (B, Ffx) \xrightarrow{(id, \varphi)} (B, y)$.
- Any vertical arrow (u, ρ) can be factored as $(A, x) \xrightarrow{(u, id)} (A', Fux) \xrightarrow{(id, \rho)} (A', x')$.
- And any double cell (α, Φ) can be factored as

$$\begin{array}{ccccc} (A, x) & \xrightarrow{(f, id)} & (B, Ffx) & \xrightarrow{(id, \varphi)} & (B, y) \\ \downarrow (u, id) & & \downarrow (v, id) & & \downarrow (v, id) \\ & (\alpha, id) & (B', FvFfx) & \xrightarrow{(id, Fv\varphi)} & (B', Fvy) \\ (A', Fux) & \xrightarrow{(f', id)} & (B', Ff'Fux) & \xrightarrow{(id, \Phi)} & (B', y') \\ \downarrow (id, \rho) & & \downarrow (id, Ff'\rho) & & \downarrow (id, \lambda) \\ (A', x') & \xrightarrow{(f', id)} & (B', Ff'x') & \xrightarrow{(id, \varphi')} & (B', y') \end{array}$$

$\mathbb{G}r F$ as a double lax cone

For $F: \mathbb{D} \rightarrow \mathbf{DblCat}$ a vertical double functor, construct $F \xrightarrow{\lambda} \Delta \mathbb{G}r F$ as follows:

- Double functors $FA \xrightarrow{\lambda_A} \mathbb{G}r F$: $\lambda_A(-) = (A, -)$.
- Vertical transformations $\lambda_A = \bullet \xrightarrow{\lambda_u} \lambda_{A'} Fu$: for each $x \in FA$, resp.
 $x \xrightarrow{\varphi} y$:

$$\begin{array}{ccc} (A, x) & & (A, x) \xrightarrow{(id, \varphi)} (A, y) \\ \downarrow \bullet (\lambda_u)_x = (u, id), & & (u, id) \bullet \downarrow (\lambda_u)_\varphi = (id, id) \quad \bullet (u, id) \\ (A', F\varphi x) & & (A', Fux) \xrightarrow{(id, Fu\varphi)} (A', Fuy) \end{array}$$

Gr F as a double lax cone

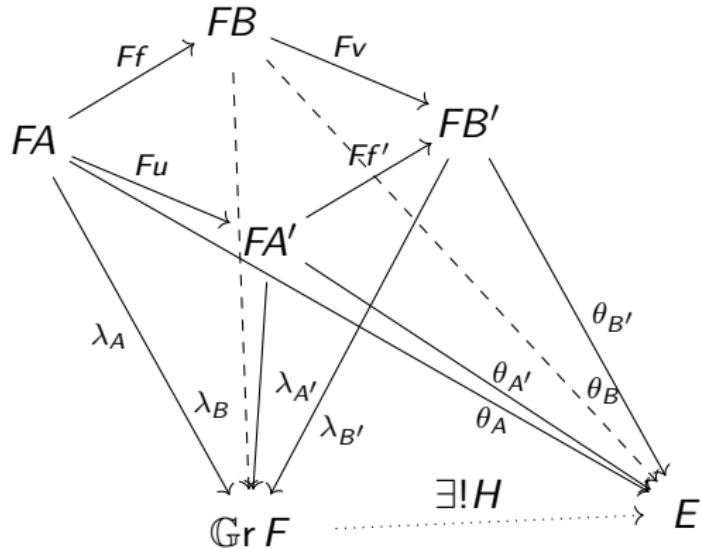
- Hor. transformations $\lambda_A \xrightarrow{\lambda_f} \lambda_B Ff$: for each $x \in FA$, resp. $x \xrightarrow{\rho} x'$:

$$\begin{array}{ccc}
 (A, x) & \xrightarrow{(f, id)} & (B, Ffx) \\
 (A, x) & \xrightarrow{(\lambda_f)_x = (f, id)} & (B, Ffx), \quad (\lambda_f)_\rho = (\text{id}, \text{id}) \\
 & \downarrow (id, \rho) & \downarrow (id, Ff\rho) \\
 (A, x') & \xrightarrow{(f, id)} & (B, Ffx')
 \end{array}$$

- The modifications λ_α , given for each $x \in FA$:

$$\begin{array}{ccc}
 \lambda_A & \xrightarrow{\lambda_f} & \lambda_B Ff \\
 \parallel & & \parallel \\
 \lambda_u & \bullet & \lambda_v Ff \\
 & \downarrow & \\
 & \lambda_\alpha & \lambda_{B'} Fv Ff \\
 & & \parallel \\
 & & \lambda_{B'} F\alpha \\
 & & \downarrow \\
 \lambda_{A'} Fu & \xrightarrow{\lambda_{f'} Fu} & \lambda_{B'} Ff' Fu
 \end{array}
 \qquad
 \begin{array}{ccc}
 (A, x) & \xrightarrow{(f, id)} & (B, Ffx) \\
 & \downarrow & \downarrow \\
 (u, id) & \bullet & (\lambda_\alpha)_x = (\alpha, id) \quad (B', Fv Ffx) \\
 & \downarrow & \\
 & & (\lambda_\alpha)_\rho = (\alpha, id) \\
 & & \downarrow \\
 (A', Fux) & \xrightarrow{(f', id)} & (B', Ff' Fux)
 \end{array}$$

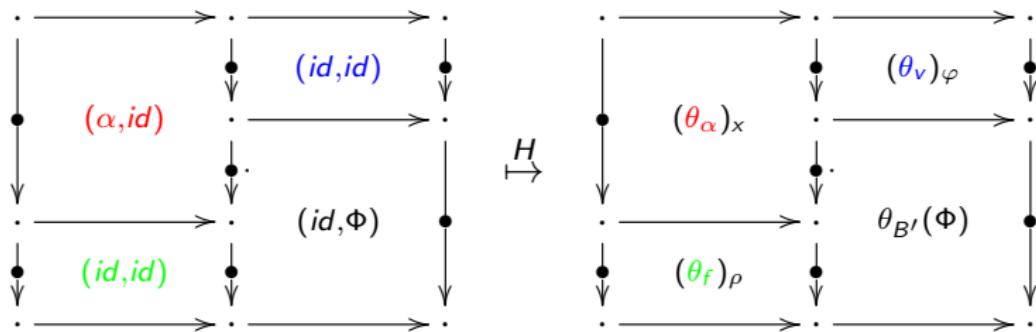
Theorem: $\mathbb{G}r F$ is the double lax colimit of F (in DblCat)



$$H\lambda_A = \theta_A, \quad H\lambda_u = \theta_u, \quad H\lambda_f = \theta_f, \quad H\lambda_\alpha = \theta_\alpha$$

Factorization gives H on double cells

Recall that any double cell (α, Φ) can be factored as



this is why this works

A corollary: Recall that $\int_{\mathbb{V}A} \mathbb{V}(\mathbb{V} \circ F) = \mathbb{V} \int_A F$. We obtain that $\int_A F$ is the *lax tricollector* of F in 2-Cat. Looking at the other examples gives other universal properties of those constructions.

Thank you!